
NOTE: This chapter is an excerpt from the *MSP430x5xx and MSP430x6xx Family User's Guide*. The most recent version of the full user's guide is available at <http://www.ti.com/lit/pdf/slau208>.

This chapter describes the extended MSP430X 16-bit RISC CPU (CPUX) with 1MB memory access, its addressing modes, and instruction set.

NOTE: The MSP430X CPIX implemented on this device family, formally called CPUXV2, has in some cases, slightly different cycle counts from the MSP430X CPIX implemented on the 2xx and 4xx families.

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1.1 MSP430X CPU (CPUX) Introduction

The MSP430X CPU incorporates features specifically designed for modern programming techniques, such as calculated branching, table processing, and the use of high-level languages such as C. The MSP430X CPU can address a 1MB address range without paging. The MSP430X CPU is completely backward compatible with the MSP430 CPU.

The MSP430X CPU features include:

- RISC architecture
- Orthogonal architecture
- Full register access including program counter (PC), status register (SR), and stack pointer (SP)
- Single-cycle register operations
- Large register file reduces fetches to memory.
- 20-bit address bus allows direct access and branching throughout the entire memory range without paging.
- 16-bit data bus allows direct manipulation of word-wide arguments.
- Constant generator provides the six most often used immediate values and reduces code size.
- Direct memory-to-memory transfers without intermediate register holding
- Byte, word, and 20-bit address-word addressing

The block diagram of the MSP430X CPU is shown in [Figure 1-1](#).

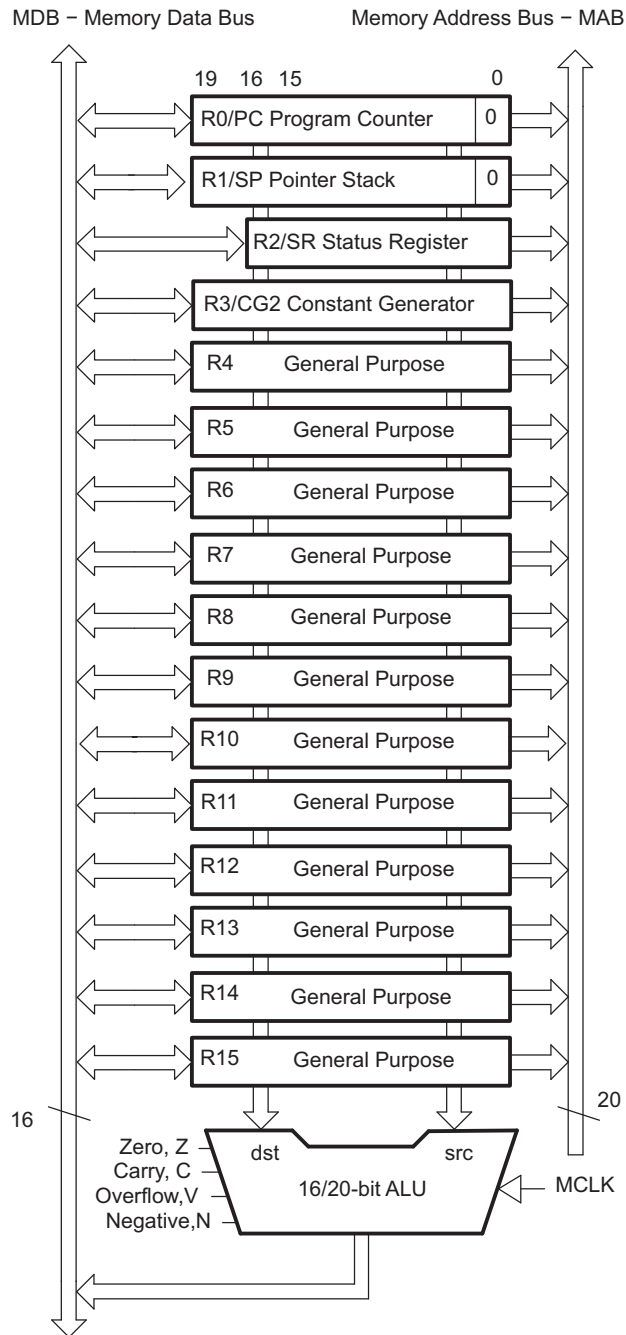


Figure 1-1. MSP430X CPU Block Diagram

1.2 Interrupts

The MSP430X has the following interrupt structure:

- Vectored interrupts with no polling necessary
- Interrupt vectors are located downward from address 0FFFEh.

The interrupt vectors contain 16-bit addresses that point into the lower 64KB memory. This means all interrupt handlers must start in the lower 64KB memory.

During an interrupt, the program counter (PC) and the status register (SR) are pushed onto the stack as shown in [Figure 1-2](#). The MSP430X architecture stores the complete 20-bit PC value efficiently by appending the PC bits 19:16 to the stored SR value automatically on the stack. When the RETI instruction is executed, the full 20-bit PC is restored making return from interrupt to any address in the memory range possible.

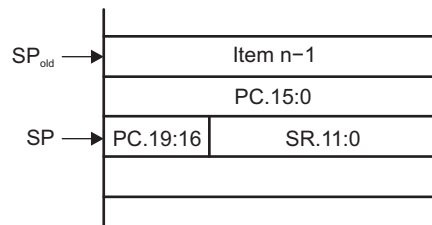


Figure 1-2. PC Storage on the Stack for Interrupts

1.3 CPU Registers

The CPU incorporates 16 registers (R0 through R15). Registers R0, R1, R2, and R3 have dedicated functions. Registers R4 through R15 are working registers for general use.

1.3.1 Program Counter (PC)

The 20-bit Program Counter (PC, also called R0) points to the next instruction to be executed. Each instruction uses an even number of bytes (2, 4, 6, or 8 bytes), and the PC is incremented accordingly. Instruction accesses are performed on word boundaries, and the PC is aligned to even addresses. Figure 1-3 shows the PC.

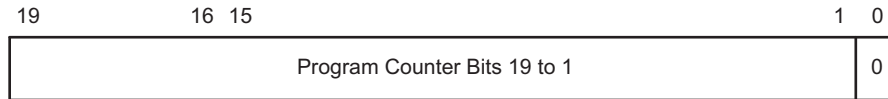


Figure 1-3. Program Counter

The PC can be addressed with all instructions and addressing modes. A few examples:

```
MOV.W #LABEL,PC ; Branch to address LABEL (lower 64KB)

MOVA #LABEL,PC ; Branch to address LABEL (1MB memory)

MOV.W LABEL,PC ; Branch to address in word LABEL
                ; (lower 64KB)

MOV.W @R14,PC ; Branch indirect to address in
                ; R14 (lower 64KB)

ADDA #4,PC ; Skip two words (1MB memory)
```

The BR and CALL instructions reset the upper four PC bits to 0. Only addresses in the lower 64KB address range can be reached with the BR or CALL instruction. When branching or calling, addresses beyond the lower 64KB range can only be reached using the BRA or CALLA instructions. Also, any instruction to directly modify the PC does so according to the used addressing mode. For example, MOV.W #value,PC clears the upper four bits of the PC, because it is a .W instruction.

The PC is automatically stored on the stack with CALL (or CALLA) instructions and during an interrupt service routine. Figure 1-4 shows the storage of the PC with the return address after a CALLA instruction. A CALL instruction stores only bits 15:0 of the PC.

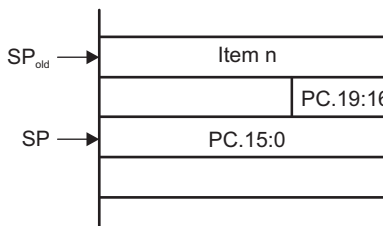


Figure 1-4. PC Storage on the Stack for CALLA

The RETA instruction restores bits 19:0 of the PC and adds 4 to the stack pointer (SP). The RET instruction restores bits 15:0 to the PC and adds 2 to the SP.

1.3.2 Stack Pointer (SP)

The 20-bit Stack Pointer (SP, also called R1) is used by the CPU to store the return addresses of subroutine calls and interrupts. It uses a predecrement, postincrement scheme. In addition, the SP can be used by software with all instructions and addressing modes. Figure 1-5 shows the SP. The SP is initialized into RAM by the user, and is always aligned to even addresses.

Figure 1-6 shows the stack usage. Figure 1-7 shows the stack usage when 20-bit address words are pushed.

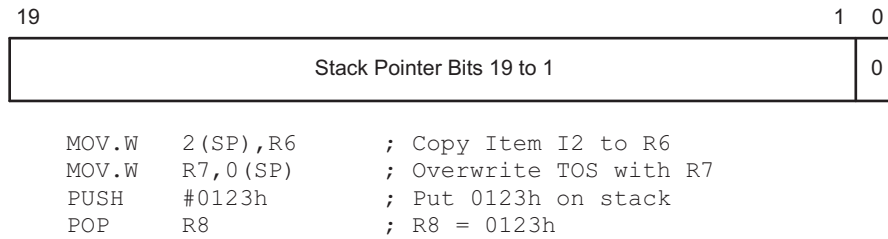


Figure 1-5. Stack Pointer

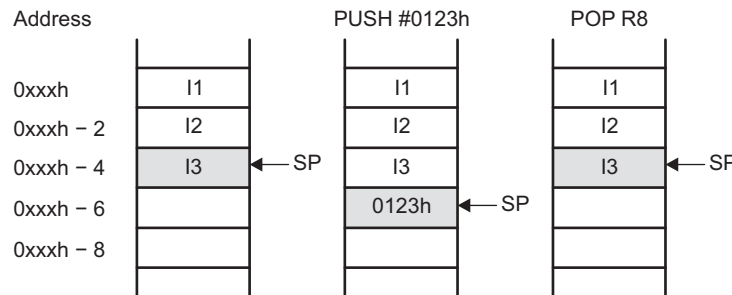


Figure 1-6. Stack Usage

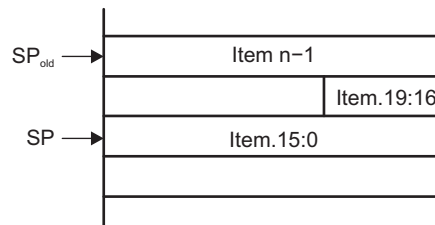


Figure 1-7. PUSHX.A Format on the Stack

The special cases of using the SP as an argument to the PUSH and POP instructions are described and shown in Figure 1-8.

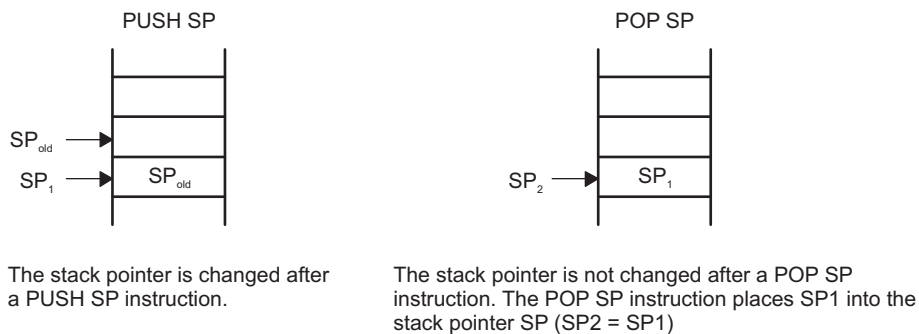


Figure 1-8. PUSH SP, POP SP Sequence

1.3.3 Status Register (SR)

The 16-bit Status Register (SR, also called R2), used as a source or destination register, can only be used in register mode addressed with word instructions. The remaining combinations of addressing modes are used to support the constant generator. Figure 1-9 shows the SR bits. Do not write 20-bit values to the SR. Unpredictable operation can result.

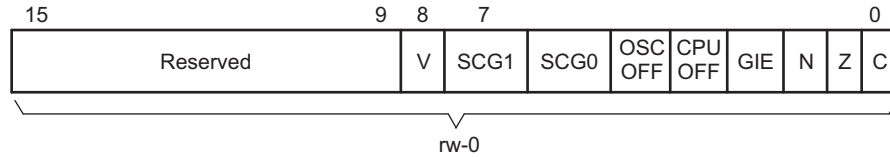


Figure 1-9. SR Bits

Table 1-1 describes the SR bits.

Table 1-1. SR Bit Description

Bit	Description
Reserved	Reserved
V	<p>Overflow. This bit is set when the result of an arithmetic operation overflows the signed-variable range.</p> <p>ADD(.B), ADDX(.B,.A), ADDC(.B), ADDCX(.B.A), ADDA</p> <p>Set when: positive + positive = negative negative + negative = positive otherwise reset</p> <p>SUB(.B), SUBX(.B,.A), SUBC(.B), SUBCX(.B,.A), SUBA, CMP(.B), CMPX(.B,.A), CMPA</p> <p>Set when: positive – negative = negative negative – positive = positive otherwise reset</p>
SCG1	System clock generator 1. This bit may be used to enable or disable functions in the clock system depending on the device family; for example, DCO bias enable or disable.
SCG0	System clock generator 0. This bit may be used to enable or disable functions in the clock system depending on the device family; for example, FLL enable or disable.
OSCOFF	Oscillator off. When this bit is set, it turns off the LFXT1 crystal oscillator when LFXT1CLK is not used for MCLK or SMCLK.
CPUOFF	CPU off. When this bit is set, it turns off the CPU.
GIE	General interrupt enable. When this bit is set, it enables maskable interrupts. When it is reset, all maskable interrupts are disabled.
N	Negative. This bit is set when the result of an operation is negative and cleared when the result is positive.
Z	Zero. This bit is set when the result of an operation is 0 and cleared when the result is not 0.
C	Carry. This bit is set when the result of an operation produced a carry and cleared when no carry occurred.

NOTE: Bit manipulations of the SR should be done by the following instructions: MOV, BIS, and BIC.

1.3.4 Constant Generator Registers (CG1 and CG2)

Six commonly-used constants are generated with the constant generator registers R2 (CG1) and R3 (CG2), without requiring an additional 16-bit word of program code. The constants are selected with the source register addressing modes (As), as described in [Table 1-2](#).

Table 1-2. Values of Constant Generators CG1, CG2

Register	As	Constant	Remarks
R2	00	–	Register mode
R2	01	(0)	Absolute address mode
R2	10	00004h	+4, bit processing
R2	11	00008h	+8, bit processing
R3	00	00000h	0, word processing
R3	01	00001h	+1
R3	10	00002h	+2, bit processing
R3	11	FFh, FFFFh, FFFFFh	–1, word processing

The constant generator advantages are:

- No special instructions required
- No additional code word for the six constants
- No code memory access required to retrieve the constant

The assembler uses the constant generator automatically if one of the six constants is used as an immediate source operand. Registers R2 and R3, used in the constant mode, cannot be addressed explicitly; they act as source-only registers.

1.3.4.1 Constant Generator – Expanded Instruction Set

The RISC instruction set of the MSP430 has only 27 instructions. However, the constant generator allows the MSP430 assembler to support 24 additional emulated instructions. For example, the single-operand instruction:

```
CLR dst
```

is emulated by the double-operand instruction with the same length:

```
MOV R3, dst
```

where the #0 is replaced by the assembler, and R3 is used with As = 00.

```
INC dst
```

is replaced by:

```
ADD #1, dst
```


1.3.5 General-Purpose Registers (R4 to R15)

The 12 CPU registers (R4 to R15) contain 8-bit, 16-bit, or 20-bit values. Any byte-write to a CPU register clears bits 19:8. Any word-write to a register clears bits 19:16. The only exception is the SXT instruction. The SXT instruction extends the sign through the complete 20-bit register.

Figure 1-10 through Figure 1-14 show the handling of byte, word, and address-word data. Note the reset of the leading most significant bits (MSBs) if a register is the destination of a byte or word instruction.

Figure 1-10 shows byte handling (8-bit data, .B suffix). The handling is shown for a source register and a destination memory byte and for a source memory byte and a destination register.

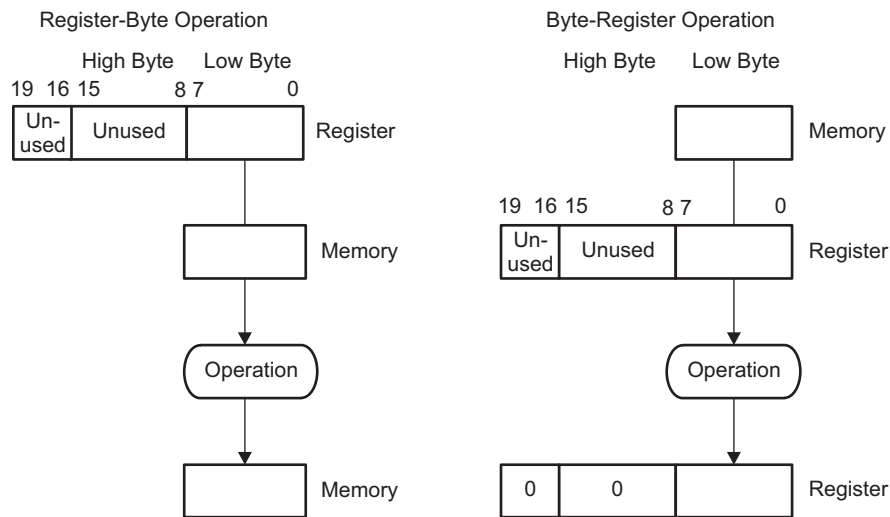


Figure 1-10. Register-Byte and Byte-Register Operation

Figure 1-11 and Figure 1-12 show 16-bit word handling (.W suffix). The handling is shown for a source register and a destination memory word and for a source memory word and a destination register.

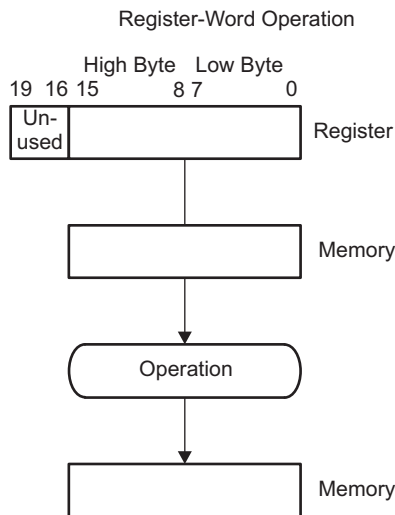


Figure 1-11. Register-Word Operation

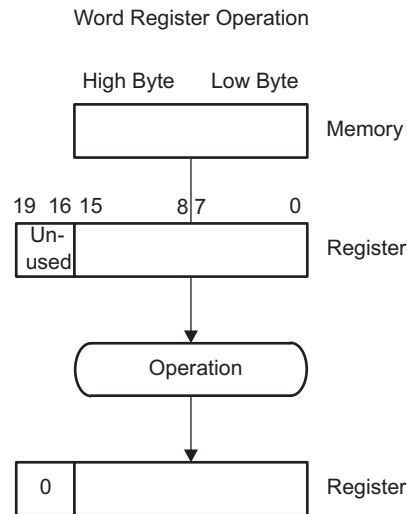


Figure 1-12. Word-Register Operation

Figure 1-13 and Figure 1-14 show 20-bit address-word handling (.A suffix). The handling is shown for a source register and a destination memory address-word and for a source memory address-word and a destination register.

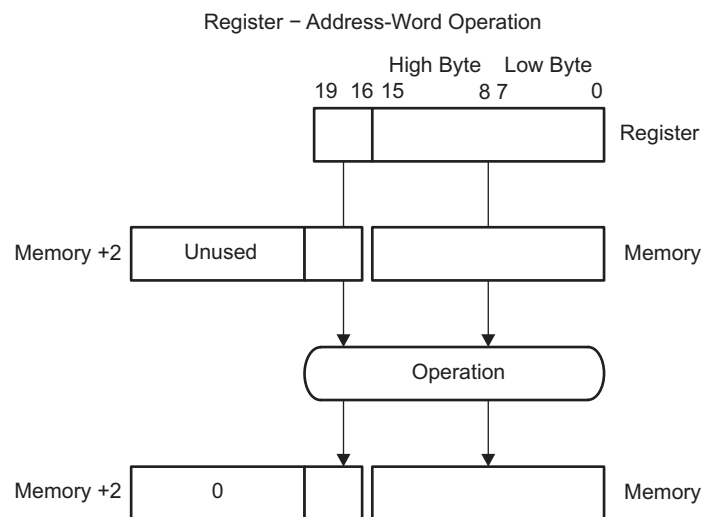


Figure 1-13. Register – Address-Word Operation

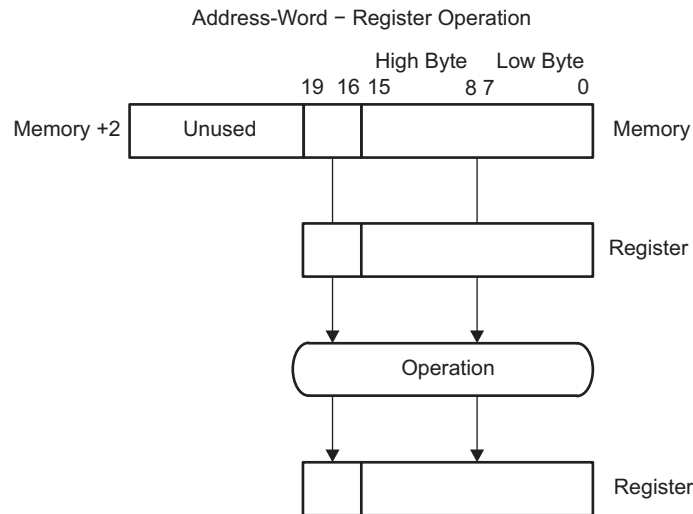


Figure 1-14. Address-Word - Register Operation

1.4 Addressing Modes

Seven addressing modes for the source operand and four addressing modes for the destination operand use 16-bit or 20-bit addresses (see Table 1-3). The MSP430 and MSP430X instructions are usable throughout the entire 1MB memory range.

Table 1-3. Source and Destination Addressing

As, Ad	Addressing Mode	Syntax	Description
00, 0	Register	Rn	Register contents are operand.
01, 1	Indexed	X(Rn)	(Rn + X) points to the operand. X is stored in the next word, or stored in combination of the preceding extension word and the next word.
01, 1	Symbolic	ADDR	(PC + X) points to the operand. X is stored in the next word, or stored in combination of the preceding extension word and the next word. Indexed mode X(PC) is used.
01, 1	Absolute	&ADDR	The word following the instruction contains the absolute address. X is stored in the next word, or stored in combination of the preceding extension word and the next word. Indexed mode X(SR) is used.
10, -	Indirect Register	@Rn	Rn is used as a pointer to the operand.
11, -	Indirect Autoincrement	@Rn+	Rn is used as a pointer to the operand. Rn is incremented afterwards by 1 for .B instructions, by 2 for .W instructions, and by 4 for .A instructions.
11, -	Immediate	#N	N is stored in the next word, or stored in combination of the preceding extension word and the next word. Indirect autoincrement mode @PC+ is used.

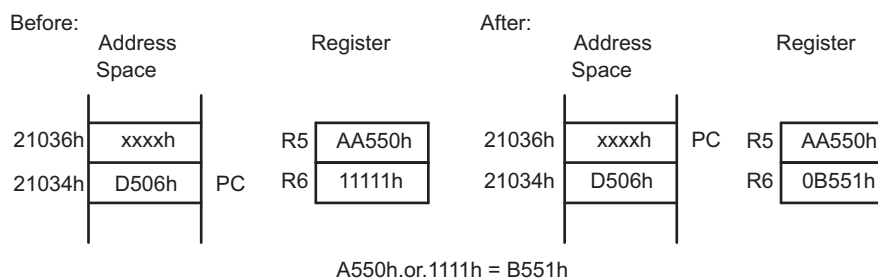
The seven addressing modes are explained in detail in the following sections. Most of the examples show the same addressing mode for the source and destination, but any valid combination of source and destination addressing modes is possible in an instruction.

NOTE: Use of Labels EDE, TONI, TOM, and LEO

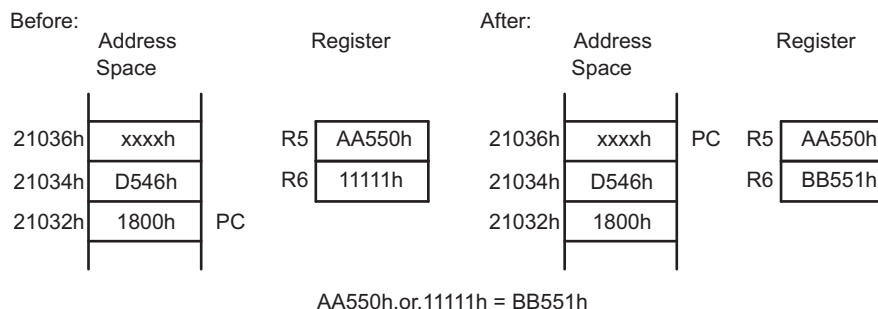
Throughout MSP430 documentation, EDE, TONI, TOM, and LEO are used as generic labels. They are only labels and have no special meaning.

1.4.1 Register Mode

Operation:	The operand is the 8-, 16-, or 20-bit content of the used CPU register.
Length:	One, two, or three words
Comment:	Valid for source and destination
Byte operation:	Byte operation reads only the eight least significant bits (LSBs) of the source register Rsrc and writes the result to the eight LSBs of the destination register Rdst. The bits Rdst.19:8 are cleared. The register Rsrc is not modified.
Word operation:	Word operation reads the 16 LSBs of the source register Rsrc and writes the result to the 16 LSBs of the destination register Rdst. The bits Rdst.19:16 are cleared. The register Rsrc is not modified.
Address-word operation:	Address-word operation reads the 20 bits of the source register Rsrc and writes the result to the 20 bits of the destination register Rdst. The register Rsrc is not modified
SXT exception:	The SXT instruction is the only exception for register operation. The sign of the low byte in bit 7 is extended to the bits Rdst.19:8.
Example:	<pre>BIS.W R5,R6 ;</pre> <p>This instruction logically ORs the 16-bit data contained in R5 with the 16-bit contents of R6. R6.19:16 is cleared.</p>



Example:	<pre>BISX.A R5,R6 ;</pre> <p>This instruction logically ORs the 20-bit data contained in R5 with the 20-bit contents of R6.</p> <p>The extension word contains the A/L bit for 20-bit data. The instruction word uses byte mode with bits A/L:B/W = 01. The result of the instruction is:</p>
----------	---



1.4.2 Indexed Mode

The Indexed mode calculates the address of the operand by adding the signed index to a CPU register. The Indexed mode has four addressing possibilities:

- MSP430 instruction with Indexed mode in lower 64KB memory (see [Section 1.4.2.1](#))
- MSP430 instruction with Indexed mode addressing memory above the lower 64KB memory (see [Section 1.4.2.2](#))
- MSP430X instruction with Indexed mode (see [Section 1.4.2.3](#))
- MSP430X address instructions with Indexed mode (see [Section 1.4.2.4](#))

1.4.2.1 MSP430 Instruction With Indexed Mode in Lower 64KB Memory

If the CPU register Rn points to an address in the lower 64KB of the memory range, the calculated memory address bits 19:16 are cleared after the addition of the CPU register Rn and the signed 16-bit index. This means the calculated memory address is always located in the lower 64KB and does not overflow or underflow out of the lower 64KB memory space. The RAM and the peripheral registers can be accessed this way and existing MSP430 software is usable without modifications as shown in [Figure 1-15](#).

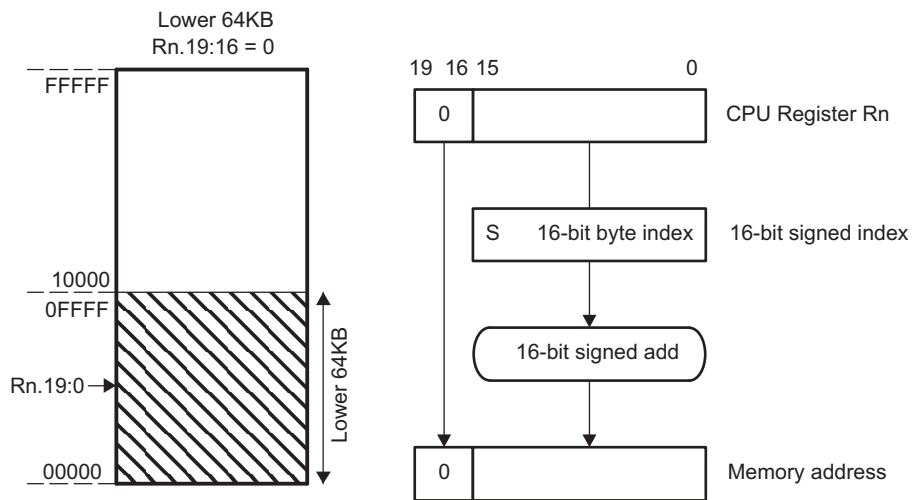
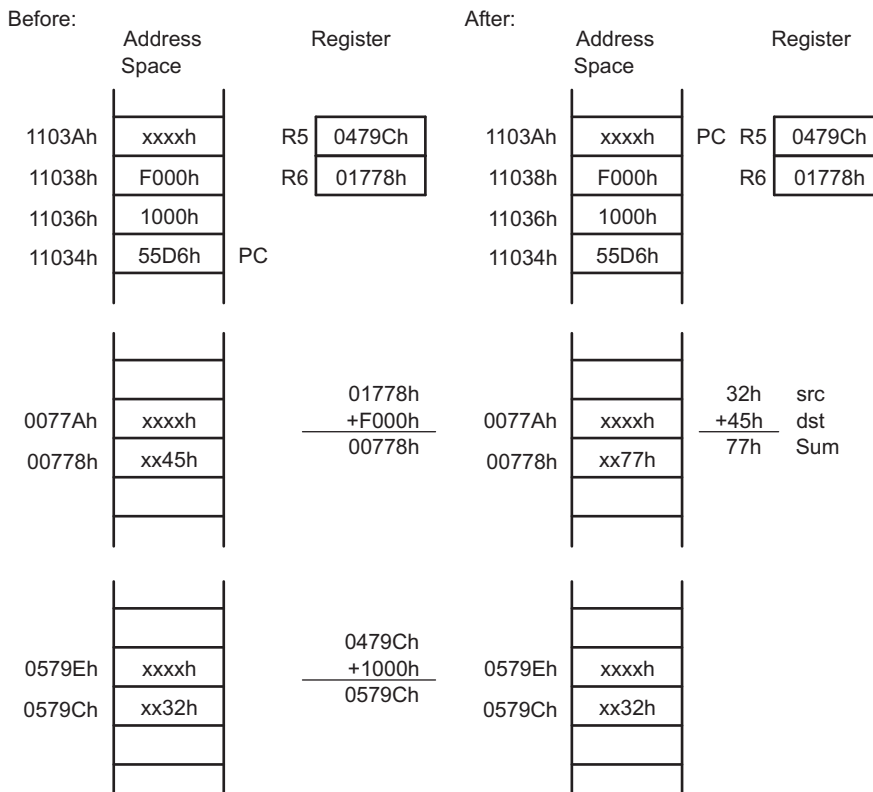


Figure 1-15. Indexed Mode in Lower 64KB

- Length:** Two or three words
- Operation:** The signed 16-bit index is located in the next word after the instruction and is added to the CPU register Rn. The resulting bits 19:16 are cleared giving a truncated 16-bit memory address, which points to an operand address in the range 00000h to 0FFFFh. The operand is the content of the addressed memory location.
- Comment:** Valid for source and destination. The assembler calculates the register index and inserts it.
- Example:** `ADD.B 1000h(R5), 0F000h(R6);`
 This instruction adds the 8-bit data contained in source byte 1000h(R5) and the destination byte 0F000h(R6) and places the result into the destination byte. Source and destination bytes are both located in the lower 64KB due to the cleared bits 19:16 of registers R5 and R6.
- Source:** The byte pointed to by R5 + 1000h results in address 0479Ch + 1000h = 0579Ch after truncation to a 16-bit address.
- Destination:** The byte pointed to by R6 + F000h results in address 01778h + F000h = 00778h after truncation to a 16-bit address.



1.4.2.2 MSP430 Instruction With Indexed Mode in Upper Memory

If the CPU register Rn points to an address above the lower 64KB memory, the Rn bits 19:16 are used for the address calculation of the operand. The operand may be located in memory in the range Rn ±32KB, because the index, X, is a signed 16-bit value. In this case, the address of the operand can overflow or underflow into the lower 64KB memory space (see Figure 1-16 and Figure 1-17).

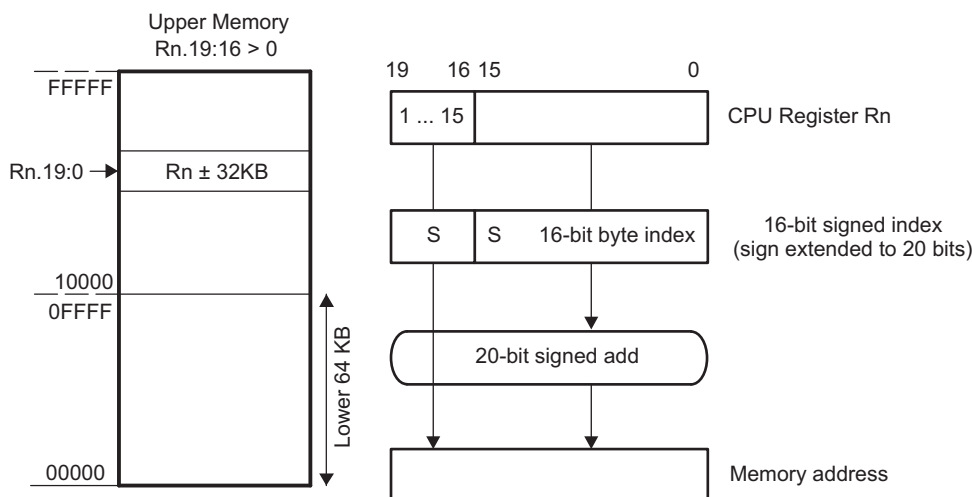


Figure 1-16. Indexed Mode in Upper Memory

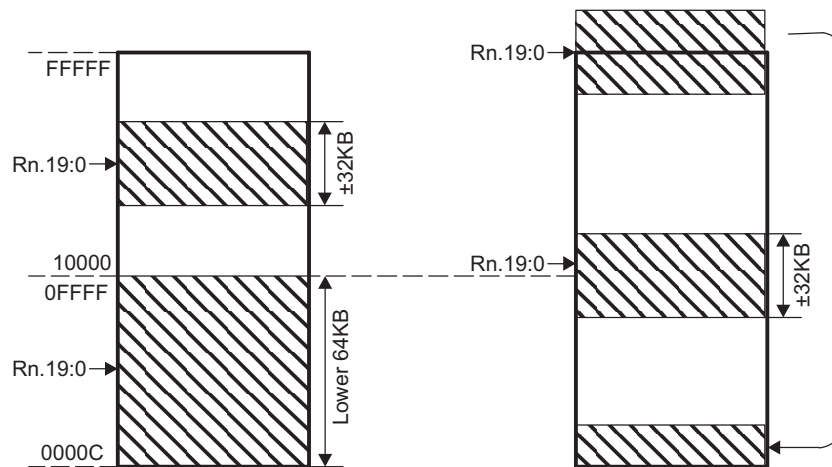


Figure 1-17. Overflow and Underflow for Indexed Mode

- Length:** Two or three words
- Operation:** The sign-extended 16-bit index in the next word after the instruction is added to the 20 bits of the CPU register Rn. This delivers a 20-bit address, which points to an address in the range 0 to FFFFFh. The operand is the content of the addressed memory location.
- Comment:** Valid for source and destination. The assembler calculates the register index and inserts it.
- Example:** `ADD.W 8346h(R5), 2100h(R6) ;`
 This instruction adds the 16-bit data contained in the source and the destination addresses and places the 16-bit result into the destination. Source and destination operand can be located in the entire address range.
- Source:** The word pointed to by R5 + 8346h. The negative index 8346h is sign extended, which results in address 23456h + F8346h = 1B79Ch.
- Destination:** The word pointed to by R6 + 2100h results in address 15678h + 2100h = 17778h.

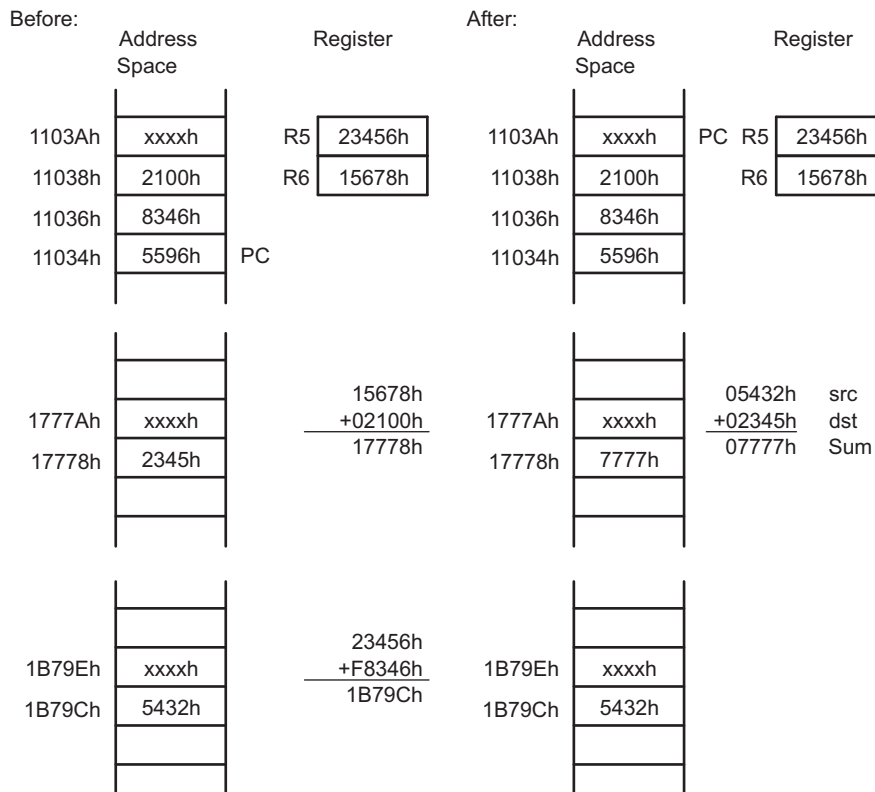


Figure 1-18. Example for Indexed Mode

1.4.2.3 MSP430X Instruction With Indexed Mode

When using an MSP430X instruction with Indexed mode, the operand can be located anywhere in the range of $R_n + 19$ bits.

Length: Three or four words

Operation: The operand address is the sum of the 20-bit CPU register content and the 20-bit index. The 4 MSBs of the index are contained in the extension word; the 16 LSBs are contained in the word following the instruction. The CPU register is not modified

Comment: Valid for source and destination. The assembler calculates the register index and inserts it.

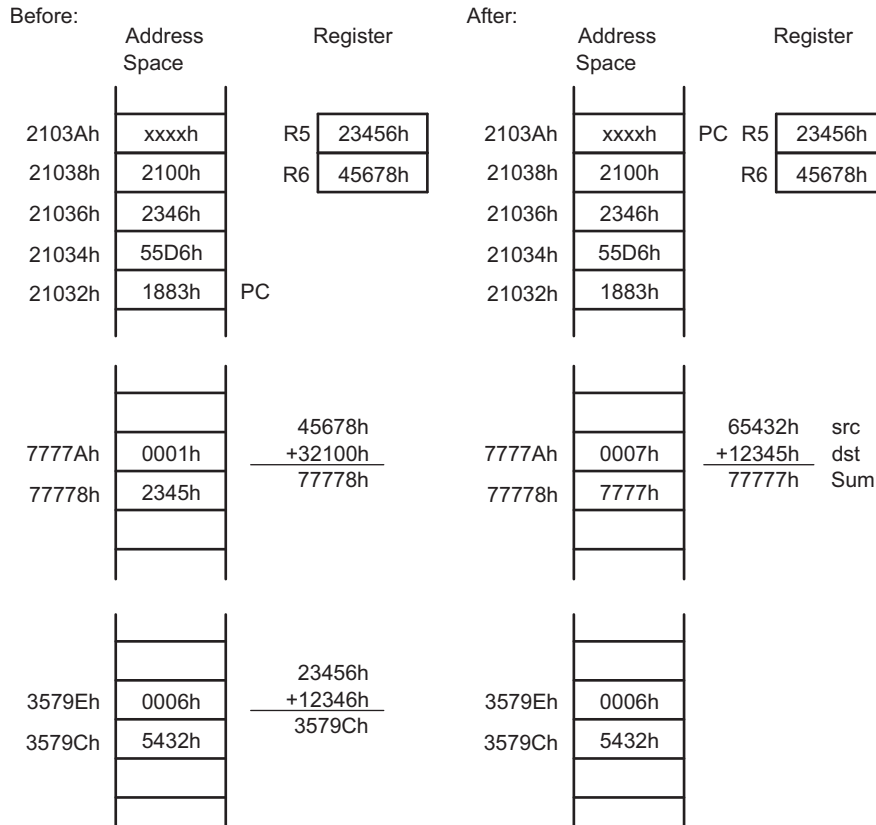
Example: `ADDX.A 12346h(R5), 32100h(R6) ;`

This instruction adds the 20-bit data contained in the source and the destination addresses and places the result into the destination.

Source: Two words pointed to by $R_5 + 12346h$ which results in address $23456h + 12346h = 3579Ch$.

Destination: Two words pointed to by $R_6 + 32100h$ which results in address $45678h + 32100h = 77778h$.

The extension word contains the MSBs of the source index and of the destination index and the A/L bit for 20-bit data. The instruction word uses byte mode due to the 20-bit data length with bits A/L:B/W = 01.



1.4.2.4 MSP430X Address Instructions With Indexed Mode

When using an MSP430X Address Instruction with Indexed mode, the operand is located in memory in the range $R_n \pm 32KB$, because the index, X, is a signed 16-bit value.

Length: Two words

Operation: The sign-extended 16-bit index in the next word after the instruction is added to the 20 bits of the CPU register R_n . This delivers a 20-bit address, which points to an address in the range 0 to FFFFh. The operand is the content of the addressed memory location.

Comment: Valid for source and destination. The assembler calculates the register index and inserts it.

Example: `MOVA 8002h(R5),R6 ; // R5 = 0x100`

This instruction loads the 20-bit data contained in the source address into destination register.

Source: Two words pointed to by $R5 + 8002h$ and $R5 + 8002h + 2h$ which results in address $00100h + F8002h (+2h) = F8102h$ and $F8104h$.

Destination: Register R6

1.4.3 Symbolic Mode

The Symbolic mode calculates the address of the operand by adding the signed index to the PC. The Symbolic mode has three addressing possibilities:

- Symbolic mode in lower 64KB of memory
- MSP430 instruction with Symbolic mode addressing memory above the lower 64KB of memory.
- MSP430X instruction with Symbolic mode

1.4.3.1 Symbolic Mode in Lower 64KB

If the PC points to an address in the lower 64KB of the memory range, the calculated memory address bits 19:16 are cleared after the addition of the PC and the signed 16-bit index. This means the calculated memory address is always located in the lower 64KB and does not overflow or underflow out of the lower 64KB memory space. The RAM and the peripheral registers can be accessed this way and existing MSP430 software is usable without modifications as shown in [Figure 1-19](#).

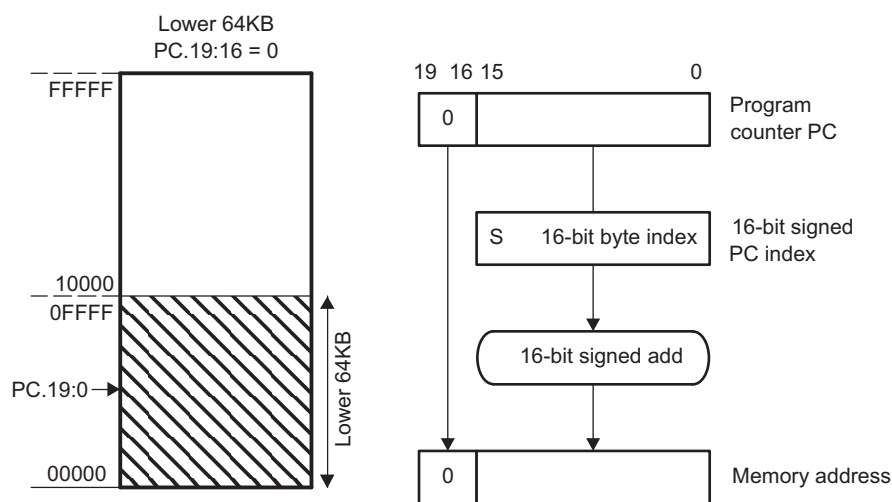
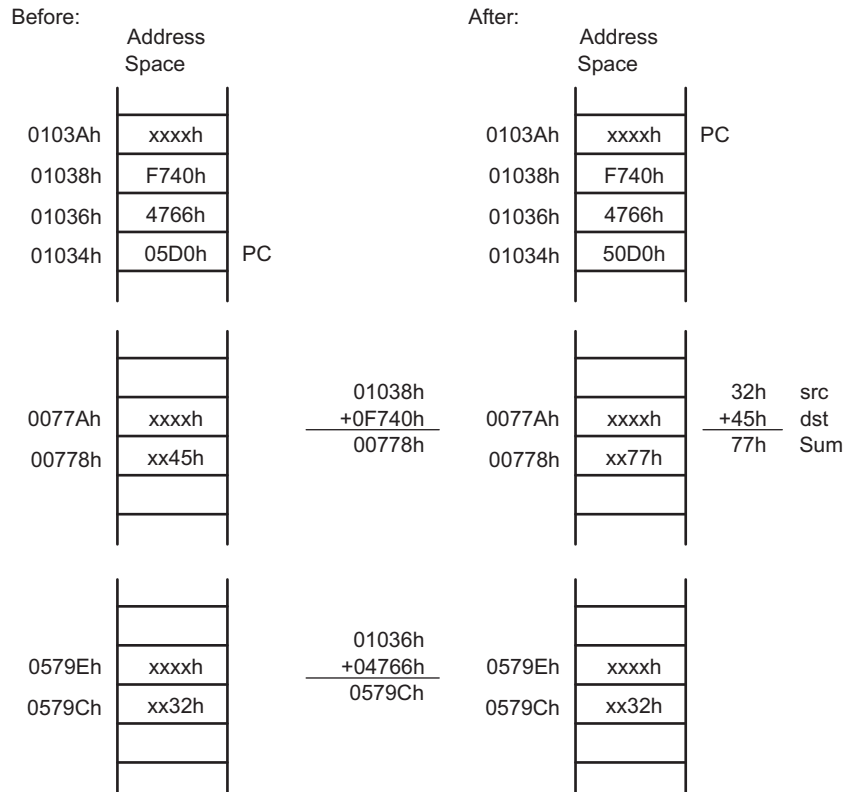


Figure 1-19. Symbolic Mode Running in Lower 64KB

Operation:	The signed 16-bit index in the next word after the instruction is added temporarily to the PC. The resulting bits 19:16 are cleared giving a truncated 16-bit memory address, which points to an operand address in the range 00000h to 0FFFFh. The operand is the content of the addressed memory location.
Length:	Two or three words
Comment:	Valid for source and destination. The assembler calculates the PC index and inserts it.
Example:	<code>ADD.B EDE, TONI ;</code> This instruction adds the 8-bit data contained in source byte EDE and destination byte TONI and places the result into the destination byte TONI. Bytes EDE and TONI and the program are located in the lower 64KB.
Source:	Byte EDE located at address 0579Ch, pointed to by PC + 4766h, where the PC index 4766h is the result of 0579Ch – 01036h = 04766h. Address 01036h is the location of the index for this example.
Destination:	Byte TONI located at address 00778h, pointed to by PC + F740h, is the truncated 16-bit result of 00778h – 1038h = FF740h. Address 01038h is the location of the index for this example.



1.4.3.2 MSP430 Instruction With Symbolic Mode in Upper Memory

If the PC points to an address above the lower 64KB memory, the PC bits 19:16 are used for the address calculation of the operand. The operand may be located in memory in the range PC ± 32KB, because the index, X, is a signed 16-bit value. In this case, the address of the operand can overflow or underflow into the lower 64KB memory space as shown in Figure 1-20 and Figure 1-21.

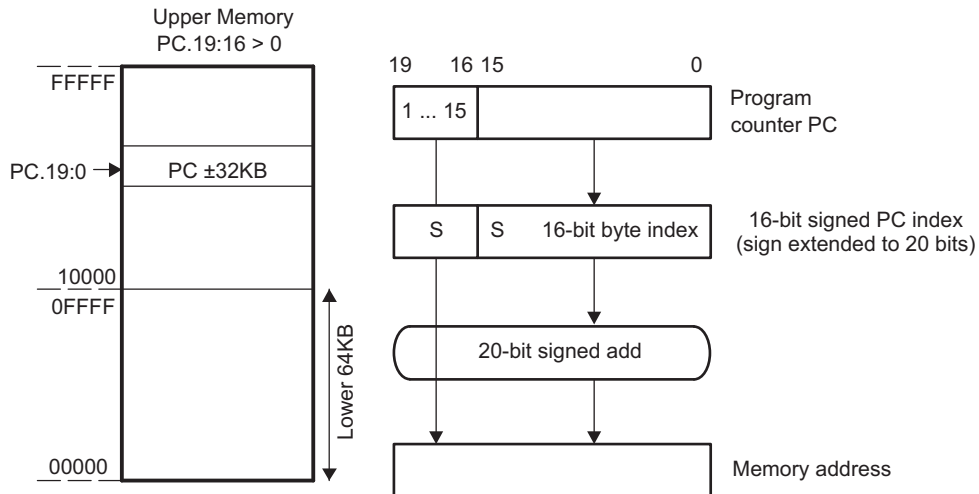


Figure 1-20. Symbolic Mode Running in Upper Memory

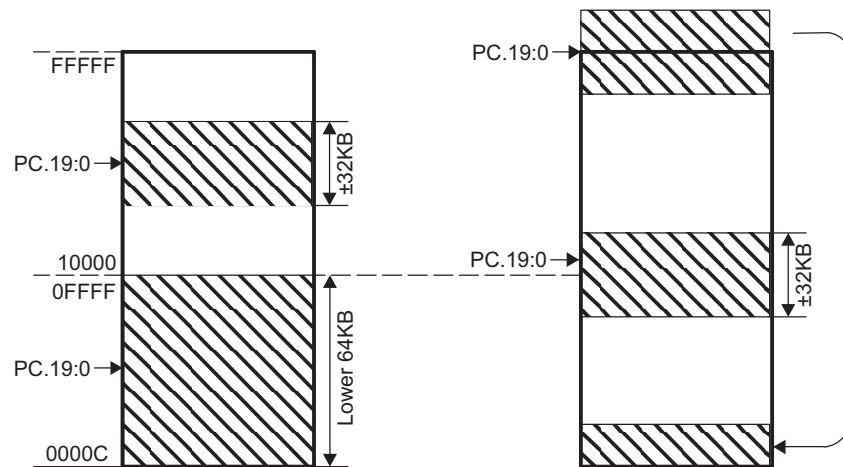
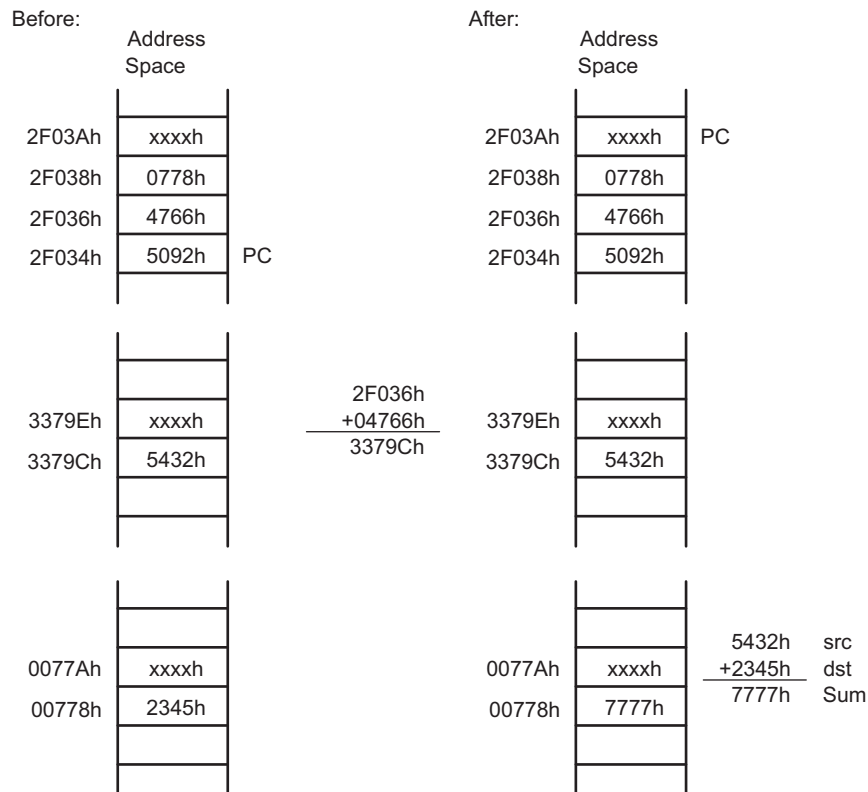


Figure 1-21. Overflow and Underflow for Symbolic Mode

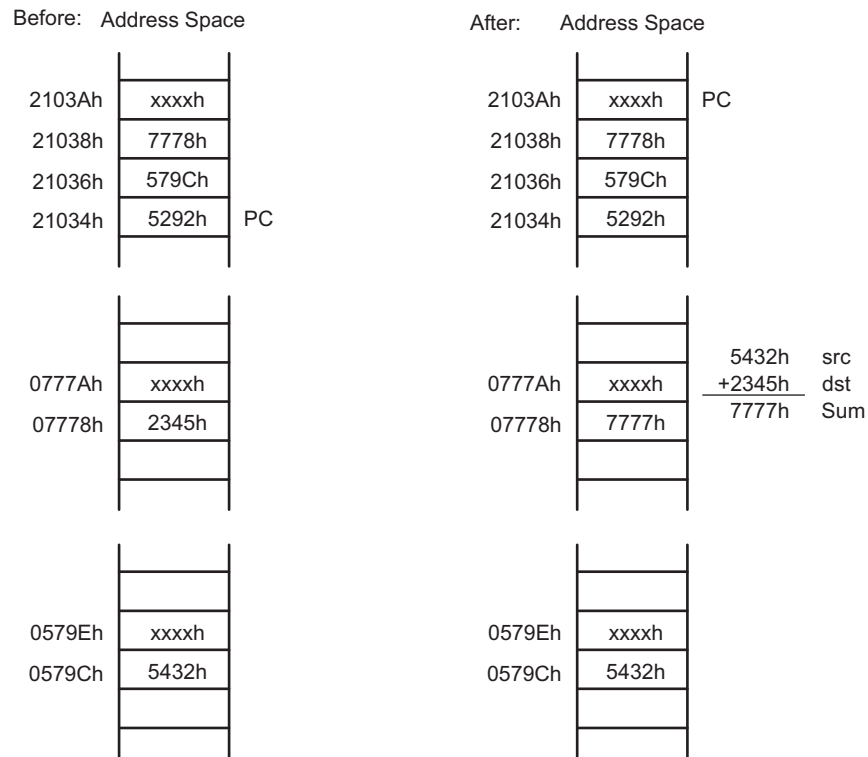
Length:	Two or three words
Operation:	The sign-extended 16-bit index in the next word after the instruction is added to the 20 bits of the PC. This delivers a 20-bit address, which points to an address in the range 0 to FFFFh. The operand is the content of the addressed memory location.
Comment:	Valid for source and destination. The assembler calculates the PC index and inserts it
Example:	<pre>ADD.W EDE, &TONI ;</pre> <p>This instruction adds the 16-bit data contained in source word EDE and destination word TONI and places the 16-bit result into the destination word TONI. For this example, the instruction is located at address 2F034h.</p>
Source:	Word EDE at address 3379Ch, pointed to by PC + 4766h, which is the 16-bit result of 3379Ch – 2F036h = 04766h. Address 2F036h is the location of the index for this example.
Destination:	Word TONI located at address 00778h pointed to by the absolute address 00778h



1.4.3.3 MSP430X Instruction With Symbolic Mode

When using an MSP430X instruction with Symbolic mode, the operand can be located anywhere in the range of PC + 19 bits.

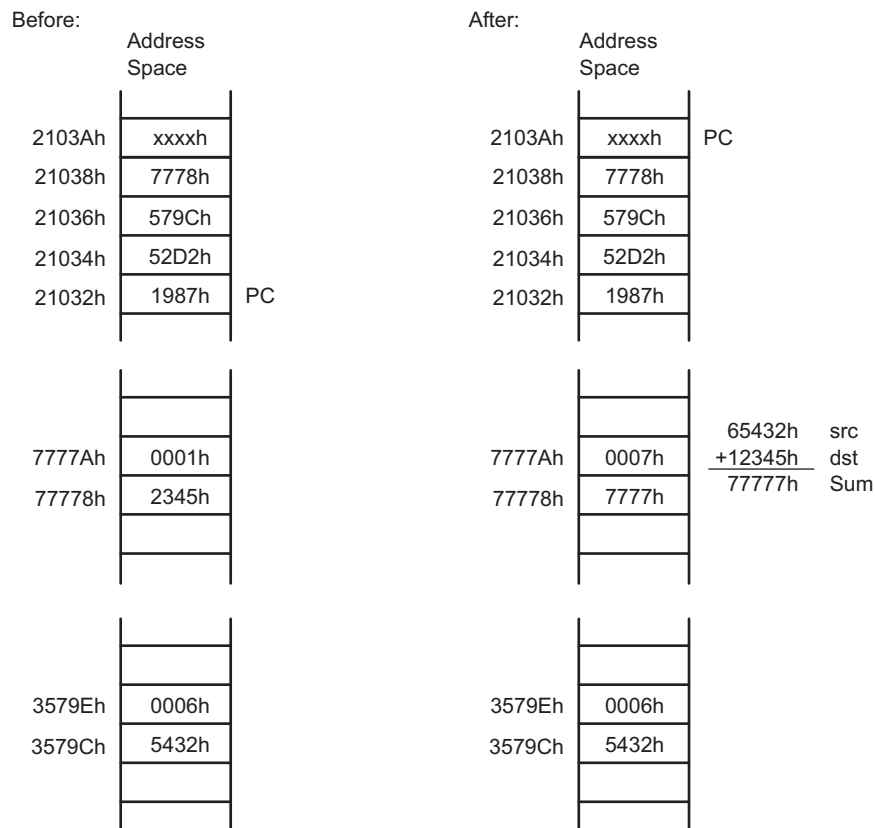
- Length:** Three or four words
- Operation:** The operand address is the sum of the 20-bit PC and the 20-bit index. The 4 MSBs of the index are contained in the extension word; the 16 LSBs are contained in the word following the instruction.
- Comment:** Valid for source and destination. The assembler calculates the register index and inserts it.
- Example:** `ADDX.B EDE,TONI ;`
 This instruction adds the 8-bit data contained in source byte EDE and destination byte TONI and places the result into the destination byte TONI.
- Source:** Byte EDE located at address 3579Ch, pointed to by PC + 14766h, is the 20-bit result of 3579Ch – 21036h = 14766h. Address 21036h is the address of the index in this example.
- Destination:** Byte TONI located at address 77778h, pointed to by PC + 56740h, is the 20-bit result of 77778h – 21038h = 56740h. Address 21038h is the address of the index in this example.



1.4.4.2 MSP430X Instruction With Absolute Mode

If an MSP430X instruction is used with Absolute addressing mode, the absolute address is a 20-bit value and, therefore, points to any address in the memory range. The address value is calculated as an index from 0. The 4 MSBs of the index are contained in the extension word, and the 16 LSBs are contained in the word following the instruction.

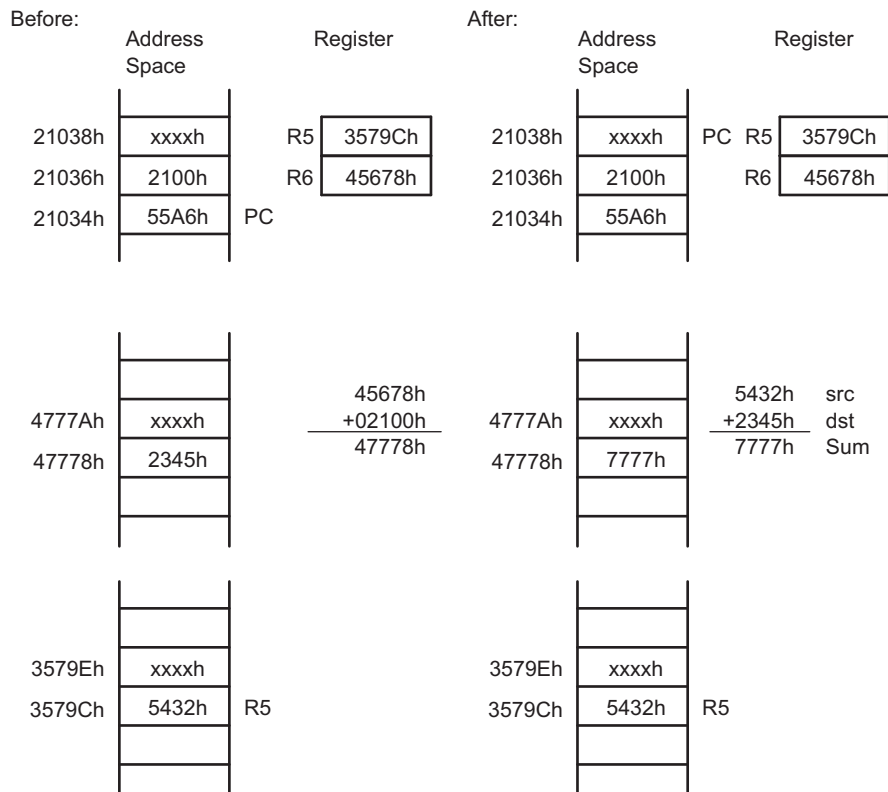
Length:	Three or four words
Operation:	The operand is the content of the addressed memory location.
Comment:	Valid for source and destination. The assembler calculates the index from 0 and inserts it.
Example:	<pre>ADDX.A &EDE, &TONI ;</pre> <p>This instruction adds the 20-bit data contained in the absolute source and destination addresses and places the result into the destination.</p>
Source:	Two words beginning with address EDE
Destination:	Two words beginning with address TONI



1.4.5 Indirect Register Mode

The Indirect Register mode uses the contents of the CPU register Rsrc as the source operand. The Indirect Register mode always uses a 20-bit address.

Length:	One, two, or three words
Operation:	The operand is the content the addressed memory location. The source register Rsrc is not modified.
Comment:	Valid only for the source operand. The substitute for the destination operand is 0(Rdst).
Example:	<p>ADDX.W @R5, 2100h(R6)</p> <p>This instruction adds the two 16-bit operands contained in the source and the destination addresses and places the result into the destination.</p>
Source:	Word pointed to by R5. R5 contains address 3579Ch for this example.
Destination:	Word pointed to by R6 + 2100h, which results in address 45678h + 2100h = 7778h



1.4.6 Indirect Autoincrement Mode

The Indirect Autoincrement mode uses the contents of the CPU register Rsrc as the source operand. Rsrc is then automatically incremented by 1 for byte instructions, by 2 for word instructions, and by 4 for address-word instructions immediately after accessing the source operand. If the same register is used for source and destination, it contains the incremented address for the destination access. Indirect Autoincrement mode always uses 20-bit addresses.

Length: One, two, or three words

Operation: The operand is the content of the addressed memory location.

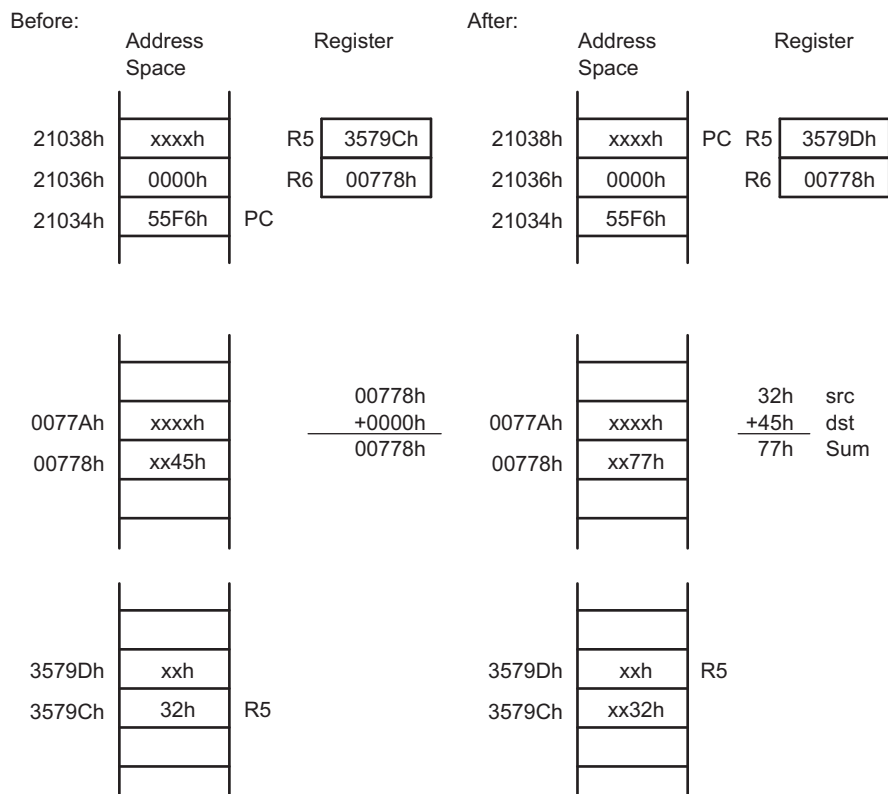
Comment: Valid only for the source operand

Example: `ADD.B @R5+,0(R6)`

This instruction adds the 8-bit data contained in the source and the destination addresses and places the result into the destination.

Source: Byte pointed to by R5. R5 contains address 3579Ch for this example.

Destination: Byte pointed to by R6 + 0h, which results in address 0778h for this example



1.4.7 Immediate Mode

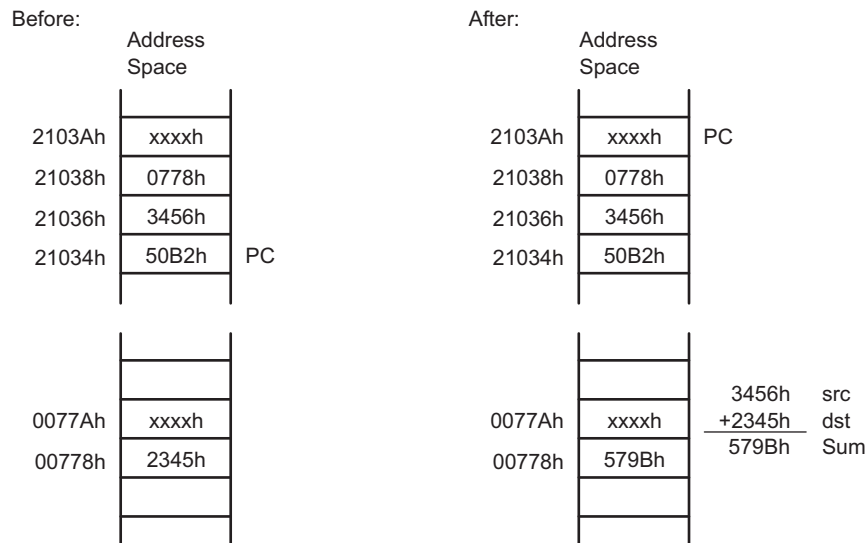
The Immediate mode allows accessing constants as operands by including the constant in the memory location following the instruction. The PC is used with the Indirect Autoincrement mode. The PC points to the immediate value contained in the next word. After the fetching of the immediate operand, the PC is incremented by 2 for byte, word, or address-word instructions. The Immediate mode has two addressing possibilities:

- 8-bit or 16-bit constants with MSP430 instructions
- 20-bit constants with MSP430X instruction

1.4.7.1 MSP430 Instructions With Immediate Mode

If an MSP430 instruction is used with Immediate addressing mode, the constant is an 8- or 16-bit value and is stored in the word following the instruction.

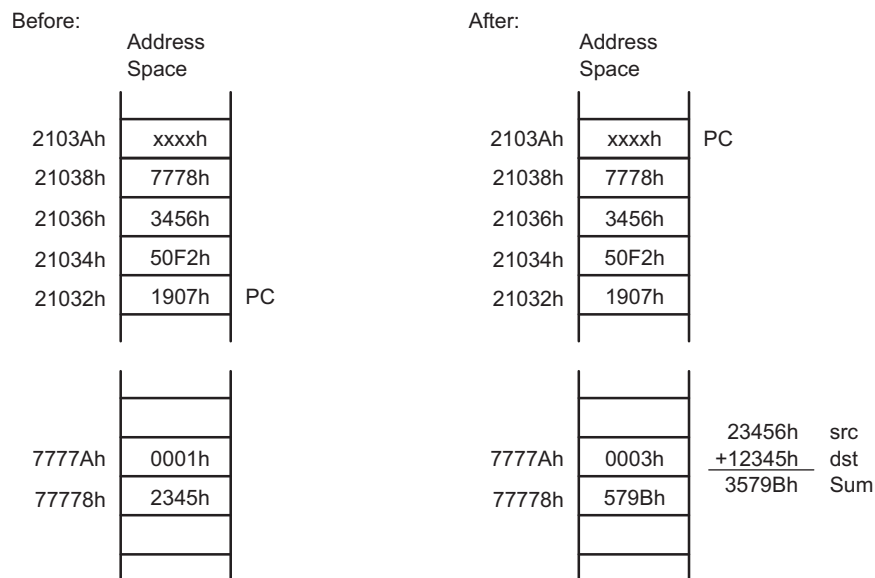
Length:	Two or three words. One word less if a constant of the constant generator can be used for the immediate operand.
Operation:	The 16-bit immediate source operand is used together with the 16-bit destination operand.
Comment:	Valid only for the source operand
Example:	ADD #3456h, &TONI This instruction adds the 16-bit immediate operand 3456h to the data in the destination address TONI.
Source:	16-bit immediate value 3456h
Destination:	Word at address TONI



1.4.7.2 MSP430X Instructions With Immediate Mode

If an MSP430X instruction is used with Immediate addressing mode, the constant is a 20-bit value. The 4 MSBs of the constant are stored in the extension word, and the 16 LSBs of the constant are stored in the word following the instruction.

Length:	Three or four words. One word less if a constant of the constant generator can be used for the immediate operand.
Operation:	The 20-bit immediate source operand is used together with the 20-bit destination operand.
Comment:	Valid only for the source operand
Example:	<p>ADDX.A #23456h, &TONI ;</p> <p>This instruction adds the 20-bit immediate operand 23456h to the data in the destination address TONI.</p>
Source:	20-bit immediate value 23456h
Destination:	Two words beginning with address TONI



1.5 MSP430 and MSP430X Instructions

MSP430 instructions are the 27 implemented instructions of the MSP430 CPU. These instructions are used throughout the 1MB memory range unless their 16-bit capability is exceeded. The MSP430X instructions are used when the addressing of the operands or the data length exceeds the 16-bit capability of the MSP430 instructions.

There are three possibilities when choosing between an MSP430 and MSP430X instruction:

- To use only the MSP430 instructions – The only exceptions are the CALLA and the RETA instruction. This can be done if a few, simple rules are met:
 - Place all constants, variables, arrays, tables, and data in the lower 64KB. This allows the use of MSP430 instructions with 16-bit addressing for all data accesses. No pointers with 20-bit addresses are needed.
 - Place subroutine constants immediately after the subroutine code. This allows the use of the symbolic addressing mode with its 16-bit index to reach addresses within the range of PC + 32KB.
- To use only MSP430X instructions – The disadvantages of this method are the reduced speed due to the additional CPU cycles and the increased program space due to the necessary extension word for any double-operand instruction.
- Use the best fitting instruction where needed.

[Section 1.5.1](#) lists and describes the MSP430 instructions, and [Section 1.5.2](#) lists and describes the MSP430X instructions.

1.5.1 MSP430 Instructions

The MSP430 instructions can be used, regardless if the program resides in the lower 64KB or beyond it. The only exceptions are the instructions CALL and RET, which are limited to the lower 64KB address range. CALLA and RETA instructions have been added to the MSP430X CPU to handle subroutines in the entire address range with no code size overhead.

1.5.1.1 MSP430 Double-Operand (Format I) Instructions

[Figure 1-22](#) shows the format of the MSP430 double-operand instructions. Source and destination words are appended for the Indexed, Symbolic, Absolute, and Immediate modes. [Table 1-4](#) lists the 12 MSP430 double-operand instructions.

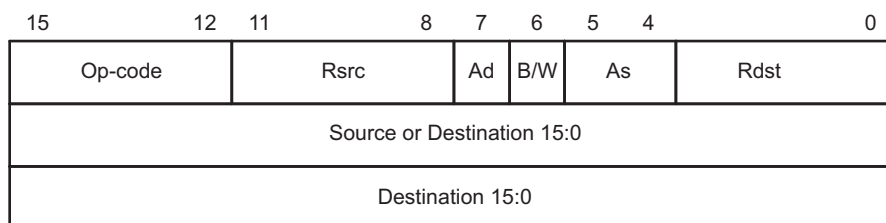


Figure 1-22. MSP430 Double-Operand Instruction Format

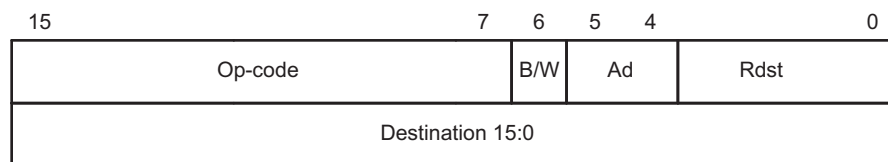
Table 1-4. MSP430 Double-Operand Instructions

Mnemonic	S-Reg, D-Reg	Operation	Status Bits ⁽¹⁾			
			V	N	Z	C
MOV (.B)	src,dst	src → dst	–	–	–	–
ADD (.B)	src,dst	src + dst → dst	*	*	*	*
ADDC (.B)	src,dst	src + dst + C → dst	*	*	*	*
SUB (.B)	src,dst	dst + .not.src + 1 → dst	*	*	*	*
SUBC (.B)	src,dst	dst + .not.src + C → dst	*	*	*	*
CMP (.B)	src,dst	dst - src	*	*	*	*
DADD (.B)	src,dst	src + dst + C → dst (decimally)	*	*	*	*
BIT (.B)	src,dst	src .and. dst	0	*	*	Z
BIC (.B)	src,dst	.not.src .and. dst → dst	–	–	–	–
BIS (.B)	src,dst	src .or. dst → dst	–	–	–	–
XOR (.B)	src,dst	src .xor. dst → dst	*	*	*	Z
AND (.B)	src,dst	src .and. dst → dst	0	*	*	Z

⁽¹⁾ * = Status bit is affected.
 – = Status bit is not affected.
 0 = Status bit is cleared.
 1 = Status bit is set.

1.5.1.2 MSP430 Single-Operand (Format II) Instructions

Figure 1-23 shows the format for MSP430 single-operand instructions, except RETI. The destination word is appended for the Indexed, Symbolic, Absolute, and Immediate modes. Table 1-5 lists the seven single-operand instructions.


Figure 1-23. MSP430 Single-Operand Instructions
Table 1-5. MSP430 Single-Operand Instructions

Mnemonic	S-Reg, D-Reg	Operation	Status Bits ⁽¹⁾			
			V	N	Z	C
RRC (.B)	dst	C → MSB →.....LSB → C	0	*	*	*
RRA (.B)	dst	MSB → MSB →....LSB → C	0	*	*	*
PUSH (.B)	src	SP - 2 → SP, src → SP	–	–	–	–
SWPB	dst	bit 15...bit 8 ↔ bit 7...bit 0	–	–	–	–
CALL	dst	Call subroutine in lower 64KB	–	–	–	–
RETI		TOS → SR, SP + 2 → SP TOS → PC, SP + 2 → SP	*	*	*	*
SXT	dst	Register mode: bit 7 → bit 8...bit 19 Other modes: bit 7 → bit 8...bit 15	0	*	*	Z

⁽¹⁾ * = Status bit is affected.
 – = Status bit is not affected.
 0 = Status bit is cleared.
 1 = Status bit is set.

1.5.1.3 Jump Instructions

Figure 1-24 shows the format for MSP430 and MSP430X jump instructions. The signed 10-bit word offset of the jump instruction is multiplied by two, sign-extended to a 20-bit address, and added to the 20-bit PC. This allows jumps in a range of -512 to $+511$ words relative to the PC in the full 20-bit address space. Jumps do not affect the status bits. Table 1-6 lists and describes the eight jump instructions.

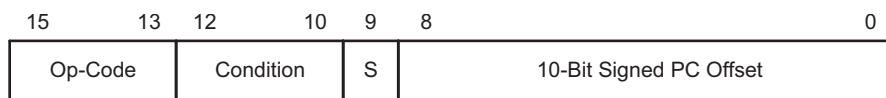


Figure 1-24. Format of Conditional Jump Instructions

Table 1-6. Conditional Jump Instructions

Mnemonic	S-Reg, D-Reg	Operation
JEQ, JZ	Label	Jump to label if zero bit is set
JNE, JNZ	Label	Jump to label if zero bit is reset
JC	Label	Jump to label if carry bit is set
JNC	Label	Jump to label if carry bit is reset
JN	Label	Jump to label if negative bit is set
JGE	Label	Jump to label if (N .XOR. V) = 0
JL	Label	Jump to label if (N .XOR. V) = 1
JMP	Label	Jump to label unconditionally

1.5.1.4 Emulated Instructions

In addition to the MSP430 and MSP430X instructions, emulated instructions are instructions that make code easier to write and read, but do not have op-codes themselves. Instead, they are replaced automatically by the assembler with a core instruction. There is no code or performance penalty for using emulated instructions. The emulated instructions are listed in Table 1-7.

Table 1-7. Emulated Instructions

Instruction	Explanation	Emulation	Status Bits ⁽¹⁾			
			V	N	Z	C
ADC(.B) dst	Add Carry to dst	ADDC(.B) #0, dst	*	*	*	*
BR dst	Branch indirectly dst	MOV dst, PC	–	–	–	–
CLR(.B) dst	Clear dst	MOV(.B) #0, dst	–	–	–	–
CLRC	Clear Carry bit	BIC #1, SR	–	–	–	0
CLRN	Clear Negative bit	BIC #4, SR	–	0	–	–
CLRZ	Clear Zero bit	BIC #2, SR	–	–	0	–
DADC(.B) dst	Add Carry to dst decimally	DADD(.B) #0, dst	*	*	*	*
DEC(.B) dst	Decrement dst by 1	SUB(.B) #1, dst	*	*	*	*
DECD(.B) dst	Decrement dst by 2	SUB(.B) #2, dst	*	*	*	*
DINT	Disable interrupt	BIC #8, SR	–	–	–	–
EINT	Enable interrupt	BIS #8, SR	–	–	–	–
INC(.B) dst	Increment dst by 1	ADD(.B) #1, dst	*	*	*	*
INCD(.B) dst	Increment dst by 2	ADD(.B) #2, dst	*	*	*	*

⁽¹⁾ * = Status bit is affected.
 – = Status bit is not affected.
 0 = Status bit is cleared.
 1 = Status bit is set.

Table 1-7. Emulated Instructions (continued)

Instruction	Explanation	Emulation	Status Bits ⁽¹⁾			
			V	N	Z	C
INV(.B) dst	Invert dst	XOR(.B) #-1, dst	*	*	*	*
NOP	No operation	MOV R3, R3	–	–	–	–
POP dst	Pop operand from stack	MOV @SP+, dst	–	–	–	–
RET	Return from subroutine	MOV @SP+, PC	–	–	–	–
RLA(.B) dst	Shift left dst arithmetically	ADD(.B) dst, dst	*	*	*	*
RLC(.B) dst	Shift left dst logically through Carry	ADDC(.B) dst, dst	*	*	*	*
SBC(.B) dst	Subtract Carry from dst	SUBC(.B) #0, dst	*	*	*	*
SETC	Set Carry bit	BIS #1, SR	–	–	–	1
SETN	Set Negative bit	BIS #4, SR	–	1	–	–
SETZ	Set Zero bit	BIS #2, SR	–	–	1	–
TST(.B) dst	Test dst (compare with 0)	CMP(.B) #0, dst	0	*	*	1

1.5.1.5 MSP430 Instruction Execution

The number of CPU clock cycles required for an instruction depends on the instruction format and the addressing modes used – not the instruction itself. The number of clock cycles refers to MCLK.

1.5.1.5.1 Instruction Cycles and Length for Interrupt, Reset, and Subroutines

Table 1-8 lists the length and the CPU cycles for reset, interrupts, and subroutines.

Table 1-8. Interrupt, Return, and Reset Cycles and Length

Action	Execution Time (MCLK Cycles)	Length of Instruction (Words)
Return from interrupt RETI	5	1
Return from subroutine RET	4	1
Interrupt request service (cycles needed before first instruction)	6	–
WDT reset	4	–
Reset (RST/NMI)	4	–

1.5.1.5.2 Format II (Single-Operand) Instruction Cycles and Lengths

Table 1-9 lists the length and the CPU cycles for all addressing modes of the MSP430 single-operand instructions.

Table 1-9. MSP430 Format II Instruction Cycles and Length

Addressing Mode	No. of Cycles			Length of Instruction	Example
	RRA, RRC SWPB, SXT	PUSH	CALL		
Rn	1	3	4	1	SWPB R5
@Rn	3	3	4	1	RRC @R9
@Rn+	3	3	4	1	SWPB @R10+
#N	N/A	3	4	2	CALL #LABEL
X(Rn)	4	4	5	2	CALL 2(R7)
EDE	4	4	5	2	PUSH EDE
&EDE	4	4	6	2	SXT &EDE

1.5.1.5.3 Jump Instructions Cycles and Lengths

All jump instructions require one code word and take two CPU cycles to execute, regardless of whether the jump is taken or not.

1.5.1.5.4 Format I (Double-Operand) Instruction Cycles and Lengths

Table 1-10 lists the length and CPU cycles for all addressing modes of the MSP430 Format I instructions.

Table 1-10. MSP430 Format I Instructions Cycles and Length

Addressing Mode		No. of Cycles	Length of Instruction	Example
Source	Destination			
Rn	Rm	1	1	MOV R5, R8
	PC	3	1	BR R9
	x(Rm)	4 ⁽¹⁾	2	ADD R5, 4(R6)
	EDE	4 ⁽¹⁾	2	XOR R8, EDE
	&EDE	4 ⁽¹⁾	2	MOV R5, &EDE
@Rn	Rm	2	1	AND @R4, R5
	PC	4	1	BR @R8
	x(Rm)	5 ⁽¹⁾	2	XOR @R5, 8(R6)
	EDE	5 ⁽¹⁾	2	MOV @R5, EDE
	&EDE	5 ⁽¹⁾	2	XOR @R5, &EDE
@Rn+	Rm	2	1	ADD @R5+, R6
	PC	4	1	BR @R9+
	x(Rm)	5 ⁽¹⁾	2	XOR @R5, 8(R6)
	EDE	5 ⁽¹⁾	2	MOV @R9+, EDE
	&EDE	5 ⁽¹⁾	2	MOV @R9+, &EDE
#N	Rm	2	2	MOV #20, R9
	PC	3	2	BR #2AEh
	x(Rm)	5 ⁽¹⁾	3	MOV #0300h, 0(SP)
	EDE	5 ⁽¹⁾	3	ADD #33, EDE
	&EDE	5 ⁽¹⁾	3	ADD #33, &EDE
x(Rn)	Rm	3	2	MOV 2(R5), R7
	PC	5	2	BR 2(R6)
	TONI	6 ⁽¹⁾	3	MOV 4(R7), TONI
	x(Rm)	6 ⁽¹⁾	3	ADD 4(R4), 6(R9)
	&TONI	6 ⁽¹⁾	3	MOV 2(R4), &TONI
EDE	Rm	3	2	AND EDE, R6
	PC	5	2	BR EDE
	TONI	6 ⁽¹⁾	3	CMP EDE, TONI
	x(Rm)	6 ⁽¹⁾	3	MOV EDE, 0(SP)
	&TONI	6 ⁽¹⁾	3	MOV EDE, &TONI
&EDE	Rm	3	2	MOV &EDE, R8
	PC	5	2	BR &EDE
	TONI	6 ⁽¹⁾	3	MOV &EDE, TONI
	x(Rm)	6 ⁽¹⁾	3	MOV &EDE, 0(SP)
	&TONI	6 ⁽¹⁾	3	MOV &EDE, &TONI

⁽¹⁾ MOV, BIT, and CMP instructions execute in one fewer cycle.

1.5.2 MSP430X Extended Instructions

The extended MSP430X instructions give the MSP430X CPU full access to its 20-bit address space. Most MSP430X instructions require an additional word of op-code called the extension word. Some extended instructions do not require an additional word and are noted in the instruction description. All addresses, indexes, and immediate numbers have 20-bit values when preceded by the extension word.

There are two types of extension words:

- Register or register mode for Format I instructions and register mode for Format II instructions
- Extension word for all other address mode combinations

1.5.2.1 Register Mode Extension Word

The register mode extension word is shown in [Figure 1-25](#) and described in [Table 1-11](#). An example is shown in [Figure 1-27](#).

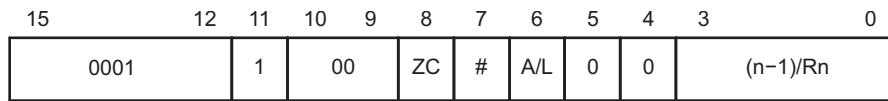


Figure 1-25. Extension Word for Register Modes

Table 1-11. Description of the Extension Word Bits for Register Mode

Bit	Description															
15:11	Extension word op-code. Op-codes 1800h to 1FFFh are extension words.															
10:9	Reserved															
ZC	Zero carry 0 The executed instruction uses the status of the carry bit C. 1 The executed instruction uses the carry bit as 0. The carry bit is defined by the result of the final operation after instruction execution.															
#	Repetition 0 The number of instruction repetitions is set by extension word bits 3:0. 1 The number of instruction repetitions is defined by the value of the four LSBs of Rn. See description for bits 3:0.															
A/L	Data length extension. Together with the B/W bits of the following MSP430 instruction, the AL bit defines the used data length of the instruction. <table border="1" style="width: 100%; margin-top: 5px;"> <thead> <tr> <th>A/L</th> <th>B/W</th> <th>Comment</th> </tr> </thead> <tbody> <tr> <td>0</td> <td>0</td> <td>Reserved</td> </tr> <tr> <td>0</td> <td>1</td> <td>20-bit address word</td> </tr> <tr> <td>1</td> <td>0</td> <td>16-bit word</td> </tr> <tr> <td>1</td> <td>1</td> <td>8-bit byte</td> </tr> </tbody> </table>	A/L	B/W	Comment	0	0	Reserved	0	1	20-bit address word	1	0	16-bit word	1	1	8-bit byte
A/L	B/W	Comment														
0	0	Reserved														
0	1	20-bit address word														
1	0	16-bit word														
1	1	8-bit byte														
5:4	Reserved															
3:0	Repetition count # = 0 These four bits set the repetition count n. These bits contain n – 1. # = 1 These four bits define the CPU register whose bits 3:0 set the number of repetitions. Rn.3:0 contain n – 1.															

1.5.2.2 Non-Register Mode Extension Word

The extension word for non-register modes is shown in [Figure 1-26](#) and described in [Table 1-12](#). An example is shown in [Figure 1-28](#).

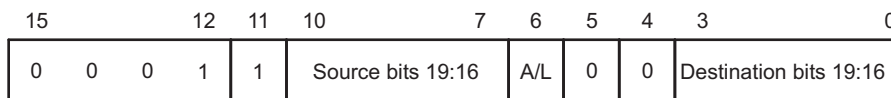


Figure 1-26. Extension Word for Non-Register Modes

Table 1-12. Description of Extension Word Bits for Non-Register Modes

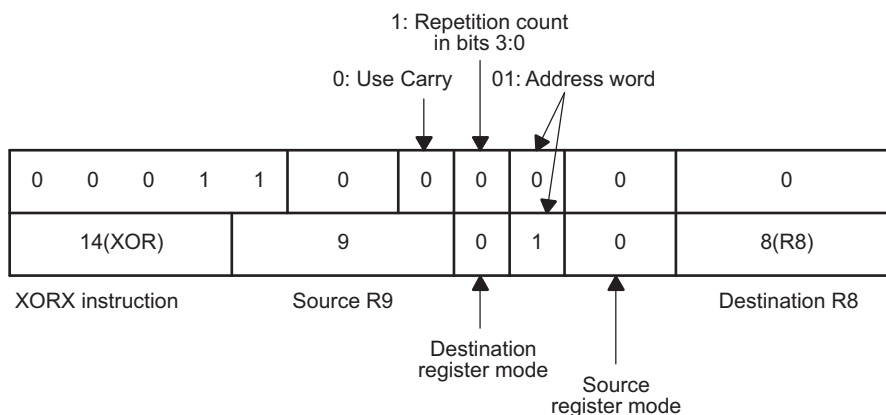
Bit	Description															
15:11	Extension word op-code. Op-codes 1800h to 1FFFh are extension words.															
Source Bits 19:16	The four MSBs of the 20-bit source. Depending on the source addressing mode, these four MSBs may belong to an immediate operand, an index, or to an absolute address.															
A/L	Data length extension. Together with the B/W bits of the following MSP430 instruction, the AL bit defines the used data length of the instruction. <table border="1"> <thead> <tr> <th>A/L</th> <th>B/W</th> <th>Comment</th> </tr> </thead> <tbody> <tr> <td>0</td> <td>0</td> <td>Reserved</td> </tr> <tr> <td>0</td> <td>1</td> <td>20-bit address word</td> </tr> <tr> <td>1</td> <td>0</td> <td>16-bit word</td> </tr> <tr> <td>1</td> <td>1</td> <td>8-bit byte</td> </tr> </tbody> </table>	A/L	B/W	Comment	0	0	Reserved	0	1	20-bit address word	1	0	16-bit word	1	1	8-bit byte
A/L	B/W	Comment														
0	0	Reserved														
0	1	20-bit address word														
1	0	16-bit word														
1	1	8-bit byte														
5:4	Reserved															
Destination Bits 19:16	The four MSBs of the 20-bit destination. Depending on the destination addressing mode, these four MSBs may belong to an index or to an absolute address.															

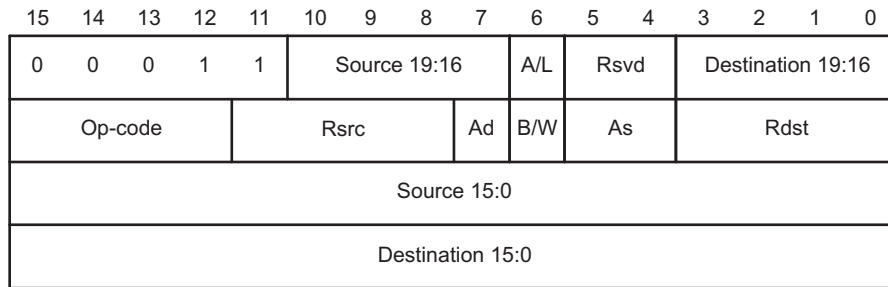
NOTE: B/W and A/L bit settings for SWPBX and SXTX

A/L	B/W	
0	0	SWPBX.A, SXTX.A
0	1	N/A
1	0	SWPB.W, SXTX.W
1	1	N/A

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
0	0	0	1	1	00	ZC	#	A/L	Rsvd	(n-1)/Rn					
Op-code				Rsrc				Ad	B/W	As	Rdst				

XORX.A R9, R8

**Figure 1-27. Example for Extended Register or Register Instruction**



XORX.A #12345h, 45678h(R15)

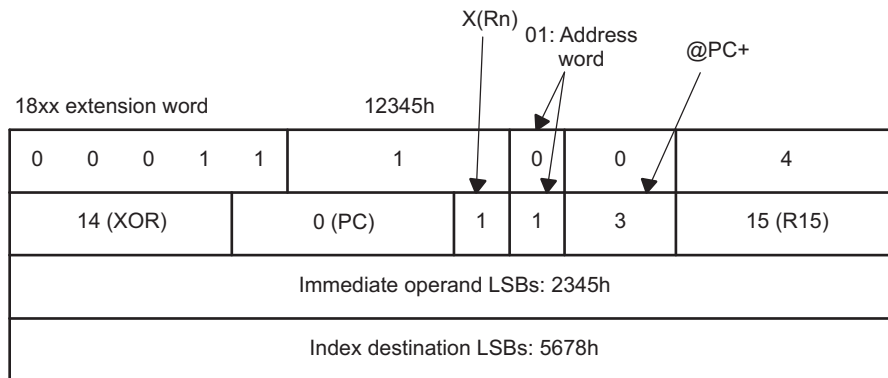


Figure 1-28. Example for Extended Immediate or Indexed Instruction

1.5.2.3 Extended Double-Operand (Format I) Instructions

All 12 double-operand instructions have extended versions as listed in [Table 1-13](#).

Table 1-13. Extended Double-Operand Instructions

Mnemonic	Operands	Operation	Status Bits ⁽¹⁾			
			V	N	Z	C
MOVX(.B, .A)	src,dst	src → dst	–	–	–	–
ADDX(.B, .A)	src,dst	src + dst → dst	*	*	*	*
ADDCX(.B, .A)	src,dst	src + dst + C → dst	*	*	*	*
SUBX(.B, .A)	src,dst	dst + .not.src + 1 → dst	*	*	*	*
SUBCX(.B, .A)	src,dst	dst + .not.src + C → dst	*	*	*	*
CMPX(.B, .A)	src,dst	dst – src	*	*	*	*
DADDX(.B, .A)	src,dst	src + dst + C → dst (decimal)	*	*	*	*
BITX(.B, .A)	src,dst	src .and. dst	0	*	*	Z
BICX(.B, .A)	src,dst	.not.src .and. dst → dst	–	–	–	–
BISX(.B, .A)	src,dst	src .or. dst → dst	–	–	–	–
XORX(.B, .A)	src,dst	src .xor. dst → dst	*	*	*	Z
ANDX(.B, .A)	src,dst	src .and. dst → dst	0	*	*	Z

⁽¹⁾ * = Status bit is affected.
 – = Status bit is not affected.
 0 = Status bit is cleared.
 1 = Status bit is set.

The four possible addressing combinations for the extension word for Format I instructions are shown in Figure 1-29.

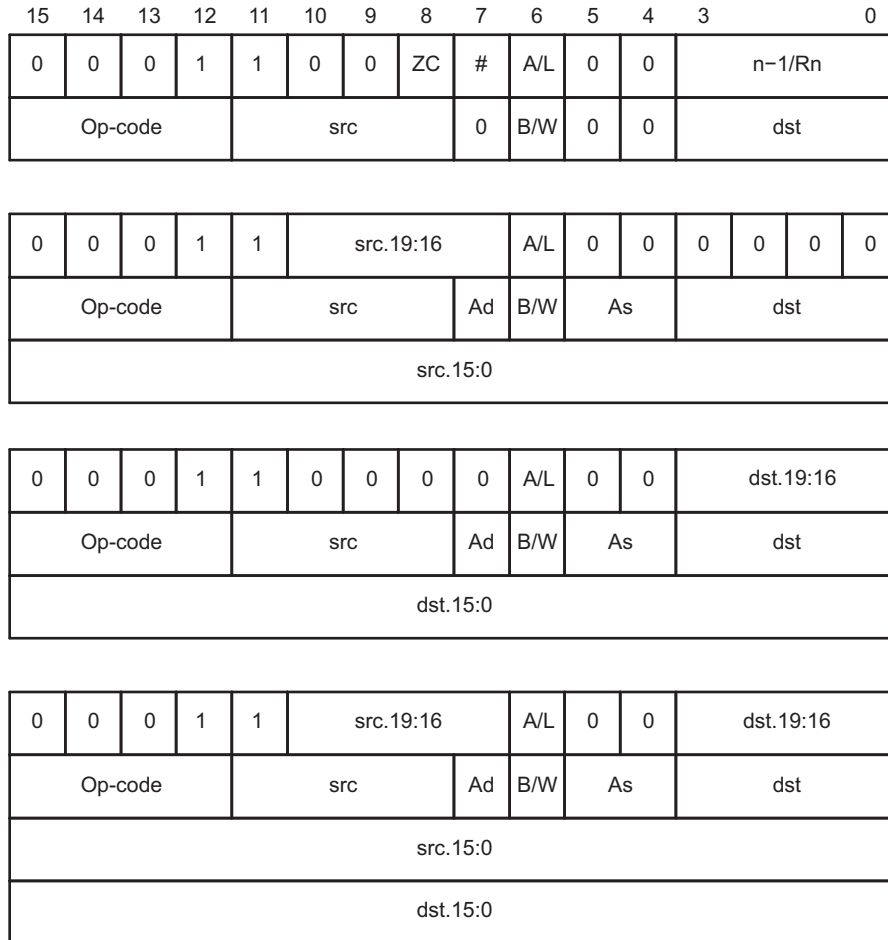


Figure 1-29. Extended Format I Instruction Formats

If the 20-bit address of a source or destination operand is located in memory, not in a CPU register, then two words are used for this operand as shown in Figure 1-30.

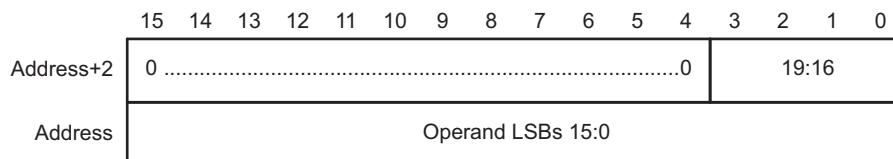


Figure 1-30. 20-Bit Addresses in Memory

1.5.2.4 Extended Single-Operand (Format II) Instructions

Extended MSP430X Format II instructions are listed in Table 1-14.

Table 1-14. Extended Single-Operand Instructions

Mnemonic	Operands	Operation	n	Status Bits ⁽¹⁾			
				V	N	Z	C
CALLA	dst	Call indirect to subroutine (20-bit address)		-	-	-	-
POPM.A	#n,Rdst	Pop n 20-bit registers from stack	1 to 16	-	-	-	-
POPM.W	#n,Rdst	Pop n 16-bit registers from stack	1 to 16	-	-	-	-
PUSHM.A	#n,Rsrc	Push n 20-bit registers to stack	1 to 16	-	-	-	-
PUSHM.W	#n,Rsrc	Push n 16-bit registers to stack	1 to 16	-	-	-	-
PUSHX(.B, .A)	src	Push 8-, 16-, or 20-bit source to stack		-	-	-	-
RRCM(.A)	#n,Rdst	Rotate right Rdst n bits through carry (16-, 20-bit register)	1 to 4	0	*	*	*
RRUM(.A)	#n,Rdst	Rotate right Rdst n bits unsigned (16-, 20-bit register)	1 to 4	0	*	*	*
RRAM(.A)	#n,Rdst	Rotate right Rdst n bits arithmetically (16-, 20-bit register)	1 to 4	0	*	*	*
RLAM(.A)	#n,Rdst	Rotate left Rdst n bits arithmetically (16-, 20-bit register)	1 to 4	*	*	*	*
RRCX(.B, .A)	dst	Rotate right dst through carry (8-, 16-, 20-bit data)	1	0	*	*	*
RRUX(.B, .A)	Rdst	Rotate right dst unsigned (8-, 16-, 20-bit)	1	0	*	*	*
RRAX(.B, .A)	dst	Rotate right dst arithmetically	1	0	*	*	*
SWPBX(.A)	dst	Exchange low byte with high byte	1	-	-	-	-
SXTX(.A)	Rdst	Bit7 → bit8 ... bit19	1	0	*	*	Z
SXTX(.A)	dst	Bit7 → bit8 ... MSB	1	0	*	*	Z

⁽¹⁾ * = Status bit is affected.
 - = Status bit is not affected.
 0 = Status bit is cleared.
 1 = Status bit is set.

The three possible addressing mode combinations for Format II instructions are shown in Figure 1-31.

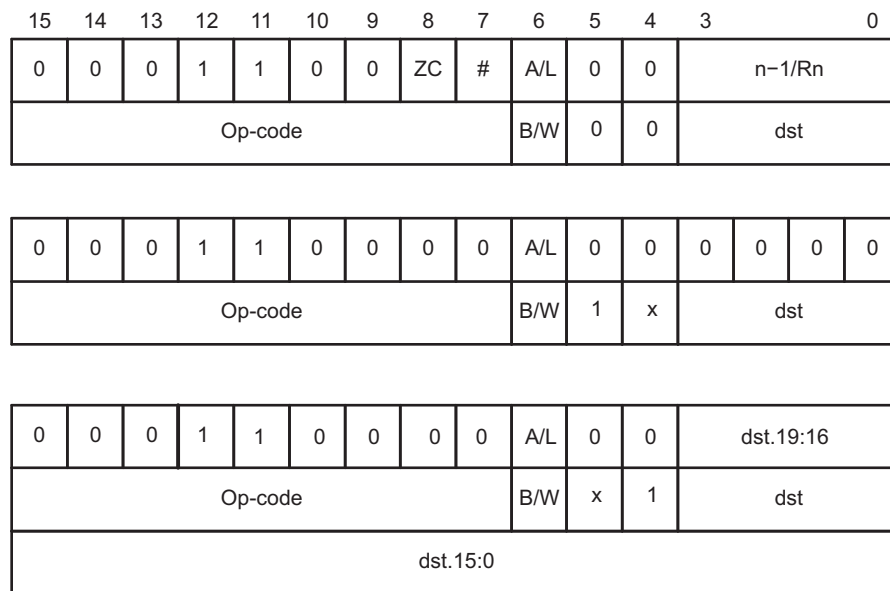


Figure 1-31. Extended Format II Instruction Format

1.5.2.4.1 Extended Format II Instruction Format Exceptions

Exceptions for the Format II instruction formats are shown in [Figure 1-32](#) through [Figure 1-35](#).

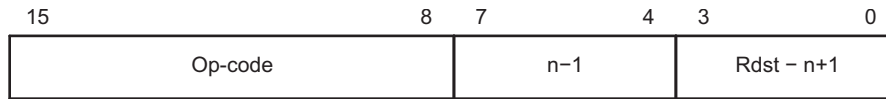


Figure 1-32. PUSHM and POPM Instruction Format



Figure 1-33. RRCM, RRAM, RRUM, and RLAM Instruction Format

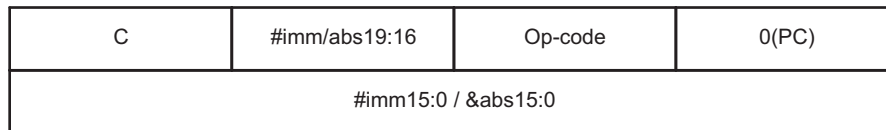
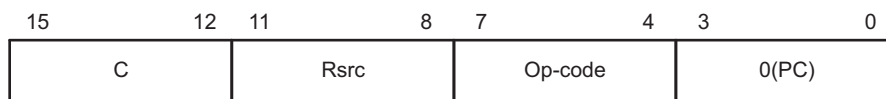


Figure 1-34. BRA Instruction Format

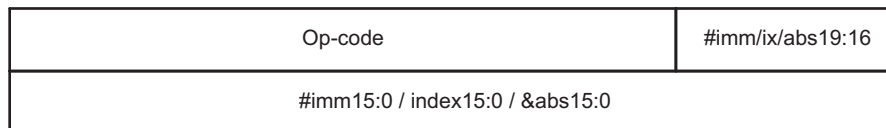
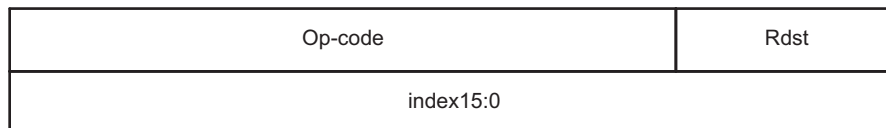
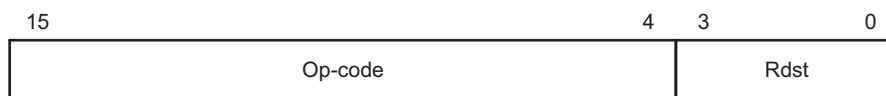


Figure 1-35. CALLA Instruction Format

1.5.2.5 Extended Emulated Instructions

The extended instructions together with the constant generator form the extended emulated instructions. [Table 1-15](#) lists the emulated instructions.

Table 1-15. Extended Emulated Instructions

Instruction	Explanation	Emulation
ADCX(.B, .A) dst	Add carry to dst	ADDCX(.B, .A) #0, dst
BRA dst	Branch indirect dst	MOVA dst, PC
RETA	Return from subroutine	MOVA @SP+, PC
CLRA Rdst	Clear Rdst	MOV #0, Rdst
CLR(.B, .A) dst	Clear dst	MOVX(.B, .A) #0, dst
DADCX(.B, .A) dst	Add carry to dst decimally	DADDX(.B, .A) #0, dst
DECX(.B, .A) dst	Decrement dst by 1	SUBX(.B, .A) #1, dst
DECDA Rdst	Decrement Rdst by 2	SUBA #2, Rdst
DECDX(.B, .A) dst	Decrement dst by 2	SUBX(.B, .A) #2, dst
INCX(.B, .A) dst	Increment dst by 1	ADDX(.B, .A) #1, dst
INCDA Rdst	Increment Rdst by 2	ADDA #2, Rdst
INCDX(.B, .A) dst	Increment dst by 2	ADDX(.B, .A) #2, dst
INVX(.B, .A) dst	Invert dst	XORX(.B, .A) #-1, dst
RLAX(.B, .A) dst	Shift left dst arithmetically	ADDX(.B, .A) dst, dst
RLCX(.B, .A) dst	Shift left dst logically through carry	ADDCX(.B, .A) dst, dst
SBCX(.B, .A) dst	Subtract carry from dst	SUBCX(.B, .A) #0, dst
TSTA Rdst	Test Rdst (compare with 0)	CMPA #0, Rdst
TSTX(.B, .A) dst	Test dst (compare with 0)	CMPX(.B, .A) #0, dst
POPX dst	Pop to dst	MOVX(.B, .A) @SP+, dst

1.5.2.6 MSP430X Address Instructions

MSP430X address instructions are instructions that support 20-bit operands but have restricted addressing modes. The addressing modes are restricted to the Register mode and the Immediate mode, except for the MOVA instruction as listed in [Table 1-16](#). Restricting the addressing modes removes the need for the additional extension-word op-code improving code density and execution time. Address instructions should be used any time an MSP430X instruction is needed with the corresponding restricted addressing mode.

Table 1-16. Address Instructions, Operate on 20-Bit Register Data

Mnemonic	Operands	Operation	Status Bits ⁽¹⁾			
			V	N	Z	C
ADDA	Rsrc, Rdst	Add source to destination register	*	*	*	*
	#imm20, Rdst					
MOVA	Rsrc, Rdst	Move source to destination	-	-	-	-
	#imm20, Rdst					
	z16(Rsrc), Rdst					
	EDE, Rdst					
	&abs20, Rdst					
	@Rsrc, Rdst					
	@Rsrc+, Rdst					
	Rsrc, z16(Rdst)					
	Rsrc, &abs20					
	Rsrc, Rdst		Compare source to destination register	*	*	*
#imm20, Rdst						
SUBA	Rsrc, Rdst	Subtract source from destination register	*	*	*	*
	#imm20, Rdst					

⁽¹⁾ * = Status bit is affected.
 - = Status bit is not affected.
 0 = Status bit is cleared.
 1 = Status bit is set.

1.5.2.7 MSP430X Instruction Execution

The number of CPU clock cycles required for an MSP430X instruction depends on the instruction format and the addressing modes used, not the instruction itself. The number of clock cycles refers to MCLK.

1.5.2.7.1 MSP430X Format II (Single-Operand) Instruction Cycles and Lengths

Table 1-17 lists the length and the CPU cycles for all addressing modes of the MSP430X extended single-operand instructions.

Table 1-17. MSP430X Format II Instruction Cycles and Length

Instruction	Execution Cycles, Length of Instruction (Words)						
	Rn	@Rn	@Rn+	#N	X(Rn)	EDE	&EDE
RRAM	n, 1	–	–	–	–	–	–
RRCM	n, 1	–	–	–	–	–	–
RRUM	n, 1	–	–	–	–	–	–
RLAM	n, 1	–	–	–	–	–	–
PUSHM	2+n, 1	–	–	–	–	–	–
PUSHM.A	2+2n, 1	–	–	–	–	–	–
POPM	2+n, 1	–	–	–	–	–	–
POPM.A	2+2n, 1	–	–	–	–	–	–
CALLA	5, 1	6, 1	6, 1	5, 2	5 ⁽¹⁾ , 2	7, 2	7, 2
RRAX(.B)	1+n, 2	4, 2	4, 2	–	5, 3	5, 3	5, 3
RRAX.A	1+n, 2	6, 2	6, 2	–	7, 3	7, 3	7, 3
RRCX(.B)	1+n, 2	4, 2	4, 2	–	5, 3	5, 3	5, 3
RRCX.A	1+n, 2	6, 2	6, 2	–	7, 3	7, 3	7, 3
PUSHX(.B)	4, 2	4, 2	4, 2	4, 3	5 ⁽¹⁾ , 3	5, 3	5, 3
PUSHX.A	5, 2	6, 2	6, 2	5, 3	7 ⁽¹⁾ , 3	7, 3	7, 3
POPX(.B)	3, 2	–	–	–	5, 3	5, 3	5, 3
POPX.A	4, 2	–	–	–	7, 3	7, 3	7, 3

⁽¹⁾ Add one cycle when Rn = SP

1.5.2.7.2 MSP430X Format I (Double-Operand) Instruction Cycles and Lengths

Table 1-18 lists the length and CPU cycles for all addressing modes of the MSP430X extended Format I instructions.

Table 1-18. MSP430X Format I Instruction Cycles and Length

Addressing Mode		No. of Cycles		Length of Instruction	Examples
Source	Destination	.B/.W	.A	.B/.W/.A	
Rn	Rm ⁽¹⁾	2	2	2	BITX.B R5,R8
	PC	4	4	2	ADDX.R9,PC
	x(Rm)	5 ⁽²⁾	7 ⁽³⁾	3	ANDX.A R5,4(R6)
	EDE	5 ⁽²⁾	7 ⁽³⁾	3	XORX.R8,EDE
	&EDE	5 ⁽²⁾	7 ⁽³⁾	3	BITX.W R5,&EDE
@Rn	Rm	3	4	2	BITX.@R5,R8
	PC	5	6	2	ADDX.@R9,PC
	x(Rm)	6 ⁽²⁾	9 ⁽³⁾	3	ANDX.A @R5,4(R6)
	EDE	6 ⁽²⁾	9 ⁽³⁾	3	XORX.@R8,EDE
	&EDE	6 ⁽²⁾	9 ⁽³⁾	3	BITX.B @R5,&EDE
@Rn+	Rm	3	4	2	BITX.@R5+,R8
	PC	5	6	2	ADDX.A @R9+,PC
	x(Rm)	6 ⁽²⁾	9 ⁽³⁾	3	ANDX.@R5+,4(R6)
	EDE	6 ⁽²⁾	9 ⁽³⁾	3	XORX.B @R8+,EDE
	&EDE	6 ⁽²⁾	9 ⁽³⁾	3	BITX.@R5+,&EDE
#N	Rm	3	3	3	BITX.#20,R8
	PC ⁽⁴⁾	4	4	3	ADDX.A #FE00h,PC
	x(Rm)	6 ⁽²⁾	8 ⁽³⁾	4	ANDX.#1234,4(R6)
	EDE	6 ⁽²⁾	8 ⁽³⁾	4	XORX.#A5A5h,EDE
	&EDE	6 ⁽²⁾	8 ⁽³⁾	4	BITX.B #12,&EDE
x(Rn)	Rm	4	5	3	BITX.2(R5),R8
	PC ⁽⁴⁾	6	7	3	SUBX.A 2(R6),PC
	TONI	7 ⁽²⁾	10 ⁽³⁾	4	ANDX.4(R7),4(R6)
	x(Rm)	7 ⁽²⁾	10 ⁽³⁾	4	XORX.B 2(R6),EDE
	&TONI	7 ⁽²⁾	10 ⁽³⁾	4	BITX.8(SP),&EDE
EDE	Rm	4	5	3	BITX.B EDE,R8
	PC ⁽⁴⁾	6	7	3	ADDX.A EDE,PC
	TONI	7 ⁽²⁾	10 ⁽³⁾	4	ANDX.EDE,4(R6)
	x(Rm)	7 ⁽²⁾	10 ⁽³⁾	4	ANDX.EDE,TONI
	&TONI	7 ⁽²⁾	10 ⁽³⁾	4	BITX.EDE,&TONI
&EDE	Rm	4	5	3	BITX.&EDE,R8
	PC ⁽⁴⁾	6	7	3	ADDX.A &EDE,PC
	TONI	7 ⁽²⁾	10 ⁽³⁾	4	ANDX.B &EDE,4(R6)
	x(Rm)	7 ⁽²⁾	10 ⁽³⁾	4	XORX.&EDE,TONI
	&TONI	7 ⁽²⁾	10 ⁽³⁾	4	BITX.&EDE,&TONI

⁽¹⁾ Repeat instructions require n + 1 cycles, where n is the number of times the instruction is executed.

⁽²⁾ Reduce the cycle count by one for MOV, BIT, and CMP instructions.

⁽³⁾ Reduce the cycle count by two for MOV, BIT, and CMP instructions.

⁽⁴⁾ Reduce the cycle count by one for MOV, ADD, and SUB instructions.

1.5.2.7.3 MSP430X Address Instruction Cycles and Lengths

Table 1-19 lists the length and the CPU cycles for all addressing modes of the MSP430X address instructions.

Table 1-19. Address Instruction Cycles and Length

Addressing Mode		Execution Time (MCLK Cycles)		Length of Instruction (Words)		Example
Source	Destination	MOVA BRA	CMPA ADDA SUBA	MOVA	CMPA ADDA SUBA	
Rn	Rn	1	1	1	1	CMPA R5, R8
	PC	3	3	1	1	SUBA R9, PC
	x(Rn)	4	–	2	–	MOVA R5, 4(R6)
	EDE	4	–	2	–	MOVA R8, EDE
	&EDE	4	–	2	–	MOVA R5, &EDE
@Rn	Rm	3	–	1	–	MOVA @R5, R8
	PC	5	–	1	–	MOVA @R9, PC
@Rn+	Rm	3	–	1	–	MOVA @R5+, R8
	PC	5	–	1	–	MOVA @R9+, PC
#N	Rm	2	3	2	2	CMPA #20, R8
	PC	3	3	2	2	SUBA #FE00h, PC
x(Rn)	Rm	4	–	2	–	MOVA 2(R5), R8
	PC	6	–	2	–	MOVA 2(R6), PC
EDE	Rm	4	–	2	–	MOVA EDE, R8
	PC	6	–	2	–	MOVA EDE, PC
&EDE	Rm	4	–	2	–	MOVA &EDE, R8
	PC	6	–	2	–	MOVA &EDE, PC

1.6 Instruction Set Description

Table 1-20 shows all available instructions:

Table 1-20. Instruction Map of MSP430X

	000	040	080	0C0	100	140	180	1C0	200	240	280	2C0	300	340	380	3C0
0xxx	MOVA, CMPA, ADDA, SUBA, RRCM, RRAM, RLAM, RRUM															
10xx	RRC	RRC. B	SWP B		RRA	RRA. B	SXT		PUS H	PUS H.B	CALL		RETI	CALL A		
14xx	PUSHM.A, POPM.A, PUSHM.W, POPM.W															
18xx	Extension word for Format I and Format II instructions															
1Cxx	Extension word for Format I and Format II instructions															
20xx	JNE, JNZ															
24xx	JEQ, JZ															
28xx	JNC															
2Cxx	JC															
30xx	JN															
34xx	JGE															
38xx	JL															
3Cxx	JMP															
4xxx	MOV, MOV.B															
5xxx	ADD, ADD.B															
6xxx	ADDC, ADDC.B															
7xxx	SUBC, SUBC.B															
8xxx	SUB, SUB.B															
9xxx	CMP, CMP.B															
Axxx	DADD, DADD.B															
Bxxx	BIT, BIT.B															
Cxxx	BIC, BIC.B															
Dxxx	BIS, BIS.B															
Exxx	XOR, XOR.B															
Fxxx	AND, AND.B															

1.6.1 Extended Instruction Binary Descriptions

Detailed MSP430X instruction binary descriptions are shown in the following tables.

Instruction	Instruction Group				src or data.19:16		Instruction Identifier				dst		
	15	12	11	8	7	4	3	0					
MOVA	0	0	0	0	src		0	0	0	0	dst		MOVA @Rsrc,Rdst
	0	0	0	0	src		0	0	0	1	dst		MOVA @Rsrc+,Rdst
	0	0	0	0	&abs.19:16		0	0	1	0	dst		MOVA &abs20,Rdst
	&abs.15:0												
	0	0	0	0	src		0	0	1	1	dst		MOVA z16(Rsrc),Rdst
	x.15:0												
	0	0	0	0	src		0	1	1	0	&abs.19:16		MOVA Rsrc,&abs20
	&abs.15:0												
	0	0	0	0	src		0	1	1	1	dst		MOVA Rsrc,z16(Rdst)
	x.15:0												
CMPA	0	0	0	0	imm.19:16		1	0	0	0	dst		MOVA #imm20,Rdst
	imm.15:0												
	0	0	0	0	imm.19:16		1	0	0	1	dst		CMPA #imm20,Rdst
ADDA	imm.15:0												
	0	0	0	0	imm.19:16		1	0	1	0	dst		ADDA #imm20,Rdst
SUBA	imm.15:0												
	0	0	0	0	imm.19:16		1	0	1	1	dst		SUBA #imm20,Rdst
MOVA	imm.15:0												
	0	0	0	0	src		1	1	0	0	dst		MOVA Rsrc,Rdst
	0	0	0	0	src		1	1	0	1	dst		CMPA Rsrc,Rdst
	0	0	0	0	src		1	1	1	0	dst		ADDA Rsrc,Rdst
	0	0	0	0	src		1	1	1	1	dst		SUBA Rsrc,Rdst

Instruction	Instruction Group				Bit Loc.		Inst. ID		Instruction Identifier				dst		
	15	12	11	10	9	8	7	4	3	0					
RRCM.A	0	0	0	0	n-1	0	0	0	1	0	0	dst		RRCM.A #n,Rdst	
RRAM.A	0	0	0	0	n-1	0	1	0	1	0	0	dst		RRAM.A #n,Rdst	
RLAM.A	0	0	0	0	n-1	1	0	0	1	0	0	dst		RLAM.A #n,Rdst	
RRUM.A	0	0	0	0	n-1	1	1	0	1	0	0	dst		RRUM.A #n,Rdst	
RRCM.W	0	0	0	0	n-1	0	0	0	1	0	1	dst		RRCM.W #n,Rdst	
RRAM.W	0	0	0	0	n-1	0	1	0	1	0	1	dst		RRAM.W #n,Rdst	
RLAM.W	0	0	0	0	n-1	1	0	0	1	0	1	dst		RLAM.W #n,Rdst	
RRUM.W	0	0	0	0	n-1	1	1	0	1	0	1	dst		RRUM.W #n,Rdst	

Instruction Set Description

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Instruction	Instruction Identifier												dst							
	15	12	11	8	7	6	5	4	3	0										
RETI	0	0	0	1	0	0	1	1	0	0	0	0	0	0	0	0				
CALLA	0	0	0	1	0	0	1	1	0	1	0	0	dst				CALLA Rdst			
	0	0	0	1	0	0	1	1	0	1	0	1	dst				CALLA x(Rdst)			
	x.15:0																			
	0	0	0	1	0	0	1	1	0	1	1	0	dst				CALLA @Rdst			
	0	0	0	1	0	0	1	1	0	1	1	1	dst				CALLA @Rdst+			
	0	0	0	1	0	0	1	1	1	0	0	0	&abs.19:16				CALLA &abs20			
	&abs.15:0																			
	0	0	0	1	0	0	1	1	1	0	0	1	x.19:16				CALLA EDE			
	x.15:0																			
	0	0	0	1	0	0	1	1	1	0	1	1	imm.19:16				CALLA #imm20			
	imm.15:0																			
Reserved	0	0	0	1	0	0	1	1	1	0	1	0	x	x	x	x				
Reserved	0	0	0	1	0	0	1	1	1	1	x	x	x	x	x	x				
PUSHM.A	0	0	0	1	0	1	0	0	n - 1				dst				PUSHM.A #n,Rdst			
PUSHM.W	0	0	0	1	0	1	0	1	n - 1				dst				PUSHM.W #n,Rdst			
POPM.A	0	0	0	1	0	1	1	0	n - 1				dst - n + 1				POPM.A #n,Rdst			
POPM.W	0	0	0	1	0	1	1	1	n - 1				dst - n + 1				POPM.W #n,Rdst			

1.6.2 MSP430 Instructions

The MSP430 instructions are listed and described on the following pages.

1.6.2.2 ADD

ADD[.W] Add source word to destination word

ADD.B Add source byte to destination byte

Syntax ADD src,dst OR ADD.W src,dst
ADD.B src,dst

Operation src + dst → dst

Description The source operand is added to the destination operand. The previous content of the destination is lost.

Status Bits N: Set if result is negative (MSB = 1), reset if positive (MSB = 0)
Z: Set if result is zero, reset otherwise
C: Set if there is a carry from the MSB of the result, reset otherwise
V: Set if the result of two positive operands is negative, or if the result of two negative numbers is positive, reset otherwise

Mode Bits OSCOFF, CPUOFF, and GIE are not affected.

Example Ten is added to the 16-bit counter CNTR located in lower 64 K.

```
ADD.W #10,&CNTR ; Add 10 to 16-bit counter
```

Example A table word pointed to by R5 (20-bit address in R5) is added to R6. The jump to label TONI is performed on a carry.

```
ADD.W @R5,R6 ; Add table word to R6. R6.19:16 = 0
JC TONI ; Jump if carry
... ; No carry
```

Example A table byte pointed to by R5 (20-bit address) is added to R6. The jump to label TONI is performed if no carry occurs. The table pointer is auto-incremented by 1. R6.19:8 = 0

```
ADD.B @R5+,R6 ; Add byte to R6. R5 + 1. R6: 000xxh
JNC TONI ; Jump if no carry
... ; Carry occurred
```

1.6.2.3 ADDC**ADDC[.W]** Add source word and carry to destination word**ADDC.B** Add source byte and carry to destination byte**Syntax** `ADDC src,dst` OR `ADDC.W src,dst`
`ADDC.B src,dst`**Operation** `src + dst + C → dst`**Description** The source operand and the carry bit C are added to the destination operand. The previous content of the destination is lost.**Status Bits** N: Set if result is negative (MSB = 1), reset if positive (MSB = 0)

Z: Set if result is zero, reset otherwise

C: Set if there is a carry from the MSB of the result, reset otherwise

V: Set if the result of two positive operands is negative, or if the result of two negative numbers is positive, reset otherwise

Mode Bits OSCOFF, CPUOFF, and GIE are not affected.**Example** Constant value 15 and the carry of the previous instruction are added to the 16-bit counter CNTR located in lower 64 K.

```
ADDC.W    #15,&CNTR    ; Add 15 + C to 16-bit CNTR
```

Example A table word pointed to by R5 (20-bit address) and the carry C are added to R6. The jump to label TONI is performed on a carry. R6.19:16 = 0

```
ADDC.W    @R5,R6      ; Add table word + C to R6
JC        TONI        ; Jump if carry
...       ; No carry
```

Example A table byte pointed to by R5 (20-bit address) and the carry bit C are added to R6. The jump to label TONI is performed if no carry occurs. The table pointer is auto-incremented by 1. R6.19:8 = 0

```
ADDC.B    @R5+,R6     ; Add table byte + C to R6. R5 + 1
JNC       TONI        ; Jump if no carry
...       ; Carry occurred
```

1.6.2.4 AND

AND[.W] Logical AND of source word with destination word

AND.B Logical AND of source byte with destination byte

Syntax AND src,dst OR AND.W src,dst
AND.B src,dst

Operation src .and. dst → dst

Description The source operand and the destination operand are logically ANDed. The result is placed into the destination. The source operand is not affected.

Status Bits N: Set if result is negative (MSB = 1), reset if positive (MSB = 0)
Z: Set if result is zero, reset otherwise
C: Set if the result is not zero, reset otherwise. C = (.not. Z)
V: Reset

Mode Bits OSCOFF, CPUOFF, and GIE are not affected.

Example The bits set in R5 (16-bit data) are used as a mask (AA55h) for the word TOM located in the lower 64 K. If the result is zero, a branch is taken to label TONI. R5.19:16 = 0

```
MOV    #AA55h,R5      ; Load 16-bit mask to R5
AND    R5,&TOM         ; TOM .and. R5 -> TOM
JZ     TONI           ; Jump if result 0
...                               ; Result > 0
```

or shorter:

```
AND    #AA55h,&TOM    ; TOM .and. AA55h -> TOM
JZ     TONI           ; Jump if result 0
```

Example A table byte pointed to by R5 (20-bit address) is logically ANDed with R6. R5 is incremented by 1 after the fetching of the byte. R6.19:8 = 0

```
AND.B  @R5+,R6        ; AND table byte with R6. R5 + 1
```

1.6.2.5 BIC

BIC[W]	Clear bits set in source word in destination word
BIC.B	Clear bits set in source byte in destination byte
Syntax	BIC src,dst OR BIC.W src,dst BIC.B src,dst
Operation	(.not. src) .and. dst → dst
Description	The inverted source operand and the destination operand are logically ANDed. The result is placed into the destination. The source operand is not affected.
Status Bits	N: Not affected Z: Not affected C: Not affected V: Not affected
Mode Bits	OSCOFF, CPUOFF, and GIE are not affected.
Example	The bits 15:14 of R5 (16-bit data) are cleared. R5.19:16 = 0
	<pre>BIC #0C000h,R5 ; Clear R5.19:14 bits</pre>
Example	A table word pointed to by R5 (20-bit address) is used to clear bits in R7. R7.19:16 = 0
	<pre>BIC.W @R5,R7 ; Clear bits in R7 set in @R5</pre>
Example	A table byte pointed to by R5 (20-bit address) is used to clear bits in Port1.
	<pre>BIC.B @R5,&P1OUT ; Clear I/O port P1 bits set in @R5</pre>

1.6.2.6 BIS

BIS[.W] Set bits set in source word in destination word

BIS.B Set bits set in source byte in destination byte

Syntax `BIS src,dst` OR `BIS.W src,dst`
`BIS.B src,dst`

Operation `src .or. dst → dst`

Description The source operand and the destination operand are logically ORed. The result is placed into the destination. The source operand is not affected.

Status Bits N: Not affected
 Z: Not affected
 C: Not affected
 V: Not affected

Mode Bits OSCOFF, CPUOFF, and GIE are not affected.

Example Bits 15 and 13 of R5 (16-bit data) are set to one. `R5.19:16 = 0`

```
BIS    #A000h,R5        ; Set R5 bits
```

Example A table word pointed to by R5 (20-bit address) is used to set bits in R7. `R7.19:16 = 0`

```
BIS.W  @R5,R7          ; Set bits in R7
```

Example A table byte pointed to by R5 (20-bit address) is used to set bits in Port1. R5 is incremented by 1 afterwards.

```
BIS.B  @R5+,&P1OUT     ; Set I/O port P1 bits. R5 + 1
```

1.6.2.7 BIT

BIT[.W]	Test bits set in source word in destination word
BIT.B	Test bits set in source byte in destination byte
Syntax	BIT src,dst OR BIT.W src,dst BIT.B src,dst
Operation	src .and. dst
Description	The source operand and the destination operand are logically ANDed. The result affects only the status bits in SR. Register mode: the register bits Rdst.19:16 (.W) resp. Rdst. 19:8 (.B) are not cleared!
Status Bits	N: Set if result is negative (MSB = 1), reset if positive (MSB = 0) Z: Set if result is zero, reset otherwise C: Set if the result is not zero, reset otherwise. C = (.not. Z) V: Reset
Mode Bits	OSCOFF, CPUOFF, and GIE are not affected.
Example	Test if one (or both) of bits 15 and 14 of R5 (16-bit data) is set. Jump to label TONI if this is the case. R5.19:16 are not affected.

```

BIT   #C000h,R5           ; Test R5.15:14 bits
JNZ   TONI                ; At least one bit is set in R5
...   ; Both bits are reset

```

Example A table word pointed to by R5 (20-bit address) is used to test bits in R7. Jump to label TONI if at least one bit is set. R7.19:16 are not affected.

```

BIT.W @R5,R7             ; Test bits in R7
JC    TONI               ; At least one bit is set
...   ; Both are reset

```

Example A table byte pointed to by R5 (20-bit address) is used to test bits in output Port1. Jump to label TONI if no bit is set. The next table byte is addressed.

```

BIT.B @R5+,&P1OUT        ; Test I/O port P1 bits. R5 + 1
JNC   TONI               ; No corresponding bit is set
...   ; At least one bit is set

```

1.6.2.8 BR, BRANCH

* BR, BRANCH	Branch to destination in lower 64K address space
Syntax	BR dst
Operation	dst → PC
Emulation	MOV dst,PC
Description	An unconditional branch is taken to an address anywhere in the lower 64K address space. All source addressing modes can be used. The branch instruction is a word instruction.
Status Bits	Status bits are not affected.
Example	Examples for all addressing modes are given.

BR	#EXEC	; Branch to label EXEC or direct branch (for example #0A4h) ; Core instruction MOV @PC+,PC
BR	EXEC	; Branch to the address contained in EXEC ; Core instruction MOV X(PC),PC ; Indirect address
BR	&EXEC	; Branch to the address contained in absolute ; address EXEC ; Core instruction MOV X(0),PC ; Indirect address
BR	R5	; Branch to the address contained in R5 ; Core instruction MOV R5,PC ; Indirect R5
BR	@R5	; Branch to the address contained in the word ; pointed to by R5. ; Core instruction MOV @R5,PC ; Indirect, indirect R5
BR	@R5+	; Branch to the address contained in the word pointed ; to by R5 and increment pointer in R5 afterwards. ; The next time-S/W flow uses R5 pointer-it can ; alter program execution due to access to ; next address in a table pointed to by R5 ; Core instruction MOV @R5,PC ; Indirect, indirect R5 with autoincrement
BR	X(R5)	; Branch to the address contained in the address ; pointed to by R5 + X (for example table with address ; starting at X). X can be an address or a label ; Core instruction MOV X(R5),PC ; Indirect, indirect R5 + X

1.6.2.9 CALL

CALL	Call a subroutine in lower 64 K
Syntax	CALL dst
Operation	dst → tmp 16-bit dst is evaluated and stored SP – 2 → SP PC → @SP updated PC with return address to TOS tmp → PC saved 16-bit dst to PC
Description	A subroutine call is made from an address in the lower 64 K to a subroutine address in the lower 64 K. All seven source addressing modes can be used. The call instruction is a word instruction. The return is made with the RET instruction.
Status Bits	Status bits are not affected. PC.19:16 cleared (address in lower 64 K)
Mode Bits	OSCOFF, CPUOFF, and GIE are not affected.
Examples	Examples for all addressing modes are given. Immediate Mode: Call a subroutine at label EXEC (lower 64 K) or call directly to address.

```
CALL #EXEC           ; Start address EXEC
CALL #0AA04h        ; Start address 0AA04h
```

Symbolic Mode: Call a subroutine at the 16-bit address contained in address EXEC. EXEC is located at the address (PC + X) where X is within PC ± 32 K.

```
CALL EXEC           ; Start address at @EXEC. z16(PC)
```

Absolute Mode: Call a subroutine at the 16-bit address contained in absolute address EXEC in the lower 64 K.

```
CALL &EXEC          ; Start address at @EXEC
```

Register mode: Call a subroutine at the 16-bit address contained in register R5.15:0.

```
CALL R5             ; Start address at R5
```

Indirect Mode: Call a subroutine at the 16-bit address contained in the word pointed to by register R5 (20-bit address).

```
CALL @R5           ; Start address at @R5
```


1.6.2.10 CLR

*** CLR[.W]** Clear destination

*** CLR.B** Clear destination

Syntax CLR dst *or* CLR.W dst
CLR.B dst

Operation 0 → dst

Emulation MOV #0,dst
MOV.B #0,dst

Description The destination operand is cleared.

Status Bits Status bits are not affected.

Example RAM word TONI is cleared.

```
CLR    TONI    ; 0 -> TONI
```

Example Register R5 is cleared.

```
CLR    R5
```

Example RAM byte TONI is cleared.

```
CLR.B  TONI    ; 0 -> TONI
```

1.6.2.11 CLRC

* CLRC	Clear carry bit
Syntax	CLRC
Operation	0 → C
Emulation	BIC #1,SR
Description	The carry bit (C) is cleared. The clear carry instruction is a word instruction.
Status Bits	N: Not affected Z: Not affected C: Cleared V: Not affected
Mode Bits	OSCOFF, CPUOFF, and GIE are not affected.
Example	The 16-bit decimal counter pointed to by R13 is added to a 32-bit counter pointed to by R12.

```

CLRC                ; C=0: defines start
DADD @R13,0(R12)    ; add 16-bit counter to low word of 32-bit counter
DADC 2(R12)         ; add carry to high word of 32-bit counter

```

1.6.2.12 CLRN

* CLRN	Clear negative bit
Syntax	CLRN
Operation	0 → N or (.NOT.src .AND. dst → dst)
Emulation	BIC #4,SR
Description	The constant 04h is inverted (0FFFBh) and is logically ANDed with the destination operand. The result is placed into the destination. The clear negative bit instruction is a word instruction.
Status Bits	N: Reset to 0 Z: Not affected C: Not affected V: Not affected
Mode Bits	OSCOFF, CPUOFF, and GIE are not affected.
Example	The negative bit in the SR is cleared. This avoids special treatment with negative numbers of the subroutine called.

```

                CLRN
                CALL SUBR
                .....
                .....
SUBR            JN      SUBRET      ; If input is negative: do nothing and return
                .....
                .....
                .....
SUBRET         RET
    
```

1.6.2.13 CLRZ

* CLRZ	Clear zero bit
Syntax	CLRZ
Operation	0 → Z or (.NOT.src .AND. dst → dst)
Emulation	BIC #2,SR
Description	The constant 02h is inverted (0FFFDh) and logically ANDed with the destination operand. The result is placed into the destination. The clear zero bit instruction is a word instruction.
Status Bits	N: Not affected Z: Reset to 0 C: Not affected V: Not affected
Mode Bits	OSCOFF, CPUOFF, and GIE are not affected.
Example	The zero bit in the SR is cleared.

```
CLRZ
```

Indirect, Auto-Increment mode: Call a subroutine at the 16-bit address contained in the word pointed to by register R5 (20-bit address) and increment the 16-bit address in R5 afterwards by 2. The next time the software uses R5 as a pointer, it can alter the program execution due to access to the next word address in the table pointed to by R5.

```
CALL @R5+          ; Start address at @R5. R5 + 2
```

Indexed mode: Call a subroutine at the 16-bit address contained in the 20-bit address pointed to by register (R5 + X); for example, a table with addresses starting at X. The address is within the lower 64KB. X is within ±32KB.

```
CALL X(R5)        ; Start address at @(R5+X). z16(R5)
```

1.6.2.14 CMP

CMP[.W]	Compare source word and destination word
CMP.B	Compare source byte and destination byte
Syntax	CMP src,dst OR CMP.W src,dst CMP.B src,dst
Operation	(.not.src) + 1 + dst or dst – src
Description	The source operand is subtracted from the destination operand. This is made by adding the 1s complement of the source + 1 to the destination. The result affects only the status bits in SR. Register mode: the register bits Rdst.19:16 (.W) resp. Rdst. 19:8 (.B) are not cleared.
Status Bits	N: Set if result is negative (src > dst), reset if positive (src ≤ dst) Z: Set if result is zero (src = dst), reset otherwise (src ≠ dst) C: Set if there is a carry from the MSB, reset otherwise V: Set if the subtraction of a negative source operand from a positive destination operand delivers a negative result, or if the subtraction of a positive source operand from a negative destination operand delivers a positive result, reset otherwise (no overflow).
Mode Bits	OSCOFF, CPUOFF, and GIE are not affected.
Example	Compare word EDE with a 16-bit constant 1800h. Jump to label TONI if EDE equals the constant. The address of EDE is within PC + 32 K.
	<pre> CMP #01800h,EDE ; Compare word EDE with 1800h JEQ TONI ; EDE contains 1800h ... ; Not equal </pre>
Example	A table word pointed to by (R5 + 10) is compared with R7. Jump to label TONI if R7 contains a lower, signed 16-bit number. R7.19:16 is not cleared. The address of the source operand is a 20-bit address in full memory range.
	<pre> CMP.W 10(R5),R7 ; Compare two signed numbers JL TONI ; R7 < 10(R5) ... ; R7 >= 10(R5) </pre>
Example	A table byte pointed to by R5 (20-bit address) is compared to the value in output Port1. Jump to label TONI if values are equal. The next table byte is addressed.
	<pre> CMP.B @R5+,&P1OUT ; Compare P1 bits with table. R5 + 1 JEQ TONI ; Equal contents ... ; Not equal </pre>

1.6.2.15 DADC

* **DADC[W]** Add carry decimally to destination

* **DADC.B** Add carry decimally to destination

Syntax DADC dst Or DADC.W dst
DADC.B dst

Operation dst + C → dst (decimally)

Emulation DADD #0, dst
DADD.B #0, dst

Description The carry bit (C) is added decimally to the destination.

Status Bits
 N: Set if MSB is 1
 Z: Set if dst is 0, reset otherwise
 C: Set if destination increments from 9999 to 0000, reset otherwise
 Set if destination increments from 99 to 00, reset otherwise
 V: Undefined

Mode Bits OSCOFF, CPUOFF, and GIE are not affected.

Example The four-digit decimal number contained in R5 is added to an eight-digit decimal number pointed to by R8.

```

CLRC                ; Reset carry
                    ; next instruction's start condition is defined
DADD R5,0(R8)      ; Add LSDs + C
DADC 2(R8)          ; Add carry to MSD

```

Example The two-digit decimal number contained in R5 is added to a four-digit decimal number pointed to by R8.

```

CLRC                ; Reset carry
                    ; next instruction's start condition is defined
DADD.B R5,0(R8)    ; Add LSDs + C
DADC 1(R8)          ; Add carry to MSDs

```

1.6.2.16 DADD

* **DADD[.W]** Add source word and carry decimally to destination word

* **DADD.B** Add source byte and carry decimally to destination byte

Syntax DADD src,dst OR DADD.W src,dst
DADD.B src,dst

Operation src + dst + C → dst (decimally)

Description The source operand and the destination operand are treated as two (.B) or four (.W) binary coded decimals (BCD) with positive signs. The source operand and the carry bit C are added decimally to the destination operand. The source operand is not affected. The previous content of the destination is lost. The result is not defined for non-BCD numbers.

Status Bits N: Set if MSB of result is 1 (word > 7999h, byte > 79h), reset if MSB is 0
Z: Set if result is zero, reset otherwise
C: Set if the BCD result is too large (word > 9999h, byte > 99h), reset otherwise
V: Undefined

Mode Bits OSCOFF, CPUOFF, and GIE are not affected.

Example Decimal 10 is added to the 16-bit BCD counter DECCNTR.

```
DADD #10h,&DECCNTR ; Add 10 to 4-digit BCD counter
```

Example The eight-digit BCD number contained in 16-bit RAM addresses BCD and BCD+2 is added decimally to an eight-digit BCD number contained in R4 and R5 (BCD+2 and R5 contain the MSDs). The carry C is added, and cleared.

```
CLRC ; Clear carry
DADD.W &BCD,R4 ; Add LSDs. R4.19:16 = 0
DADD.W &BCD+2,R5 ; Add MSDs with carry. R5.19:16 = 0
JC OVERFLOW ; Result >9999,9999: go to error routine
... ; Result ok
```

Example The two-digit BCD number contained in word BCD (16-bit address) is added decimally to a two-digit BCD number contained in R4. The carry C is added, also. R4.19:8 = 0

```
CLRC ; Clear carry
DADD.B &BCD,R4 ; Add BCD to R4 decimally.
R4: 0,00ddh
```

1.6.2.17 DEC

* DEC.W]	Decrement destination
* DEC.B	Decrement destination
Syntax	DEC dst OR DEC.W dst DEC.B dst
Operation	dst - 1 → dst
Emulation	SUB #1, dst SUB.B #1, dst
Description	The destination operand is decremented by one. The original contents are lost.
Status Bits	N: Set if result is negative, reset if positive Z: Set if dst contained 1, reset otherwise C: Reset if dst contained 0, set otherwise V: Set if an arithmetic overflow occurs, otherwise reset. Set if initial value of destination was 08000h, otherwise reset. Set if initial value of destination was 080h, otherwise reset.
Mode Bits	OSCOFF, CPUOFF, and GIE are not affected.
Example	R10 is decremented by 1.

```

DEC    R10                ; Decrement R10

; Move a block of 255 bytes from memory location starting with EDE to
; memory location starting with TONI. Tables should not overlap: start of
; destination address TONI must not be within the range EDE to EDE+0FEh

MOV    #EDE, R6
MOV    #255, R10
L$1   MOV.B  @R6+, TONI-EDE-1(R6)
DEC    R10
JNZ   L$1

```

Do not transfer tables using the routine above with the overlap shown in [Figure 1-36](#).

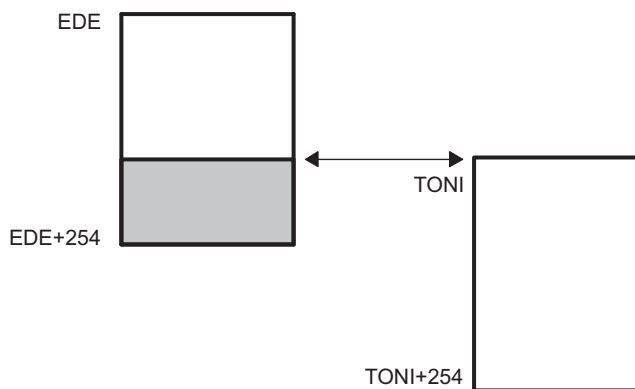


Figure 1-36. Decrement Overlap

1.6.2.18 DECD

* DECD[.W]	Double-decrement destination
* DECD.B	Double-decrement destination
Syntax	DECD dst Or DECD.W dst DECD.B dst
Operation	dst – 2 → dst
Emulation	SUB #2, dst SUB.B #2, dst
Description	The destination operand is decremented by two. The original contents are lost.
Status Bits	N: Set if result is negative, reset if positive Z: Set if dst contained 2, reset otherwise C: Reset if dst contained 0 or 1, set otherwise V: Set if an arithmetic overflow occurs, otherwise reset Set if initial value of destination was 08001 or 08000h, otherwise reset Set if initial value of destination was 081 or 080h, otherwise reset
Mode Bits	OSCOFF, CPUOFF, and GIE are not affected.
Example	R10 is decremented by 2.

```
DECD    R10                ; Decrement R10 by two
```

```
; Move a block of 255 bytes from memory location starting with EDE to
; memory location starting with TONI.
; Tables should not overlap: start of destination address TONI must not
; be within the range EDE to EDE+0FEh
```

```
MOV     #EDE, R6
MOV     #255, R10
L$1    MOV.B  @R6+, TONI-EDE-2(R6)
DECD   R10
JNZ    L$1
```

Example Memory at location LEO is decremented by two.

```
DECD.B  LEO                ; Decrement MEM(LEO)
```

Decrement status byte STATUS by two

```
DECD.B  STATUS
```

1.6.2.19 DINT

* DINT	Disable (general) interrupts
Syntax	DINT
Operation	0 → GIE or (0FFF7h .AND. SR → SR / .NOT.src .AND. dst → dst)
Emulation	BIC #8,SR
Description	All interrupts are disabled. The constant 08h is inverted and logically ANDed with the SR. The result is placed into the SR.
Status Bits	Status bits are not affected.
Mode Bits	GIE is reset. OSCOFF and CPUOFF are not affected.
Example	The general interrupt enable (GIE) bit in the SR is cleared to allow a nondisrupted move of a 32-bit counter. This ensures that the counter is not modified during the move by any interrupt.

```

DINT                ; All interrupt events using the GIE bit are disabled
NOP                 ; Required due to pipelined CPU architecture
MOV    COUNTHI,R5   ; Copy counter
MOV    COUNTLO,R6
EINT                ; All interrupt events using the GIE bit are enabled

```

NOTE: Disable interrupt

Due to the pipelined CPU architecture, clearing the general interrupt enable (GIE) requires special care.

- Include at least one instruction between DINT and the start of an code sequence that requires protection from interrupts. For example: Insert a NOP instruction after the DINT.
- Never clear the general interrupt enable (GIE) immediately after setting it. Insert at least one instruction in between such sequence.

The rules above apply to all instructions that clear the general interrupt enable bit. Not following these rules might result in unexpected CPU execution.

1.6.2.20 EINT

* EINT	Enable (general) interrupts
Syntax	EINT
Operation	1 → GIE or (0008h .OR. SR → SR / .src .OR. dst → dst)
Emulation	BIS #8,SR
Description	All interrupts are enabled. The constant #08h and the SR are logically ORed. The result is placed into the SR.
Status Bits	Status bits are not affected.
Mode Bits	GIE is set. OSCOFF and CPUOFF are not affected.
Example	The general interrupt enable (GIE) bit in the SR is set.

```

PUSH.B    &PLIN
BIC.B     @SP,&P1IFG ; Reset only accepted flags
NOP       ; Required due to pipelined CPU architecture
EINT      ; Preset port 1 interrupt flags stored on stack
          ; other interrupts are allowed

BIT       #Mask,@SP
JEQ       MaskOK    ; Flags are present identically to mask: jump
.....
MaskOK    BIC       #Mask,@SP
          .....
          INCD     SP ; Housekeeping: inverse to PUSH instruction
          ; at the start of interrupt subroutine. Corrects
          ; the stack pointer.

RETI
    
```

NOTE: Enable interrupt

Due to the pipelined CPU architecture, setting the general interrupt enable (GIE) requires special care.

- The instruction immediately after the enable interrupts instruction (EINT) is always executed, even if an interrupt service request is pending.
- Include at least one instruction between the clear of an interrupt enable or interrupt flag and the EINT instruction. For example: Insert a NOP instruction in front of the EINT instruction.
- Never clear the general interrupt enable (GIE) immediately after setting it. Insert at least one instruction in between such sequence.

The rules above apply to all instructions that set the general interrupt enable bit. Not following these rules might result in unexpected CPU execution.

1.6.2.21 INC

* **INC.W** Increment destination

* **INC.B** Increment destination

Syntax `INC dst OR INC.W dst`
`INC.B dst`

Operation `dst + 1 → dst`

Emulation `ADD #1, dst`

Description The destination operand is incremented by one. The original contents are lost.

Status Bits

- N: Set if result is negative, reset if positive
- Z: Set if dst contained 0FFFFh, reset otherwise
Set if dst contained 0FFh, reset otherwise
- C: Set if dst contained 0FFFFh, reset otherwise
Set if dst contained 0FFh, reset otherwise
- V: Set if dst contained 07FFFh, reset otherwise
Set if dst contained 07Fh, reset otherwise

Mode Bits OSCOFF, CPUOFF, and GIE are not affected.

Example The status byte, STATUS, of a process is incremented. When it is equal to 11, a branch to OVFL is taken.

```
INC.B STATUS
CMP.B #11, STATUS
JEQ OVFL
```

1.6.2.22 INCD

* INCD[.W]	Double-increment destination
* INCD.B	Double-increment destination
Syntax	INCD dst Or INCD.W dst INCD.B dst
Operation	dst + 2 → dst
Emulation	ADD #2, dst
Description	The destination operand is incremented by two. The original contents are lost.
Status Bits	N: Set if result is negative, reset if positive Z: Set if dst contained 0FFFEh, reset otherwise Set if dst contained 0FEh, reset otherwise C: Set if dst contained 0FFFEh or 0FFFFh, reset otherwise Set if dst contained 0FEh or 0FFh, reset otherwise V: Set if dst contained 07FFEh or 07FFFh, reset otherwise Set if dst contained 07Eh or 07Fh, reset otherwise
Mode Bits	OSCOFF, CPUOFF, and GIE are not affected.
Example	The item on the top of the stack (TOS) is removed without using a register.

```

.....
PUSH  R5      ; R5 is the result of a calculation, which is stored
          ; in the system stack
INCD  SP      ; Remove TOS by double-increment from stack
          ; Do not use INCD.B, SP is a word-aligned register
RET
    
```

Example The byte on the top of the stack is incremented by two.

```
INCD.B 0(SP) ; Byte on TOS is increment by two
```

1.6.2.23 INV

* INV[.W]	Invert destination
* INV.B	Invert destination
Syntax	INV dst OR INV.W dst INV.B dst
Operation	.not.dst → dst
Emulation	XOR #0FFFFh, dst XOR.B #0FFh, dst
Description	The destination operand is inverted. The original contents are lost.
Status Bits	N: Set if result is negative, reset if positive Z: Set if dst contained 0FFFFh, reset otherwise Set if dst contained 0FFh, reset otherwise C: Set if result is not zero, reset otherwise (= .NOT. Zero) V: Set if initial destination operand was negative, otherwise reset
Mode Bits	OSCOFF, CPUOFF, and GIE are not affected.
Example	Content of R5 is negated (twos complement).

```

MOV    #00AEh, R5    ; R5 = 000AEh
INV    R5             ; Invert R5,   R5 = 0FF51h
INC    R5             ; R5 is now negated, R5 = 0FF52h

```

Example Content of memory byte LEO is negated.

```

MOV.B  #0AEh, LEO    ; MEM(LEO) = 0AEh
INV.B  LEO           ; Invert LEO,   MEM(LEO) = 051h
INC.B  LEO           ; MEM(LEO) is negated, MEM(LEO) = 052h

```

1.6.2.24 JC, JHS

JC	Jump if carry
JHS	Jump if higher or same (unsigned)
Syntax	JC label JHS label
Operation	If C = 1: PC + (2 × Offset) → PC If C = 0: execute the following instruction
Description	The carry bit C in the SR is tested. If it is set, the signed 10-bit word offset contained in the instruction is multiplied by two, sign extended, and added to the 20-bit PC. This means a jump in the range –511 to +512 words relative to the PC in the full memory range. If C is reset, the instruction after the jump is executed. JC is used for the test of the carry bit C. JHS is used for the comparison of unsigned numbers.
Status Bits	Status bits are not affected
Mode Bits	OSCOFF, CPUOFF, and GIE are not affected.
Example	The state of the port 1 pin P1IN.1 bit defines the program flow.

```

BIT.B #2,&P1IN      ; Port 1, bit 1 set? Bit -> C
JC    Label1       ; Yes, proceed at Label1
...      ; No, continue
    
```

Example If R5 ≥ R6 (unsigned), the program continues at Label2.

```

CMP   R6,R 5      ; Is R5 >= R6? Info to C
JHS   Label2      ; Yes, C = 1
...      ; No, R5 < R6. Continue
    
```

Example If R5 ≥ 12345h (unsigned operands), the program continues at Label2.

```

CMPA  #12345h,R5  ; Is R5 >= 12345h? Info to C
JHS   Label2      ; Yes, 12344h < R5 <= F,FFFh. C = 1
...      ; No, R5 < 12345h. Continue
    
```

1.6.2.25 JEQ, JZ

JEQ	Jump if equal
JZ	Jump if zero
Syntax	JEQ label JZ label
Operation	If Z = 1: PC + (2 × Offset) → PC If Z = 0: execute following instruction
Description	The zero bit Z in the SR is tested. If it is set, the signed 10-bit word offset contained in the instruction is multiplied by two, sign extended, and added to the 20-bit PC. This means a jump in the range –511 to +512 words relative to the PC in the full memory range. If Z is reset, the instruction after the jump is executed. JZ is used for the test of the zero bit Z. JEQ is used for the comparison of operands.
Status Bits	Status bits are not affected
Mode Bits	OSCOFF, CPUOFF, and GIE are not affected.
Example	The state of the P2IN.0 bit defines the program flow.

```

BIT.B  #1,&P2IN      ; Port 2, bit 0 reset?
JZ     Label1       ; Yes, proceed at Label1
...    ; No, set, continue

```

Example If R5 = 15000h (20-bit data), the program continues at Label2.

```

CMPA   #15000h,R5   ; Is R5 = 15000h? Info to SR
JEQ    Label2       ; Yes, R5 = 15000h. Z = 1
...    ; No, R5 not equal 15000h. Continue

```

Example R7 (20-bit counter) is incremented. If its content is zero, the program continues at Label4.

```

ADDA   #1,R7        ; Increment R7
JZ     Label4       ; Zero reached: Go to Label4
...    ; R7 not equal 0. Continue here.

```


1.6.2.26 JGE

JGE	Jump if greater or equal (signed)
Syntax	JGE label
Operation	If (N .xor. V) = 0: PC + (2 × Offset) → PC If (N .xor. V) = 1: execute following instruction
Description	<p>The negative bit N and the overflow bit V in the SR are tested. If both bits are set or both are reset, the signed 10-bit word offset contained in the instruction is multiplied by two, sign extended, and added to the 20-bit PC. This means a jump in the range -511 to +512 words relative to the PC in full Memory range. If only one bit is set, the instruction after the jump is executed.</p> <p>JGE is used for the comparison of signed operands: also for incorrect results due to overflow, the decision made by the JGE instruction is correct.</p> <p>Note that JGE emulates the nonimplemented JP (jump if positive) instruction if used after the instructions AND, BIT, RRA, SCTX, and TST. These instructions clear the V bit.</p>
Status Bits	Status bits are not affected.
Mode Bits	OSCOFF, CPUOFF, and GIE are not affected.
Example	If byte EDE (lower 64 K) contains positive data, go to Label1. Software can run in the full memory range.

```
TST.B   &EDE           ; Is EDE positive? V <- 0
JGE     Label1        ; Yes, JGE emulates JP
...     ; No, 80h <= EDE <= FFh
```

Example	If the content of R6 is greater than or equal to the memory pointed to by R7, the program continues a Label5. Signed data. Data and program in full memory range.
----------------	---

```
CMP     @R7,R6        ; Is R6 >= @R7?
JGE     Label5        ; Yes, go to Label5
...     ; No, continue here
```

Example	If R5 ≥ 12345h (signed operands), the program continues at Label2. Program in full memory range.
----------------	--

```
CMPA   #12345h,R5    ; Is R5 >= 12345h?
JGE     Label2        ; Yes, 12344h < R5 <= 7FFFFh
...     ; No, 80000h <= R5 < 12345h
```

1.6.2.27 JL**JL** Jump if less (signed)**Syntax** JL label**Operation** If (N .xor. V) = 1: PC + (2 × Offset) → PC
If (N .xor. V) = 0: execute following instruction**Description** The negative bit N and the overflow bit V in the SR are tested. If only one is set, the signed 10-bit word offset contained in the instruction is multiplied by two, sign extended, and added to the 20-bit PC. This means a jump in the range –511 to +512 words relative to the PC in full memory range. If both bits N and V are set or both are reset, the instruction after the jump is executed.

JL is used for the comparison of signed operands: also for incorrect results due to overflow, the decision made by the JL instruction is correct.

Status Bits Status bits are not affected.**Mode Bits** OSCOFF, CPUOFF, and GIE are not affected.**Example** If byte EDE contains a smaller, signed operand than byte TONI, continue at Label1. The address EDE is within PC ± 32 K.

```

CMP.B  &TONI,EDE      ; Is EDE < TONI
JL     Label1         ; Yes
...                               ; No, TONI <= EDE

```

Example If the signed content of R6 is less than the memory pointed to by R7 (20-bit address), the program continues at Label5. Data and program in full memory range.

```

CMP    @R7,R6         ; Is R6 < @R7?
JL     Label5         ; Yes, go to Label5
...                               ; No, continue here

```

Example If R5 < 12345h (signed operands), the program continues at Label2. Data and program in full memory range.

```

CMPA   #12345h,R5     ; Is R5 < 12345h?
JL     Label2         ; Yes, 80000h =< R5 < 12345h
...                               ; No, 12344h < R5 <= 7FFFFh

```

1.6.2.28 JMP

JMP Jump unconditionally

Syntax JMP label

Operation PC + (2 × Offset) → PC

Description The signed 10-bit word offset contained in the instruction is multiplied by two, sign extended, and added to the 20-bit PC. This means an unconditional jump in the range –511 to +512 words relative to the PC in the full memory. The JMP instruction may be used as a BR or BRA instruction within its limited range relative to the PC.

Status Bits Status bits are not affected

Mode Bits OSCOFF, CPUOFF, and GIE are not affected.

Example The byte STATUS is set to 10. Then a jump to label MAINLOOP is made. Data in lower 64 K, program in full memory range.

```
MOV.B #10,&STATUS ; Set STATUS to 10
JMP MAINLOOP ; Go to main loop
```

Example The interrupt vector TAIV of Timer_A3 is read and used for the program flow. Program in full memory range, but interrupt handlers always starts in lower 64 K.

```
ADD &TAIV,PC ; Add Timer_A interrupt vector to PC
RETI ; No Timer_A interrupt pending
JMP IHCCR1 ; Timer block 1 caused interrupt
JMP IHCCR2 ; Timer block 2 caused interrupt
RETI ; No legal interrupt, return
```

1.6.2.29 JN**JN** Jump if negative**Syntax** JN label**Operation** If N = 1: PC + (2 × Offset) → PC
If N = 0: execute following instruction**Description** The negative bit N in the SR is tested. If it is set, the signed 10-bit word offset contained in the instruction is multiplied by two, sign extended, and added to the 20-bit program PC. This means a jump in the range -511 to +512 words relative to the PC in the full memory range. If N is reset, the instruction after the jump is executed.**Status Bits** Status bits are not affected.**Mode Bits** OSCOFF, CPUOFF, and GIE are not affected.**Example** The byte COUNT is tested. If it is negative, program execution continues at Label0. Data in lower 64 K, program in full memory range.

```
TST.B  &COUNT    ; Is byte COUNT negative?
JN     Label0     ; Yes, proceed at Label0
...     ; COUNT >= 0
```

Example R6 is subtracted from R5. If the result is negative, program continues at Label2. Program in full memory range.

```
SUB    R6,R5      ; R5 - R6 -> R5
JN     Label2     ; R5 is negative: R6 > R5 (N = 1)
...     ; R5 >= 0. Continue here.
```

Example R7 (20-bit counter) is decremented. If its content is below zero, the program continues at Label4. Program in full memory range.

```
SUBA   #1,R7     ; Decrement R7
JN     Label4     ; R7 < 0: Go to Label4
...     ; R7 >= 0. Continue here.
```

1.6.2.30 JNC, JLO

JNC	Jump if no carry
JLO	Jump if lower (unsigned)
Syntax	JNC label JLO label
Operation	If C = 0: PC + (2 × Offset) → PC If C = 1: execute following instruction
Description	The carry bit C in the SR is tested. If it is reset, the signed 10-bit word offset contained in the instruction is multiplied by two, sign extended, and added to the 20-bit PC. This means a jump in the range –511 to +512 words relative to the PC in the full memory range. If C is set, the instruction after the jump is executed. JNC is used for the test of the carry bit C. JLO is used for the comparison of unsigned numbers.
Status Bits	Status bits are not affected.
Mode Bits	OSCOFF, CPUOFF, and GIE are not affected.
Example	If byte EDE < 15, the program continues at Label2. Unsigned data. Data in lower 64 K, program in full memory range.

```

CMP.B #15,&EDE ; Is EDE < 15? Info to C
JLO Label2 ; Yes, EDE < 15. C = 0
... ; No, EDE >= 15. Continue
    
```

Example	The word TONI is added to R5. If no carry occurs, continue at Label0. The address of TONI is within PC ± 32 K.
----------------	--

```

ADD TONI,R5 ; TONI + R5 -> R5. Carry -> C
JNC Label0 ; No carry
... ; Carry = 1: continue here
    
```

1.6.2.31 JNZ, JNE**JNZ** Jump if not zero**JNE** Jump if not equal**Syntax** JNZ label

JNE label

Operation If Z = 0: PC + (2 × Offset) → PC
If Z = 1: execute following instruction**Description** The zero bit Z in the SR is tested. If it is reset, the signed 10-bit word offset contained in the instruction is multiplied by two, sign extended, and added to the 20-bit PC. This means a jump in the range –511 to +512 words relative to the PC in the full memory range. If Z is set, the instruction after the jump is executed.

JNZ is used for the test of the zero bit Z.

JNE is used for the comparison of operands.

Status Bits Status bits are not affected.**Mode Bits** OSCOFF, CPUOFF, and GIE are not affected.**Example** The byte STATUS is tested. If it is not zero, the program continues at Label3. The address of STATUS is within PC ± 32 K.

```
TST.B STATUS          ; Is STATUS = 0?
JNZ Label3           ; No, proceed at Label3
...                  ; Yes, continue here
```

Example If word EDE ≠ 1500, the program continues at Label2. Data in lower 64 K, program in full memory range.

```
CMP #1500,&EDE        ; Is EDE = 1500? Info to SR
JNE Label2           ; No, EDE not equal 1500.
...                  ; Yes, R5 = 1500. Continue
```

Example R7 (20-bit counter) is decremented. If its content is not zero, the program continues at Label4. Program in full memory range.

```
SUBA #1,R7           ; Decrement R7
JNZ Label4           ; Zero not reached: Go to Label4
...                  ; Yes, R7 = 0. Continue here.
```

1.6.2.32 MOV

MOV[.W]	Move source word to destination word
MOV.B	Move source byte to destination byte
Syntax	MOV src,dst OR MOV.W src,dst MOV.B src,dst
Operation	src → dst
Description	The source operand is copied to the destination. The source operand is not affected.
Status Bits	N: Not affected Z: Not affected C: Not affected V: Not affected
Mode Bits	OSCOFF, CPUOFF, and GIE are not affected.
Example	Move a 16-bit constant 1800h to absolute address-word EDE (lower 64 K)

```
MOV    #01800h,&EDE           ; Move 1800h to EDE
```

Example The contents of table EDE (word data, 16-bit addresses) are copied to table TOM. The length of the tables is 030h words. Both tables reside in the lower 64 K.

```
MOV    #EDE,R10              ; Prepare pointer (16-bit address)
Loop   MOV    @R10+,TOM-EDE-2(R10) ; R10 points to both tables.
                                   ; R10+2
CMP    #EDE+60h,R10          ; End of table reached?
JLO    Loop                  ; Not yet
...    ; Copy completed
```

Example The contents of table EDE (byte data, 16-bit addresses) are copied to table TOM. The length of the tables is 020h bytes. Both tables may reside in full memory range, but must be within $R10 \pm 32$ K.

```
MOVA   #EDE,R10              ; Prepare pointer (20-bit)
MOV    #20h,R9               ; Prepare counter
Loop   MOV.B  @R10+,TOM-EDE-1(R10) ; R10 points to both tables.
                                   ; R10+1
DEC    R9                    ; Decrement counter
JNZ    Loop                  ; Not yet done
...    ; Copy completed
```

1.6.2.33 NOP

* NOP	No operation
Syntax	NOP
Operation	None
Emulation	MOV #0, R3
Description	No operation is performed. The instruction may be used for the elimination of instructions during the software check or for defined waiting times.
Status Bits	Status bits are not affected.

1.6.2.34 POP

* **POP[W]** Pop word from stack to destination

* **POP.B** Pop byte from stack to destination

Syntax POP dst

POP.B dst

Operation @SP → temp

SP + 2 → SP

temp → dst

Emulation MOV @SP+,dst OR MOV.W @SP+,dst

MOV.B @SP+,dst

Description The stack location pointed to by the SP (TOS) is moved to the destination. The SP is incremented by two afterwards.

Status Bits Status bits are not affected.

Example The contents of R7 and the SR are restored from the stack.

```
POP    R7        ; Restore R7
POP    SR        ; Restore status register
```

Example The contents of RAM byte LEO is restored from the stack.

```
POP.B  LEO      ; The low byte of the stack is moved to LEO.
```

Example The contents of R7 is restored from the stack.

```
POP.B  R7        ; The low byte of the stack is moved to R7,
                 ; the high byte of R7 is 00h
```

Example The contents of the memory pointed to by R7 and the SR are restored from the stack.

```
POP.B  0(R7)     ; The low byte of the stack is moved to the
                 ; the byte which is pointed to by R7
                 ; Example:   R7 = 203h
                 ;           Mem(R7) = low byte of system stack
                 ; Example:   R7 = 20Ah
                 ;           Mem(R7) = low byte of system stack
POP    SR        ; Last word on stack moved to the SR
```

NOTE: System stack pointer

The system SP is always incremented by two, independent of the byte suffix.

1.6.2.35 PUSH**PUSH[.W]** Save a word on the stack**PUSH.B** Save a byte on the stack**Syntax** PUSH dst OR PUSH.W dst
PUSH.B dst**Operation** SP – 2 → SP
dst → @SP**Description** The 20-bit SP is decremented by two. The operand is then copied to the RAM word addressed by the SP. A pushed byte is stored in the low byte; the high byte is not affected.**Status Bits** Status bits are not affected.**Mode Bits** OSCOFF, CPUOFF, and GIE are not affected.**Example** Save the two 16-bit registers R9 and R10 on the stack

```

PUSH    R9        ; Save R9 and R10 XXXXh
PUSH    R10       ; YYYh

```

Example Save the two bytes EDE and TONI on the stack. The addresses EDE and TONI are within PC ± 32 K.

```

PUSH.B  EDE      ; Save EDE   xxXXh
PUSH.B  TONI     ; Save TONI  xxYYh

```

1.6.2.36 RET

*** RET** Return from subroutine

Syntax RET

Operation @SP → PC.15:0 Saved PC to PC.15:0. PC.19:16 ← 0
 SP + 2 → SP

Description The 16-bit return address (lower 64 K), pushed onto the stack by a CALL instruction is restored to the PC. The program continues at the address following the subroutine call. The four MSBs of the PC.19:16 are cleared.

Status Bits Status bits are not affected.
 PC.19:16: Cleared

Mode Bits OSCOFF, CPUOFF, and GIE are not affected.

Example Call a subroutine SUBR in the lower 64 K and return to the address in the lower 64 K after the CALL.

```

CALL    #SUBR    ; Call subroutine starting at SUBR
...
SUBR    PUSH    R14    ; Save R14 (16 bit data)
...
        ; Subroutine code
SUBR    POP     R14    ; Restore R14
SUBR    RET     ; Return to lower 64 K
    
```

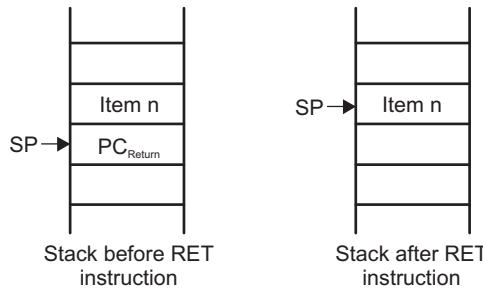


Figure 1-37. Stack After a RET Instruction

1.6.2.37 RETI

RETI	Return from interrupt
Syntax	RETI
Operation	<p>@SP → SR.15:0 Restore saved SR with PC.19:16 SP + 2 → SP</p> <p>@SP → PC.15:0 Restore saved PC.15:0 SP + 2 → SP Housekeeping</p>
Description	<p>The SR is restored to the value at the beginning of the interrupt service routine. This includes the four MSBs of the PC.19:16. The SP is incremented by two afterward. The 20-bit PC is restored from PC.19:16 (from same stack location as the status bits) and PC.15:0. The 20-bit PC is restored to the value at the beginning of the interrupt service routine. The program continues at the address following the last executed instruction when the interrupt was granted. The SP is incremented by two afterward. No interrupt flags are modified by this command.</p>
Status Bits	<p>N: Restored from stack C: Restored from stack Z: Restored from stack V: Restored from stack</p>
Mode Bits	OSCOFF, CPUOFF, and GIE are restored from stack.
Example	Interrupt handler in the lower 64 K. A 20-bit return address is stored on the stack.

```

INTRPT  PUSHM.A  #2,R14    ; Save R14 and R13 (20-bit data)
        ...           ; Interrupt handler code
        POPM.A   #2,R14    ; Restore R13 and R14 (20-bit data)
        RETI         ; Return to 20-bit address in full memory range

```

1.6.2.38 RLA

* **RLA[.W]** Rotate left arithmetically

* **RLA.B** Rotate left arithmetically

Syntax RLA dst OF RLA.W dst
RLA.B dst

Operation $C \leftarrow MSB \leftarrow MSB-1 \dots LSB+1 \leftarrow LSB \leftarrow 0$

Emulation ADD dst, dst
ADD.B dst, dst

Description The destination operand is shifted left one position as shown in Figure 1-38. The MSB is shifted into the carry bit (C) and the LSB is filled with 0. The RLA instruction acts as a signed multiplication by 2.
An overflow occurs if $dst \geq 04000h$ and $dst < 0C000h$ before operation is performed; the result has changed sign.

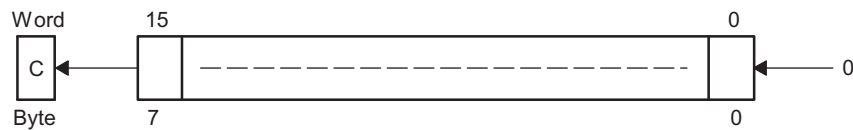


Figure 1-38. Destination Operand—Arithmetic Shift Left

An overflow occurs if $dst \geq 040h$ and $dst < 0C0h$ before the operation is performed; the result has changed sign.

Status Bits

- N: Set if result is negative, reset if positive
- Z: Set if result is zero, reset otherwise
- C: Loaded from the MSB
- V: Set if an arithmetic overflow occurs; the initial value is $04000h \leq dst < 0C000h$, reset otherwise
Set if an arithmetic overflow occurs; the initial value is $040h \leq dst < 0C0h$, reset otherwise

Mode Bits OSCOFF, CPUOFF, and GIE are not affected.

Example R7 is multiplied by 2.

```
RLA R7 ; Shift left R7 (x 2)
```

Example The low byte of R7 is multiplied by 4.

```
RLA.B R7 ; Shift left low byte of R7 (x 2)
RLA.B R7 ; Shift left low byte of R7 (x 4)
```

NOTE: RLA substitution

The assembler does not recognize the instructions:

```
RLA @R5+ RLA.B @R5+ RLA(.B) @R5
```

They must be substituted by:

```
ADD @R5+, -2(R5) ADD.B @R5+, -1(R5) ADD(.B) @R5
```

1.6.2.39 RLC

* **RLC[.W]** Rotate left through carry

* **RLC.B** Rotate left through carry

Syntax RLC dst OR RLC.W dst
RLC.B dst

Operation $C \leftarrow MSB \leftarrow MSB-1 \dots LSB+1 \leftarrow LSB \leftarrow C$

Emulation ADDC dst, dst

Description The destination operand is shifted left one position as shown in Figure 1-39. The carry bit (C) is shifted into the LSB, and the MSB is shifted into the carry bit (C).

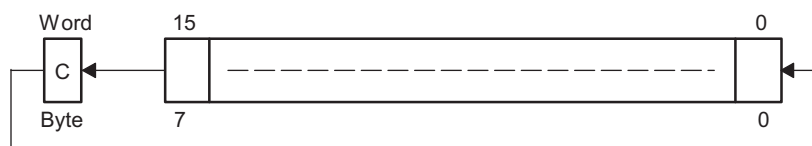


Figure 1-39. Destination Operand—Carry Left Shift

Status Bits

- N: Set if result is negative, reset if positive
- Z: Set if result is zero, reset otherwise
- C: Loaded from the MSB
- V: Set if an arithmetic overflow occurs; the initial value is $04000h \leq dst < 0C000h$, reset otherwise
Set if an arithmetic overflow occurs; the initial value is $040h \leq dst < 0C0h$, reset otherwise

Mode Bits OSCOFF, CPUOFF, and GIE are not affected.

Example R5 is shifted left one position.

```
RLC R5 ; (R5 x 2) + C -> R5
```

Example The input P1IN.1 information is shifted into the LSB of R5.

```
BIT.B #2, &P1IN ; Information -> Carry
RLC R5 ; Carry=P0in.1 -> LSB of R5
```

Example The MEM(LEO) content is shifted left one position.

```
RLC.B LEO ; Mem(LEO) x 2 + C -> Mem(LEO)
```

NOTE: RLA substitution

The assembler does not recognize the instructions:

```
RLC @R5+ RLC.B @R5+ RLC(.B) @R5
```

They must be substituted by:

```
ADDC @R5+, -2(R5) ADDC.B @R5+, -1(R5) ADDC(.B) @R5
```

1.6.2.40 RRA

RRA[W] Rotate right arithmetically destination word
RRA.B Rotate right arithmetically destination byte
Syntax RRA.B dst OR RRA.W dst
Operation MSB → MSB → MSB-1 → ... LSB+1 → LSB → C
Description The destination operand is shifted right arithmetically by one bit position as shown in Figure 1-40. The MSB retains its value (sign). RRA operates equal to a signed division by 2. The MSB is retained and shifted into the MSB-1. The LSB+1 is shifted into the LSB. The previous LSB is shifted into the carry bit C.
Status Bits
 N: Set if result is negative (MSB = 1), reset otherwise (MSB = 0)
 Z: Set if result is zero, reset otherwise
 C: Loaded from the LSB
 V: Reset
Mode Bits OSCOFF, CPUOFF, and GIE are not affected.
Example The signed 16-bit number in R5 is shifted arithmetically right one position.

```
RRA R5 ; R5/2 -> R5
```

Example The signed RAM byte EDE is shifted arithmetically right one position.

```
RRA.B EDE ; EDE/2 -> EDE
```

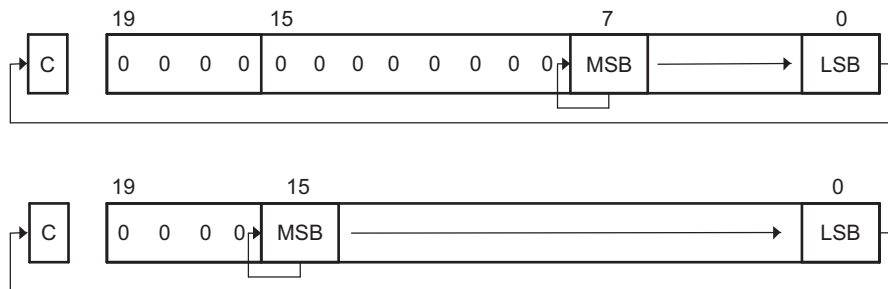


Figure 1-40. Rotate Right Arithmetically RRA.B and RRA.W

1.6.2.41 RRC

RRC[.W] Rotate right through carry destination word**RRC.B** Rotate right through carry destination byte

Syntax RRC dst OF RRC.W dst
RRC.B dst

Operation $C \rightarrow \text{MSB} \rightarrow \text{MSB}-1 \rightarrow \dots \text{LSB}+1 \rightarrow \text{LSB} \rightarrow C$ **Description** The destination operand is shifted right by one bit position as shown in Figure 1-41. The carry bit C is shifted into the MSB and the LSB is shifted into the carry bit C.**Status Bits** N: Set if result is negative (MSB = 1), reset otherwise (MSB = 0)

Z: Set if result is zero, reset otherwise

C: Loaded from the LSB

V: Reset

Mode Bits OSCOFF, CPUOFF, and GIE are not affected.**Example** RAM word EDE is shifted right one bit position. The MSB is loaded with 1.

```
SETC          ; Prepare carry for MSB
RRC  EDE      ; EDE = EDE >> 1 + 8000h
```

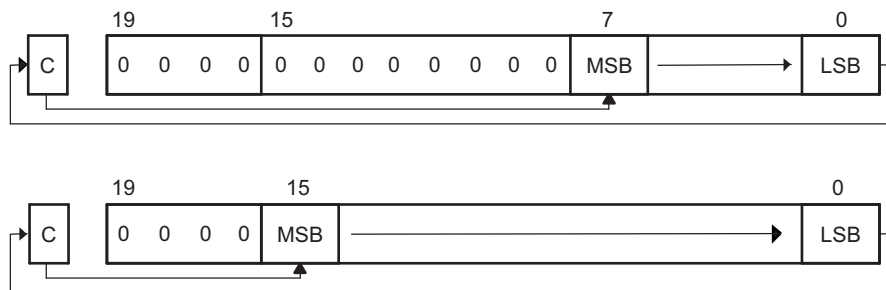


Figure 1-41. Rotate Right Through Carry RRC.B and RRC.W

1.6.2.42 SBC

* SBC[W]	Subtract borrow (.NOT. carry) from destination
* SBC.B	Subtract borrow (.NOT. carry) from destination
Syntax	SBC dst OR SBC.W dst SBC.B dst
Operation	dst + 0FFFFh + C → dst dst + 0FFh + C → dst
Emulation	SUBC #0, dst SUBC.B #0, dst
Description	The carry bit (C) is added to the destination operand minus one. The previous contents of the destination are lost.
Status Bits	N: Set if result is negative, reset if positive Z: Set if result is zero, reset otherwise C: Set if there is a carry from the MSB of the result, reset otherwise Set to 1 if no borrow, reset if borrow V: Set if an arithmetic overflow occurs, reset otherwise
Mode Bits	OSCOFF, CPUOFF, and GIE are not affected.
Example	The 16-bit counter pointed to by R13 is subtracted from a 32-bit counter pointed to by R12.

```

SUB    @R13,0(R12)    ; Subtract LSDs
SBC    2(R12)         ; Subtract carry from MSD
    
```

Example The 8-bit counter pointed to by R13 is subtracted from a 16-bit counter pointed to by R12.

```

SUB.B  @R13,0(R12)    ; Subtract LSDs
SBC.B  1(R12)         ; Subtract carry from MSD
    
```

NOTE: Borrow implementation

The borrow is treated as a .NOT. carry:

Borrow	Carry Bit
Yes	0
No	1

1.6.2.43 SETC

* SETC	Set carry bit
Syntax	SETC
Operation	1 → C
Emulation	BIS #1,SR
Description	The carry bit (C) is set.
Status Bits	N: Not affected Z: Not affected C: Set V: Not affected
Mode Bits	OSCOFF, CPUOFF, and GIE are not affected.
Example	Emulation of the decimal subtraction: Subtract R5 from R6 decimally. Assume that R5 = 03987h and R6 = 04137h.

```

DSUB  ADD    #06666h,R5    ; Move content R5 from 0-9 to 6-0Fh
                                ; R5 = 03987h + 06666h = 09FEDh
      INV    R5              ; Invert this (result back to 0-9)
                                ; R5 = .NOT. R5 = 06012h
      SETC                                ; Prepare carry = 1
      DADD   R5,R6          ; Emulate subtraction by addition of:
                                ; (010000h - R5 - 1)
                                ; R6 = R6 + R5 + 1
                                ; R6 = 0150h

```

1.6.2.44 SETN

* SETN	Set negative bit
Syntax	SETN
Operation	1 → N
Emulation	BIS #4, SR
Description	The negative bit (N) is set.
Status Bits	N: Set Z: Not affected C: Not affected V: Not affected
Mode Bits	OSCOFF, CPUOFF, and GIE are not affected.

1.6.2.45 SETZ

* SETZ	Set zero bit
Syntax	SETZ
Operation	1 → N
Emulation	BIS #2, SR
Description	The zero bit (Z) is set.
Status Bits	N: Not affected Z: Set C: Not affected V: Not affected
Mode Bits	OSCOFF, CPUOFF, and GIE are not affected.

1.6.2.46 SUB

SUB[W]	Subtract source word from destination word
SUB.B	Subtract source byte from destination byte
Syntax	SUB src,dst OR SUB.W src,dst SUB.B src,dst
Operation	(.not.src) + 1 + dst → dst or dst – src → dst
Description	The source operand is subtracted from the destination operand. This is made by adding the 1s complement of the source + 1 to the destination. The source operand is not affected, the result is written to the destination operand.
Status Bits	N: Set if result is negative (src > dst), reset if positive (src ≤ dst) Z: Set if result is zero (src = dst), reset otherwise (src ≠ dst) C: Set if there is a carry from the MSB, reset otherwise V: Set if the subtraction of a negative source operand from a positive destination operand delivers a negative result, or if the subtraction of a positive source operand from a negative destination operand delivers a positive result, reset otherwise (no overflow)
Mode Bits	OSCOFF, CPUOFF, and GIE are not affected.
Example	A 16-bit constant 7654h is subtracted from RAM word EDE.
	<pre>SUB #7654h,&EDE ; Subtract 7654h from EDE</pre>
Example	A table word pointed to by R5 (20-bit address) is subtracted from R7. Afterwards, if R7 contains zero, jump to label TONI. R5 is then auto-incremented by 2. R7.19:16 = 0.
	<pre>SUB @R5+,R7 ; Subtract table number from R7. R5 + 2 JZ TONI ; R7 = @R5 (before subtraction) ... ; R7 <> @R5 (before subtraction)</pre>
Example	Byte CNT is subtracted from byte R12 points to. The address of CNT is within PC ± 32K. The address R12 points to is in full memory range.
	<pre>SUB.B CNT,0(R12) ; Subtract CNT from @R12</pre>

1.6.2.47 SUBC

SUBC[W]	Subtract source word with carry from destination word
SUBC.B	Subtract source byte with carry from destination byte
Syntax	SUBC src,dst OR SUBC.W src,dst SUBC.B src,dst
Operation	$(.not.src) + C + dst \rightarrow dst$ or $dst - (src - 1) + C \rightarrow dst$
Description	The source operand is subtracted from the destination operand. This is done by adding the 1s complement of the source + carry to the destination. The source operand is not affected, the result is written to the destination operand. Used for 32, 48, and 64-bit operands.
Status Bits	N: Set if result is negative (MSB = 1), reset if positive (MSB = 0) Z: Set if result is zero, reset otherwise C: Set if there is a carry from the MSB, reset otherwise V: Set if the subtraction of a negative source operand from a positive destination operand delivers a negative result, or if the subtraction of a positive source operand from a negative destination operand delivers a positive result, reset otherwise (no overflow)
Mode Bits	OSCOFF, CPUOFF, and GIE are not affected.
Example	A 16-bit constant 7654h is subtracted from R5 with the carry from the previous instruction. R5.19:16 = 0

```
SUBC.W #7654h,R5 ; Subtract 7654h + C from R5
```

Example A 48-bit number (3 words) pointed to by R5 (20-bit address) is subtracted from a 48-bit counter in RAM, pointed to by R7. R5 points to the next 48-bit number afterwards. The address R7 points to is in full memory range.

```
SUB @R5+,0(R7) ; Subtract LSBs. R5 + 2
SUBC @R5+,2(R7) ; Subtract MIDs with C. R5 + 2
SUBC @R5+,4(R7) ; Subtract MSBs with C. R5 + 2
```

Example Byte CNT is subtracted from the byte, R12 points to. The carry of the previous instruction is used. The address of CNT is in lower 64 K.

```
SUBC.B &CNT,0(R12) ; Subtract byte CNT from @R12
```

1.6.2.48 SWPB

SWPB Swap bytes
Syntax SWPB *dst*
Operation *dst.15:8* ↔ *dst.7:0*
Description The high and the low byte of the operand are exchanged. PC.19:16 bits are cleared in register mode.
Status Bits Status bits are not affected
Mode Bits OSCOFF, CPUOFF, and GIE are not affected.
Example Exchange the bytes of RAM word EDE (lower 64 K)

```
MOV    #1234h,&EDE    ; 1234h -> EDE
SWPB  &EDE           ; 3412h -> EDE
```

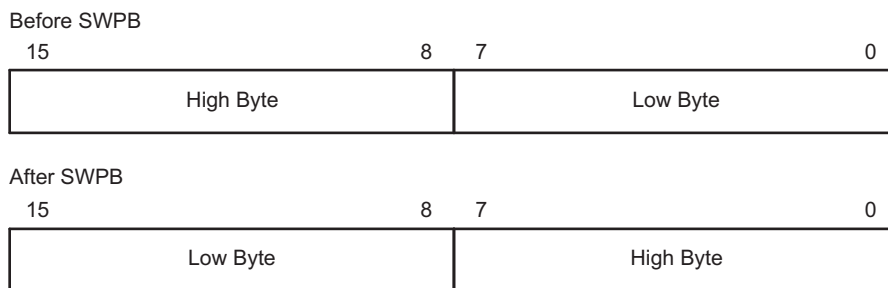


Figure 1-42. Swap Bytes in Memory

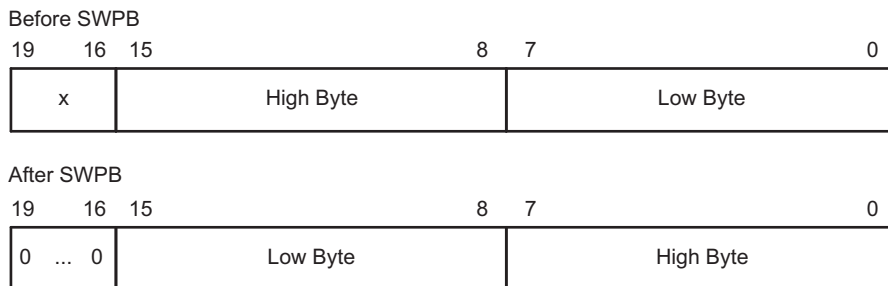


Figure 1-43. Swap Bytes in a Register

1.6.2.49 SXT

SXT	Extend sign
Syntax	SXT dst
Operation	dst.7 → dst.15:8, dst.7 → dst.19:8 (register mode)
Description	Register mode: the sign of the low byte of the operand is extended into the bits Rdst.19:8. Rdst.7 = 0: Rdst.19:8 = 000h afterwards Rdst.7 = 1: Rdst.19:8 = FFFh afterwards Other modes: the sign of the low byte of the operand is extended into the high byte. dst.7 = 0: high byte = 00h afterwards dst.7 = 1: high byte = FFh afterwards
Status Bits	N: Set if result is negative, reset otherwise Z: Set if result is zero, reset otherwise C: Set if result is not zero, reset otherwise (C = .not.Z) V: Reset
Mode Bits	OSCOFF, CPUOFF, and GIE are not affected.
Example	The signed 8-bit data in EDE (lower 64 K) is sign extended and added to the 16-bit signed data in R7.

```
MOV.B  &EDE,R5      ; EDE -> R5. 00XXh
SXT    R5           ; Sign extend low byte to R5.19:8
ADD    R5,R7       ; Add signed 16-bit values
```

Example The signed 8-bit data in EDE (PC +32 K) is sign extended and added to the 20-bit data in R7.

```
MOV.B  EDE,R5      ; EDE -> R5. 00XXh
SXT    R5           ; Sign extend low byte to R5.19:8
ADDA   R5,R7       ; Add signed 20-bit values
```


1.6.2.50 TST

*** TST.W]** Test destination

*** TST.B** Test destination

Syntax TST dst OR TST.W dst
TST.B dst

Operation dst + 0FFFFh + 1

dst + 0FFh + 1

Emulation CMP #0, dst
CMP.B #0, dst

Description The destination operand is compared with zero. The status bits are set according to the result. The destination is not affected.

Status Bits
N: Set if destination is negative, reset if positive
Z: Set if destination contains zero, reset otherwise
C: Set
V: Reset

Mode Bits OSCOFF, CPUOFF, and GIE are not affected.

Example R7 is tested. If it is negative, continue at R7NEG; if it is positive but not zero, continue at R7POS.

```

TST    R7        ; Test R7
JN     R7NEG     ; R7 is negative
JZ     R7ZERO    ; R7 is zero
R7POS  .....    ; R7 is positive but not zero
R7NEG  .....    ; R7 is negative
R7ZERO .....    ; R7 is zero
    
```

Example The low byte of R7 is tested. If it is negative, continue at R7NEG; if it is positive but not zero, continue at R7POS.

```

TST.B  R7        ; Test low byte of R7
JN     R7NEG     ; Low byte of R7 is negative
JZ     R7ZERO    ; Low byte of R7 is zero
R7POS  .....    ; Low byte of R7 is positive but not zero
R7NEG  .....    ; Low byte of R7 is negative
R7ZERO .....    ; Low byte of R7 is zero
    
```

1.6.2.51 XOR

XOR[.W]	Exclusive OR source word with destination word
XOR.B	Exclusive OR source byte with destination byte
Syntax	XOR src,dst OR XOR.W src,dst XOR.B src,dst
Operation	src .xor. dst → dst
Description	The source and destination operands are exclusively ORed. The result is placed into the destination. The source operand is not affected. The previous content of the destination is lost.
Status Bits	N: Set if result is negative (MSB = 1), reset if positive (MSB = 0) Z: Set if result is zero, reset otherwise C: Set if result is not zero, reset otherwise (C = .not. Z) V: Set if both operands are negative before execution, reset otherwise
Mode Bits	OSCOFF, CPUOFF, and GIE are not affected.
Example	Toggle bits in word CNTR (16-bit data) with information (bit = 1) in address-word TONI. Both operands are located in lower 64 K.
	<pre>XOR &TONI,&CNTR ; Toggle bits in CNTR</pre>
Example	A table word pointed to by R5 (20-bit address) is used to toggle bits in R6. R6.19:16 = 0.
	<pre>XOR @R5,R6 ; Toggle bits in R6</pre>
Example	Reset to zero those bits in the low byte of R7 that are different from the bits in byte EDE. R7.19:8 = 0. The address of EDE is within PC ± 32 K.
	<pre>XOR.B EDE,R7 ; Set different bits to 1 in R7. INV.B R7 ; Invert low byte of R7, high byte is 0h</pre>

1.6.3 Extended Instructions

The extended MSP430X instructions give the MSP430X CPU full access to its 20-bit address space. MSP430X instructions require an additional word of op-code called the extension word. All addresses, indexes, and immediate numbers have 20-bit values when preceded by the extension word. The MSP430X extended instructions are listed and described in the following pages.

1.6.3.2 ADDX

ADDX.A Add source address-word to destination address-word

ADDX.[W] Add source word to destination word

ADDX.B Add source byte to destination byte

Syntax `ADDX.A src,dst`

`ADDX src,dst` Or `ADDX.W src,dst`

`ADDX.B src,dst`

Operation `src + dst → dst`

Description The source operand is added to the destination operand. The previous contents of the destination are lost. Both operands can be located in the full address space.

Status Bits N: Set if result is negative (MSB = 1), reset if positive (MSB = 0)

Z: Set if result is zero, reset otherwise

C: Set if there is a carry from the MSB of the result, reset otherwise

V: Set if the result of two positive operands is negative, or if the result of two negative numbers is positive, reset otherwise

Mode Bits OSCOFF, CPUOFF, and GIE are not affected.

Example Ten is added to the 20-bit pointer CNTR located in two words CNTR (LSBs) and CNTR+2 (MSBs).

```
ADDX.A    #10,CNTR    ; Add 10 to 20-bit pointer
```

Example A table word (16-bit) pointed to by R5 (20-bit address) is added to R6. The jump to label TONI is performed on a carry.

```
ADDX.W    @R5,R6     ; Add table word to R6
JC        TONI       ; Jump if carry
...       ; No carry
```

Example A table byte pointed to by R5 (20-bit address) is added to R6. The jump to label TONI is performed if no carry occurs. The table pointer is auto-incremented by 1.

```
ADDX.B    @R5+,R6    ; Add table byte to R6. R5 + 1. R6: 000xxh
JNC       TONI       ; Jump if no carry
...       ; Carry occurred
```

Note: Use ADDA for the following two cases for better code density and execution.

```
ADDX.A    Rsrc,Rdst
ADDX.A    #imm20,Rdst
```

1.6.3.3 ADDCX**ADDCX.A** Add source address-word and carry to destination address-word**ADDCX.[W]** Add source word and carry to destination word**ADDCX.B** Add source byte and carry to destination byte**Syntax** `ADDCX.A src,dst``ADDCX src,dst Or ADDCX.W src,dst``ADDCX.B src,dst`**Operation** `src + dst + C → dst`**Description** The source operand and the carry bit C are added to the destination operand. The previous contents of the destination are lost. Both operands may be located in the full address space.**Status Bits** N: Set if result is negative (MSB = 1), reset if positive (MSB = 0)

Z: Set if result is zero, reset otherwise

C: Set if there is a carry from the MSB of the result, reset otherwise

V: Set if the result of two positive operands is negative, or if the result of two negative numbers is positive, reset otherwise

Mode Bits OSCOFF, CPUOFF, and GIE are not affected.**Example** Constant 15 and the carry of the previous instruction are added to the 20-bit counter CNTR located in two words.

```
ADDCX.A #15,&CNTR ; Add 15 + C to 20-bit CNTR
```

Example A table word pointed to by R5 (20-bit address) and the carry C are added to R6. The jump to label TONI is performed on a carry.

```
ADDCX.W @R5,R6 ; Add table word + C to R6
JC TONI ; Jump if carry
... ; No carry
```

Example A table byte pointed to by R5 (20-bit address) and the carry bit C are added to R6. The jump to label TONI is performed if no carry occurs. The table pointer is auto-incremented by 1.

```
ADDCX.B @R5+,R6 ; Add table byte + C to R6. R5 + 1
JNC TONI ; Jump if no carry
... ; Carry occurred
```

1.6.3.4 ANDX

ANDX.A Logical AND of source address-word with destination address-word

ANDX.[W] Logical AND of source word with destination word

ANDX.B Logical AND of source byte with destination byte

Syntax
 ANDX.A *src,dst*
 ANDX *src,dst* Or ANDX.W *src,dst*
 ANDX.B *src,dst*

Operation *src .and. dst* → *dst*

Description The source operand and the destination operand are logically ANDed. The result is placed into the destination. The source operand is not affected. Both operands may be located in the full address space.

Status Bits
 N: Set if result is negative (MSB = 1), reset if positive (MSB = 0)
 Z: Set if result is zero, reset otherwise
 C: Set if the result is not zero, reset otherwise. C = (.not. Z)
 V: Reset

Mode Bits OSCOFF, CPUOFF, and GIE are not affected.

Example The bits set in R5 (20-bit data) are used as a mask (AAA55h) for the address-word TOM located in two words. If the result is zero, a branch is taken to label TONI.

```
MOVA    #AAA55h,R5      ; Load 20-bit mask to R5
ANDX.A  R5,TOM          ; TOM .and. R5 -> TOM
JZ      TONI            ; Jump if result 0
...     ; Result > 0
```

or shorter:

```
ANDX.A  #AAA55h,TOM     ; TOM .and. AAA55h -> TOM
JZ      TONI            ; Jump if result 0
```

Example A table byte pointed to by R5 (20-bit address) is logically ANDed with R6. R6.19:8 = 0. The table pointer is auto-incremented by 1.

```
ANDX.B  @R5+,R6        ; AND table byte with R6. R5 + 1
```

1.6.3.5 BICX

BICX.A Clear bits set in source address-word in destination address-word

BICX.[W] Clear bits set in source word in destination word

BICX.B Clear bits set in source byte in destination byte

Syntax
 BICX.A *src,dst*
 BICX *src,dst* OR BICX.W *src,dst*
 BICX.B *src,dst*

Operation (.not. *src*) .and. *dst* → *dst*

Description The inverted source operand and the destination operand are logically ANDed. The result is placed into the destination. The source operand is not affected. Both operands may be located in the full address space.

Status Bits
 N: Not affected
 Z: Not affected
 C: Not affected
 V: Not affected

Mode Bits OSCOFF, CPUOFF, and GIE are not affected.

Example The bits 19:15 of R5 (20-bit data) are cleared.

```
BICX.A #0F8000h,R5 ; Clear R5.19:15 bits
```

Example A table word pointed to by R5 (20-bit address) is used to clear bits in R7. R7.19:16 = 0.

```
BICX.W @R5,R7 ; Clear bits in R7
```

Example A table byte pointed to by R5 (20-bit address) is used to clear bits in output Port1.

```
BICX.B @R5,&P1OUT ; Clear I/O port P1 bits
```


1.6.3.6 BISX

BISX.A Set bits set in source address-word in destination address-word

BISX.[W] Set bits set in source word in destination word

BISX.B Set bits set in source byte in destination byte

Syntax
 BISX.A *src,dst*
 BISX *src,dst* OR BISX.W *src,dst*
 BISX.B *src,dst*

Operation *src .or. dst* → *dst*

Description The source operand and the destination operand are logically ORed. The result is placed into the destination. The source operand is not affected. Both operands may be located in the full address space.

Status Bits
 N: Not affected
 Z: Not affected
 C: Not affected
 V: Not affected

Mode Bits OSCOFF, CPUOFF, and GIE are not affected.

Example Bits 16 and 15 of R5 (20-bit data) are set to one.

```
BISX.A    #018000h,R5    ; Set R5.16:15 bits
```

Example A table word pointed to by R5 (20-bit address) is used to set bits in R7.

```
BISX.W    @R5,R7        ; Set bits in R7
```

Example A table byte pointed to by R5 (20-bit address) is used to set bits in output Port1.

```
BISX.B    @R5,&P1OUT    ; Set I/O port P1 bits
```

1.6.3.7 BITX**BITX.A** Test bits set in source address-word in destination address-word**BITX.[W]** Test bits set in source word in destination word**BITX.B** Test bits set in source byte in destination byte**Syntax** BITX.A *src,dst*BITX *src,dst* OR BITX.W *src,dst*BITX.B *src,dst***Operation** *src* .and. *dst* → *dst***Description** The source operand and the destination operand are logically ANDed. The result affects only the status bits. Both operands may be located in the full address space.**Status Bits** N: Set if result is negative (MSB = 1), reset if positive (MSB = 0)

Z: Set if result is zero, reset otherwise

C: Set if the result is not zero, reset otherwise. C = (.not. Z)

V: Reset

Mode Bits OSCOFF, CPUOFF, and GIE are not affected.**Example** Test if bit 16 or 15 of R5 (20-bit data) is set. Jump to label TONI if so.

```

BITX.A    #018000h,R5      ; Test R5.16:15 bits
JNZ      TONI             ; At least one bit is set
...      ; Both are reset

```

Example A table word pointed to by R5 (20-bit address) is used to test bits in R7. Jump to label TONI if at least one bit is set.

```

BITX.W    @R5,R7          ; Test bits in R7: C = .not.Z
JC       TONI             ; At least one is set
...      ; Both are reset

```

Example A table byte pointed to by R5 (20-bit address) is used to test bits in input Port1. Jump to label TONI if no bit is set. The next table byte is addressed.

```

BITX.B    @R5+,&P1IN      ; Test input P1 bits. R5 + 1
JNC      TONI             ; No corresponding input bit is set
...      ; At least one bit is set

```

1.6.3.8 CLRX

* **CLRX.A** Clear destination address-word

* **CLRX.[W]** Clear destination word

* **CLRX.B** Clear destination byte

Syntax CLRX.A dst
 CLRX dst or CLRX.W dst
 CLRX.B dst

Operation 0 → dst

Emulation MOVX.A #0, dst

MOVX #0, dst

MOVX.B #0, dst

Description The destination operand is cleared.

Status Bits Status bits are not affected.

Mode Bits OSCOFF, CPUOFF, and GIE are not affected.

Example RAM address-word TONI is cleared.

```
CLRX.A TONI ; 0 -> TONI
```

1.6.3.9 CMPX

CMPX.A	Compare source address-word and destination address-word
CMPX.[W]	Compare source word and destination word
CMPX.B	Compare source byte and destination byte
Syntax	CMPX.A src,dst CMPX src,dst OR CMPX.W src,dst CMPX.B src,dst
Operation	(.not. src) + 1 + dst or dst – src
Description	The source operand is subtracted from the destination operand by adding the 1s complement of the source + 1 to the destination. The result affects only the status bits. Both operands may be located in the full address space.
Status Bits	N: Set if result is negative (src > dst), reset if positive (src ≤ dst) Z: Set if result is zero (src = dst), reset otherwise (src ≠ dst) C: Set if there is a carry from the MSB, reset otherwise V: Set if the subtraction of a negative source operand from a positive destination operand delivers a negative result, or if the subtraction of a positive source operand from a negative destination operand delivers a positive result, reset otherwise (no overflow)
Mode Bits	OSCOFF, CPUOFF, and GIE are not affected.
Example	Compare EDE with a 20-bit constant 18000h. Jump to label TONI if EDE equals the constant.

```
CMPX.A  #018000h,EDE      ; Compare EDE with 18000h
JEQ     TONI              ; EDE contains 18000h
...     ; Not equal
```

Example A table word pointed to by R5 (20-bit address) is compared with R7. Jump to label TONI if R7 contains a lower, signed, 16-bit number.

```
CMPX.W  @R5,R7           ; Compare two signed numbers
JL      TONI              ; R7 < @R5
...     ; R7 >= @R5
```

Example A table byte pointed to by R5 (20-bit address) is compared to the input in I/O Port1. Jump to label TONI if the values are equal. The next table byte is addressed.

```
CMPX.B  @R5+,&P1IN       ; Compare P1 bits with table. R5 + 1
JEQ     TONI              ; Equal contents
...     ; Not equal
```

Note: Use CMPA for the following two cases for better density and execution.

```
CMPA    Rsrc,Rdst
CMPA    #imm20,Rdst
```


1.6.3.11 DADDX

DADDX.A	Add source address-word and carry decimally to destination address-word
DADDX.[W]	Add source word and carry decimally to destination word
DADDX.B	Add source byte and carry decimally to destination byte
Syntax	DADDX.A src,dst DADDX src,dst Or DADDX.W src,dst DADDX.B src,dst
Operation	src + dst + C → dst (decimally)
Description	The source operand and the destination operand are treated as two (.B), four (.W), or five (.A) binary coded decimals (BCD) with positive signs. The source operand and the carry bit C are added decimally to the destination operand. The source operand is not affected. The previous contents of the destination are lost. The result is not defined for non-BCD numbers. Both operands may be located in the full address space.
Status Bits	N: Set if MSB of result is 1 (address-word > 79999h, word > 7999h, byte > 79h), reset if MSB is 0. Z: Set if result is zero, reset otherwise C: Set if the BCD result is too large (address-word > 99999h, word > 9999h, byte > 99h), reset otherwise V: Undefined
Mode Bits	OSCOFF, CPUOFF, and GIE are not affected.
Example	Decimal 10 is added to the 20-bit BCD counter DECCNTR located in two words.

```
DADDX.A #10h,&DECCNTR ; Add 10 to 20-bit BCD counter
```

Example The eight-digit BCD number contained in 20-bit addresses BCD and BCD+2 is added decimally to an eight-digit BCD number contained in R4 and R5 (BCD+2 and R5 contain the MSDs).

```
CLRC ; Clear carry
DADDX.W BCD,R4 ; Add LSDs
DADDX.W BCD+2,R5 ; Add MSDs with carry
JC OVERFLOW ; Result >99999999: go to error routine
... ; Result ok
```

Example The two-digit BCD number contained in 20-bit address BCD is added decimally to a two-digit BCD number contained in R4.

```
CLRC ; Clear carry
DADDX.B BCD,R4 ; Add BCD to R4 decimally.
; R4: 000ddh
```

1.6.3.12 DECX

* DECX.A	Decrement destination address-word
* DECX.[W]	Decrement destination word
* DECX.B	Decrement destination byte
Syntax	DECX.A dst DECX dst or DECX.W dst DECX.B dst
Operation	$dst - 1 \rightarrow dst$
Emulation	SUBX.A #1, dst SUBX #1, dst SUBX.B #1, dst
Description	The destination operand is decremented by one. The original contents are lost.
Status Bits	N: Set if result is negative, reset if positive Z: Set if dst contained 1, reset otherwise C: Reset if dst contained 0, set otherwise V: Set if an arithmetic overflow occurs, otherwise reset
Mode Bits	OSCOFF, CPUOFF, and GIE are not affected.
Example	RAM address-word TONI is decremented by one.

```
DECX.A TONI ; Decrement TONI
```

1.6.3.13 DECDX

* DECDX.A	Double-decrement destination address-word
* DECDX.[W]	Double-decrement destination word
* DECDX.B	Double-decrement destination byte
Syntax	DECDX.A dst DECDX dst or DECDX.W dst DECDX.B dst
Operation	dst – 2 → dst
Emulation	SUBX.A #2, dst SUBX #2, dst SUBX.B #2, dst
Description	The destination operand is decremented by two. The original contents are lost.
Status Bits	N: Set if result is negative, reset if positive Z: Set if dst contained 2, reset otherwise C: Reset if dst contained 0 or 1, set otherwise V: Set if an arithmetic overflow occurs, otherwise reset
Mode Bits	OSCOFF, CPUOFF, and GIE are not affected.
Example	RAM address-word TONI is decremented by two.

```
DECDX.A TONI ; Decrement TONI
```


1.6.3.14 INCX

* INCX.A	Increment destination address-word
* INCX.[W]	Increment destination word
* INCX.B	Increment destination byte
Syntax	INCX.A dst INCX dst or INCX.W dst INCX.B dst
Operation	dst + 1 → dst
Emulation	ADDX.A #1, dst ADDX #1, dst ADDX.B #1, dst
Description	The destination operand is incremented by one. The original contents are lost.
Status Bits	N: Set if result is negative, reset if positive Z: Set if dst contained 0FFFFFFh, reset otherwise Set if dst contained 0FFFFh, reset otherwise Set if dst contained 0FFh, reset otherwise C: Set if dst contained 0FFFFFFh, reset otherwise Set if dst contained 0FFFFh, reset otherwise Set if dst contained 0FFh, reset otherwise V: Set if dst contained 07FFFh, reset otherwise Set if dst contained 07FFFh, reset otherwise Set if dst contained 07Fh, reset otherwise
Mode Bits	OSCOFF, CPUOFF, and GIE are not affected.
Example	RAM address-wordTONI is incremented by one.
	INCX.A TONI ; Increment TONI (20-bits)

1.6.3.15 INCDX

* INCDX.A	Double-increment destination address-word
* INCDX.[W]	Double-increment destination word
* INCDX.B	Double-increment destination byte
Syntax	INCDX.A dst INCDX dst or INCDX.W dst INCDX.B dst
Operation	dst + 2 → dst
Emulation	ADDX.A #2, dst ADDX #2, dst ADDX.B #2, dst
Description	The destination operand is incremented by two. The original contents are lost.
Status Bits	N: Set if result is negative, reset if positive Z: Set if dst contained 0FFFFEh, reset otherwise Set if dst contained 0FFFEh, reset otherwise Set if dst contained 0FEh, reset otherwise C: Set if dst contained 0FFFFEh or 0FFFFFFh, reset otherwise Set if dst contained 0FFFEh or 0FFFFh, reset otherwise Set if dst contained 0FEh or 0FFh, reset otherwise V: Set if dst contained 07FFFEh or 07FFFFh, reset otherwise Set if dst contained 07FFEh or 07FFFh, reset otherwise Set if dst contained 07Eh or 07Fh, reset otherwise
Mode Bits	OSCOFF, CPUOFF, and GIE are not affected.
Example	RAM byte LEO is incremented by two; PC points to upper memory.

```
INCDX.B LEO ; Increment LEO by two
```

1.6.3.16 INVX

* INVX.A	Invert destination
* INVX.[W]	Invert destination
* INVX.B	Invert destination
Syntax	INVX.A dst INVX dst or INVX.W dst INVX.B dst
Operation	.NOT.dst → dst
Emulation	XORX.A #FFFFFFh, dst XORX #FFFFFFh, dst XORX.B #0FFh, dst
Description	The destination operand is inverted. The original contents are lost.
Status Bits	N: Set if result is negative, reset if positive Z: Set if dst contained 0FFFFFFh, reset otherwise Set if dst contained 0FFFFh, reset otherwise Set if dst contained 0FFh, reset otherwise C: Set if result is not zero, reset otherwise (= .NOT. Zero) V: Set if initial destination operand was negative, otherwise reset
Mode Bits	OSCOFF, CPUOFF, and GIE are not affected.
Example	20-bit content of R5 is negated (twos complement).

```

INVX.A  R5      ; Invert R5
INCX.A  R5      ; R5 is now negated
    
```

Example Content of memory byte LEO is negated. PC is pointing to upper memory.

```

INVX.B  LEO     ; Invert LEO
INCX.B  LEO     ; MEM(LEO) is negated
    
```

1.6.3.17 MOVX

MOVX.A	Move source address-word to destination address-word
MOVX.[W]	Move source word to destination word
MOVX.B	Move source byte to destination byte
Syntax	MOVX.A src,dst MOVX src,dst OR MOVX.W src,dst MOVX.B src,dst
Operation	src → dst
Description	The source operand is copied to the destination. The source operand is not affected. Both operands may be located in the full address space.
Status Bits	N: Not affected Z: Not affected C: Not affected V: Not affected
Mode Bits	OSCOFF, CPUOFF, and GIE are not affected.
Example	Move a 20-bit constant 18000h to absolute address-word EDE

```
MOVX.A    #018000h,&EDE          ; Move 18000h to EDE
```

Example The contents of table EDE (word data, 20-bit addresses) are copied to table TOM. The length of the table is 030h words.

```

Loop      MOVA    #EDE,R10          ; Prepare pointer (20-bit address)
          MOVX.W  @R10+,TOM-EDE-2(R10) ; R10 points to both tables.
          ; R10+2
          CMPA    #EDE+60h,R10      ; End of table reached?
          JLO    Loop              ; Not yet
          ...                      ; Copy completed
```

Example The contents of table EDE (byte data, 20-bit addresses) are copied to table TOM. The length of the table is 020h bytes.

```

Loop      MOVA    #EDE,R10          ; Prepare pointer (20-bit)
          MOV     #20h,R9           ; Prepare counter
          MOVX.W  @R10+,TOM-EDE-2(R10) ; R10 points to both tables.
          ; R10+1
          DEC     R9                ; Decrement counter
          JNZ    Loop              ; Not yet done
          ...                      ; Copy completed
```

Ten of the 28 possible addressing combinations of the MOVX.A instruction can use the MOVA instruction. This saves two bytes and code cycles. Examples for the addressing combinations are:

MOVX.A	Rsrc,Rdst	MOVA	Rsrc,Rdst	; Reg/Reg
MOVX.A	#imm20,Rdst	MOVA	#imm20,Rdst	; Immediate/Reg
MOVX.A	&abs20,Rdst	MOVA	&abs20,Rdst	; Absolute/Reg
MOVX.A	@Rsrc,Rdst	MOVA	@Rsrc,Rdst	; Indirect/Reg
MOVX.A	@Rsrc+,Rdst	MOVA	@Rsrc+,Rdst	; Indirect,Auto/Reg
MOVX.A	Rsrc,&abs20	MOVA	Rsrc,&abs20	; Reg/Absolute

The next four replacements are possible only if 16-bit indexes are sufficient for the addressing:

MOVX.A	z20(Rsrc),Rdst	MOVA	z16(Rsrc),Rdst	; Indexed/Reg
MOVX.A	Rsrc,z20(Rdst)	MOVA	Rsrc,z16(Rdst)	; Reg/Indexed
MOVX.A	symb20,Rdst	MOVA	symb16,Rdst	; Symbolic/Reg
MOVX.A	Rsrc,symb20	MOVA	Rsrc,symb16	; Reg/Symbolic

1.6.3.18 POPM

POPM.A	Restore n CPU registers (20-bit data) from the stack
POPM.[W]	Restore n CPU registers (16-bit data) from the stack
Syntax	<pre>POPM.A #n,Rdst 1 ≤ n ≤ 16 POPM.W #n,Rdst OR POPM #n,Rdst 1 ≤ n ≤ 16</pre>
Operation	<p>POPM.A: Restore the register values from stack to the specified CPU registers. The SP is incremented by four for each register restored from stack. The 20-bit values from stack (two words per register) are restored to the registers.</p> <p>POPM.W: Restore the 16-bit register values from stack to the specified CPU registers. The SP is incremented by two for each register restored from stack. The 16-bit values from stack (one word per register) are restored to the CPU registers.</p> <p>Note : This instruction does not use the extension word.</p>
Description	<p>POPM.A: The CPU registers pushed on the stack are moved to the extended CPU registers, starting with the CPU register (Rdst – n + 1). The SP is incremented by (n × 4) after the operation.</p> <p>POPM.W: The 16-bit registers pushed on the stack are moved back to the CPU registers, starting with CPU register (Rdst – n + 1). The SP is incremented by (n × 2) after the instruction. The MSBs (Rdst.19:16) of the restored CPU registers are cleared.</p>
Status Bits	Status bits are not affected, except SR is included in the operation.
Mode Bits	OSCOFF, CPUOFF, and GIE are not affected.
Example	Restore the 20-bit registers R9, R10, R11, R12, R13 from the stack
	<pre>POPM.A #5,R13 ; Restore R9, R10, R11, R12, R13</pre>
Example	Restore the 16-bit registers R9, R10, R11, R12, R13 from the stack.
	<pre>POPM.W #5,R13 ; Restore R9, R10, R11, R12, R13</pre>

1.6.3.19 PUSHM

PUSHM.A	Save n CPU registers (20-bit data) on the stack
PUSHM.[W]	Save n CPU registers (16-bit words) on the stack
Syntax	PUSHM.A #n,Rdst $1 \leq n \leq 16$ PUSHM.W #n,Rdst OR PUSHM #n,Rdst $1 \leq n \leq 16$
Operation	PUSHM.A: Save the 20-bit CPU register values on the stack. The SP is decremented by four for each register stored on the stack. The MSBs are stored first (higher address). PUSHM.W: Save the 16-bit CPU register values on the stack. The SP is decremented by two for each register stored on the stack.
Description	PUSHM.A: The n CPU registers, starting with Rdst backwards, are stored on the stack. The SP is decremented by (n × 4) after the operation. The data (Rn.19:0) of the pushed CPU registers is not affected. PUSHM.W: The n registers, starting with Rdst backwards, are stored on the stack. The SP is decremented by (n × 2) after the operation. The data (Rn.19:0) of the pushed CPU registers is not affected. Note : This instruction does not use the extension word.
Status Bits	Status bits are not affected.
Mode Bits	OSCOFF, CPUOFF, and GIE are not affected.
Example	Save the five 20-bit registers R9, R10, R11, R12, R13 on the stack
	PUSHM.A #5,R13 ; Save R13, R12, R11, R10, R9
Example	Save the five 16-bit registers R9, R10, R11, R12, R13 on the stack
	PUSHM.W #5,R13 ; Save R13, R12, R11, R10, R9

1.6.3.20 POPX

* **POPX.A** Restore single address-word from the stack

* **POPX.[W]** Restore single word from the stack

* **POPX.B** Restore single byte from the stack

Syntax POPX.A dst

POPX dst **or** POPX.W dst

POPX.B dst

Operation Restore the 8-, 16-, 20-bit value from the stack to the destination. 20-bit addresses are possible. The SP is incremented by two (byte and word operands) and by four (address-word operand).

Emulation MOVX(.B,.A) @SP+,dst

Description The item on TOS is written to the destination operand. Register mode, Indexed mode, Symbolic mode, and Absolute mode are possible. The SP is incremented by two or four.

Note: the SP is incremented by two also for byte operations.

Status Bits Status bits are not affected.

Mode Bits OSCOFF, CPUOFF, and GIE are not affected.

Example Write the 16-bit value on TOS to the 20-bit address &EDE

```
POPX.W    &EDE    ; Write word to address EDE
```

Example Write the 20-bit value on TOS to R9

```
POPX.A    R9      ; Write address-word to R9
```


1.6.3.21 PUSHX

PUSHX.A Save single address-word to the stack

PUSHX.[W] Save single word to the stack

PUSHX.B Save single byte to the stack

Syntax `PUSHX.A src`

`PUSHX src OR PUSHX.W src`

`PUSHX.B src`

Operation Save the 8-, 16-, 20-bit value of the source operand on the TOS. 20-bit addresses are possible. The SP is decremented by two (byte and word operands) or by four (address-word operand) before the write operation.

Description The SP is decremented by two (byte and word operands) or by four (address-word operand). Then the source operand is written to the TOS. All seven addressing modes are possible for the source operand.

Status Bits Status bits are not affected.

Mode Bits OSCOFF, CPUOFF, and GIE are not affected.

Example Save the byte at the 20-bit address &EDE on the stack

```
PUSHX.B    &EDE    ; Save byte at address EDE
```

Example Save the 20-bit value in R9 on the stack.

```
PUSHX.A    R9      ; Save address-word in R9
```

1.6.3.22 RLAM

RLAM.A Rotate left arithmetically the 20-bit CPU register content

RLAM.[W] Rotate left arithmetically the 16-bit CPU register content

Syntax `RLAM.A #n,Rdst` $1 \leq n \leq 4$

`RLAM.W #n,Rdst` **OR** `RLAM #n,Rdst` $1 \leq n \leq 4$

Operation $C \leftarrow MSB \leftarrow MSB-1 \dots LSB+1 \leftarrow LSB \leftarrow 0$

Description The destination operand is shifted arithmetically left one, two, three, or four positions as shown in Figure 1-44. RLAM works as a multiplication (signed and unsigned) with 2, 4, 8, or 16. The word instruction RLAM.W clears the bits Rdst.19:16.

Note : This instruction does not use the extension word.

Status Bits N: Set if result is negative

.A: Rdst.19 = 1, reset if Rdst.19 = 0

.W: Rdst.15 = 1, reset if Rdst.15 = 0

Z: Set if result is zero, reset otherwise

C: Loaded from the MSB (n = 1), MSB-1 (n = 2), MSB-2 (n = 3), MSB-3 (n = 4)

V: Undefined

Mode Bits OSCOFF, CPUOFF, and GIE are not affected.

Example The 20-bit operand in R5 is shifted left by three positions. It operates equal to an arithmetic multiplication by 8.

`RLAM.A #3,R5 ; R5 = R5 x 8`

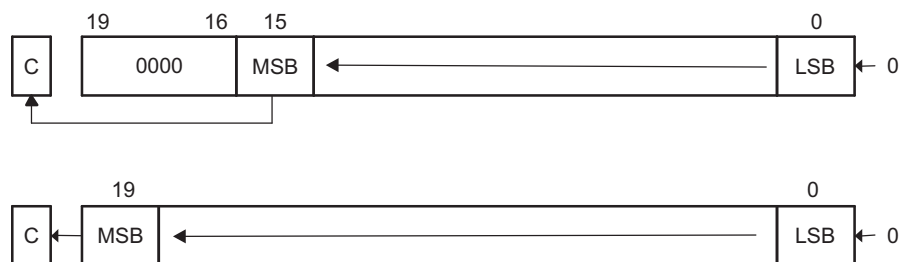


Figure 1-44. Rotate Left Arithmetically—RLAM[W] and RLAM.A

1.6.3.23 RLAX

* **RLAX.A** Rotate left arithmetically address-word

* **RLAX.[W]** Rotate left arithmetically word

* **RLAX.B** Rotate left arithmetically byte

Syntax RLAX.A dst

RLAX dst OR RLAX.W dst

RLAX.B dst

Operation C ← MSB ← MSB-1 ... LSB+1 ← LSB ← 0

Emulation ADDX.A dst, dst

ADDX dst, dst

ADDX.B dst, dst

Description The destination operand is shifted left one position as shown in Figure 1-45. The MSB is shifted into the carry bit (C) and the LSB is filled with 0. The RLAX instruction acts as a signed multiplication by 2.

Status Bits N: Set if result is negative, reset if positive

Z: Set if result is zero, reset otherwise

C: Loaded from the MSB

V: Set if an arithmetic overflow occurs: the initial value is 040000h ≤ dst < 0C0000h; reset otherwise

Set if an arithmetic overflow occurs: the initial value is 04000h ≤ dst < 0C000h; reset otherwise

Set if an arithmetic overflow occurs: the initial value is 040h ≤ dst < 0C0h; reset otherwise

Mode Bits OSCOFF, CPUOFF, and GIE are not affected.

Example The 20-bit value in R7 is multiplied by 2

RLAX.A R7 ; Shift left R7 (20-bit)

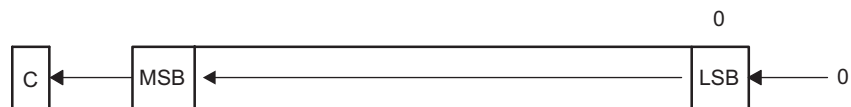


Figure 1-45. Destination Operand-Arithmetic Shift Left

1.6.3.24 RLCX

* **RLCX.A** Rotate left through carry address-word

* **RLCX.[W]** Rotate left through carry word

* **RLCX.B** Rotate left through carry byte

Syntax RLCX.A dst

RLCX dst OR RLCX.W dst

RLCX.B dst

Operation $C \leftarrow \text{MSB} \leftarrow \text{MSB}-1 \dots \text{LSB}+1 \leftarrow \text{LSB} \leftarrow C$

Emulation ADDCX.A dst, dst

ADDCX dst, dst

ADDCX.B dst, dst

Description The destination operand is shifted left one position as shown in Figure 1-46. The carry bit (C) is shifted into the LSB and the MSB is shifted into the carry bit (C).

Status Bits N: Set if result is negative, reset if positive

Z: Set if result is zero, reset otherwise

C: Loaded from the MSB

V: Set if an arithmetic overflow occurs: the initial value is $040000\text{h} \leq \text{dst} < 0\text{C}0000\text{h}$; reset otherwise

Set if an arithmetic overflow occurs: the initial value is $04000\text{h} \leq \text{dst} < 0\text{C}000\text{h}$; reset otherwise

Set if an arithmetic overflow occurs: the initial value is $040\text{h} \leq \text{dst} < 0\text{C}0\text{h}$; reset otherwise

Mode Bits OSCOFF, CPUOFF, and GIE are not affected.

Example The 20-bit value in R5 is shifted left one position.

```
RLCX.A R5 ; (R5 x 2) + C -> R5
```

Example The RAM byte LEO is shifted left one position. PC is pointing to upper memory.

```
RLCX.B LEO ; RAM(LEO) x 2 + C -> RAM(LEO)
```

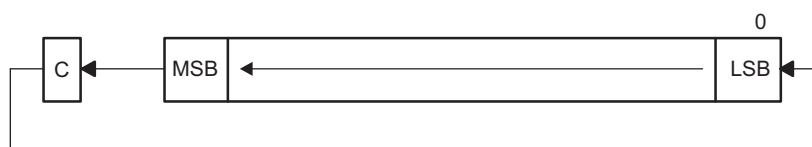


Figure 1-46. Destination Operand-Carry Left Shift

1.6.3.25 RRAM

RRAM.A Rotate right arithmetically the 20-bit CPU register content

RRAM.[W] Rotate right arithmetically the 16-bit CPU register content

Syntax
 RRAM.A #n,Rdst 1 ≤ n ≤ 4
 RRAM.W #n,Rdst OR RRAM #n,Rdst 1 ≤ n ≤ 4

Operation MSB → MSB → MSB-1 ... LSB+1 → LSB → C

Description The destination operand is shifted right arithmetically by one, two, three, or four bit positions as shown in Figure 1-47. The MSB retains its value (sign). RRAM operates equal to a signed division by 2, 4, 8, or 16. The MSB is retained and shifted into MSB-1. The LSB+1 is shifted into the LSB, and the LSB is shifted into the carry bit C. The word instruction RRAM.W clears the bits Rdst.19:16.

Note : This instruction does not use the extension word.

Status Bits
 N: Set if result is negative
 .A: Rdst.19 = 1, reset if Rdst.19 = 0
 .W: Rdst.15 = 1, reset if Rdst.15 = 0
 Z: Set if result is zero, reset otherwise
 C: Loaded from the LSB (n = 1), LSB+1 (n = 2), LSB+2 (n = 3), or LSB+3 (n = 4)
 V: Reset

Mode Bits OSCOFF, CPUOFF, and GIE are not affected.

Example The signed 20-bit number in R5 is shifted arithmetically right two positions.

```
RRAM.A #2,R5 ; R5/4 -> R5
```

Example The signed 20-bit value in R15 is multiplied by 0.75. (0.5 + 0.25) × R15.

```
PUSHM.A #1,R15 ; Save extended R15 on stack
RRAM.A #1,R15 ; R15 y 0.5 -> R15
ADDX.A @SP+,R15 ; R15 y 0.5 + R15 = 1.5 y R15 -> R15
RRAM.A #1,R15 ; (1.5 y R15) y 0.5 = 0.75 y R15 -> R15
```

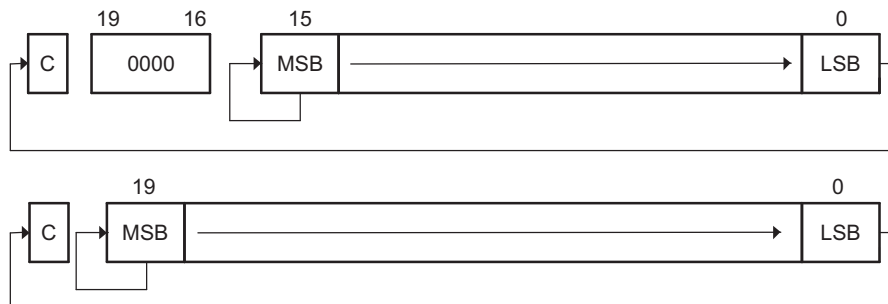


Figure 1-47. Rotate Right Arithmetically RRAM.[W] and RRAM.A

1.6.3.26 RRAX

RRAX.A Rotate right arithmetically the 20-bit operand

RRAX.[W] Rotate right arithmetically the 16-bit operand

RRAX.B Rotate right arithmetically the 8-bit operand

Syntax RRAX.A Rdst

RRAX.W Rdst

RRAX Rdst

RRAX.B Rdst

RRAX.A dst

RRAX dst **or** RRAX.W dst

RRAX.B dst

Operation MSB → MSB → MSB−1 ... LSB+1 → LSB → C

Description Register mode for the destination: the destination operand is shifted right by one bit position as shown in [Figure 1-48](#). The MSB retains its value (sign). The word instruction RRAX.W clears the bits Rdst.19:16, the byte instruction RRAX.B clears the bits Rdst.19:8. The MSB retains its value (sign), the LSB is shifted into the carry bit. RRAX here operates equal to a signed division by 2.

All other modes for the destination: the destination operand is shifted right arithmetically by one bit position as shown in [Figure 1-49](#). The MSB retains its value (sign), the LSB is shifted into the carry bit. RRAX here operates equal to a signed division by 2. All addressing modes, with the exception of the Immediate mode, are possible in the full memory.

Status Bits N: Set if result is negative, reset if positive

.A: dst.19 = 1, reset if dst.19 = 0

.W: dst.15 = 1, reset if dst.15 = 0

.B: dst.7 = 1, reset if dst.7 = 0

Z: Set if result is zero, reset otherwise

C: Loaded from the LSB

V: Reset

Mode Bits OSCOFF, CPUOFF, and GIE are not affected.

Example The signed 20-bit number in R5 is shifted arithmetically right four positions.

```
RPT      #4
RRAX.A  R5      ; R5/16 -> R5
```

Example The signed 8-bit value in EDE is multiplied by 0.5.

RRAX.B &EDE ; EDE/2 -> EDE

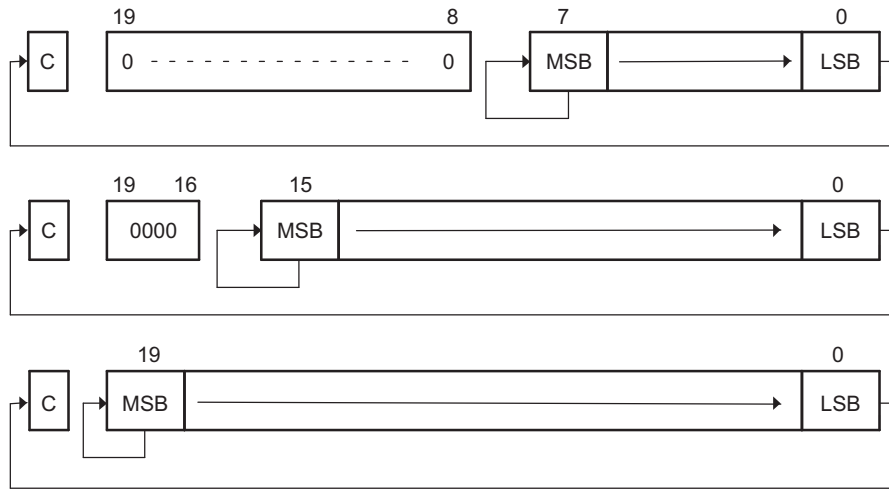


Figure 1-48. Rotate Right Arithmetically RRAX(B,A) – Register Mode

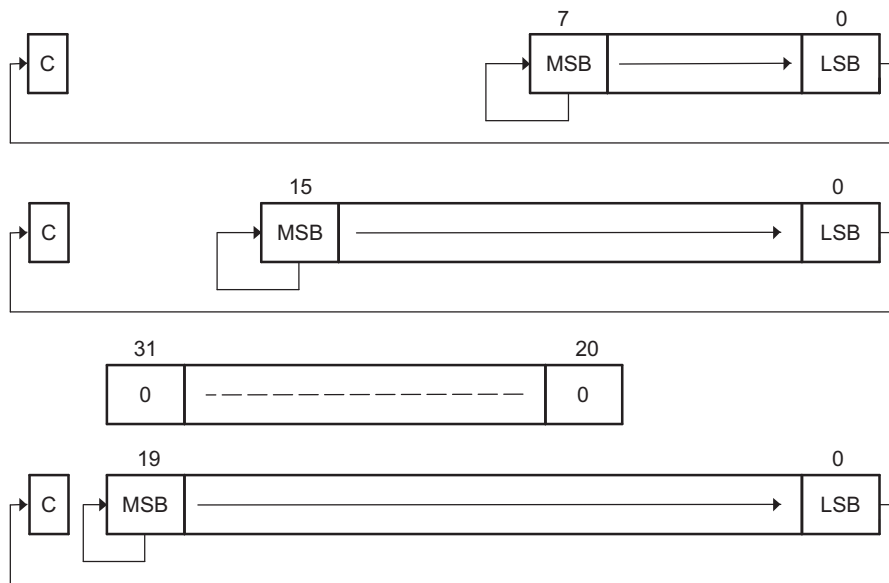


Figure 1-49. Rotate Right Arithmetically RRAX(B,A) – Non-Register Mode

1.6.3.27 RRCM

RRCM.A	Rotate right through carry the 20-bit CPU register content
RRCM.[W]	Rotate right through carry the 16-bit CPU register content
Syntax	<p>RRCM.A #n,Rdst 1 ≤ n ≤ 4</p> <p>RRCM.W #n,Rdst OR RRCM #n,Rdst 1 ≤ n ≤ 4</p>
Operation	C → MSB → MSB−1 ... LSB+1 → LSB → C
Description	<p>The destination operand is shifted right by one, two, three, or four bit positions as shown in Figure 1-50. The carry bit C is shifted into the MSB, the LSB is shifted into the carry bit. The word instruction RRCM.W clears the bits Rdst.19:16.</p> <p>Note : This instruction does not use the extension word.</p>
Status Bits	<p>N: Set if result is negative .A: Rdst.19 = 1, reset if Rdst.19 = 0 .W: Rdst.15 = 1, reset if Rdst.15 = 0</p> <p>Z: Set if result is zero, reset otherwise</p> <p>C: Loaded from the LSB (n = 1), LSB+1 (n = 2), LSB+2 (n = 3), or LSB+3 (n = 4)</p> <p>V: Reset</p>
Mode Bits	OSCOFF, CPUOFF, and GIE are not affected.

Example The address-word in R5 is shifted right by three positions. The MSB-2 is loaded with 1.

```
SETC                ; Prepare carry for MSB-2
RRCM.A #3,R5        ; R5 = R5 » 3 + 20000h
```

Example The word in R6 is shifted right by two positions. The MSB is loaded with the LSB. The MSB-1 is loaded with the contents of the carry flag.

```
RRCM.W #2,R6        ; R6 = R6 » 2. R6.19:16 = 0
```

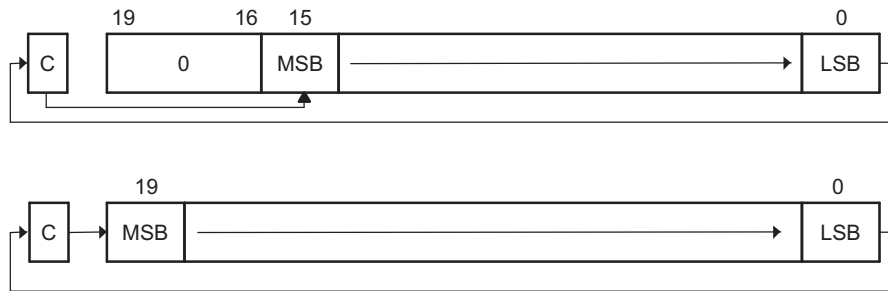


Figure 1-50. Rotate Right Through Carry RRCM[.W] and RRCM.A

1.6.3.28 RRCX

RRCX.A	Rotate right through carry the 20-bit operand
RRCX.[W]	Rotate right through carry the 16-bit operand
RRCX.B	Rotate right through carry the 8-bit operand
Syntax	<pre>RRCX.A Rdst RRCX.W Rdst RRCX Rdst RRCX.B Rdst RRCX.A dst RRCX dst or RRCX.W dst RRCX.B dst</pre>
Operation Description	<p>$C \rightarrow \text{MSB} \rightarrow \text{MSB}-1 \dots \text{LSB}+1 \rightarrow \text{LSB} \rightarrow C$</p> <p>Register mode for the destination: the destination operand is shifted right by one bit position as shown in Figure 1-51. The word instruction RRCX.W clears the bits Rdst.19:16, the byte instruction RRCX.B clears the bits Rdst.19:8. The carry bit C is shifted into the MSB, the LSB is shifted into the carry bit.</p> <p>All other modes for the destination: the destination operand is shifted right by one bit position as shown in Figure 1-52. The carry bit C is shifted into the MSB, the LSB is shifted into the carry bit. All addressing modes, with the exception of the Immediate mode, are possible in the full memory.</p>
Status Bits	<p>N: Set if result is negative .A: dst.19 = 1, reset if dst.19 = 0 .W: dst.15 = 1, reset if dst.15 = 0 .B: dst.7 = 1, reset if dst.7 = 0</p> <p>Z: Set if result is zero, reset otherwise</p> <p>C: Loaded from the LSB</p> <p>V: Reset</p>
Mode Bits	OSCOFF, CPUOFF, and GIE are not affected.
Example	The 20-bit operand at address EDE is shifted right by one position. The MSB is loaded with 1.
	<pre>SETC ; Prepare carry for MSB RRCX.A EDE ; EDE = EDE » 1 + 80000h</pre>
Example	The word in R6 is shifted right by 12 positions.

```
RPT      #12
RRCX.W  R6      ; R6 = R6 » 12. R6.19:16 = 0
```

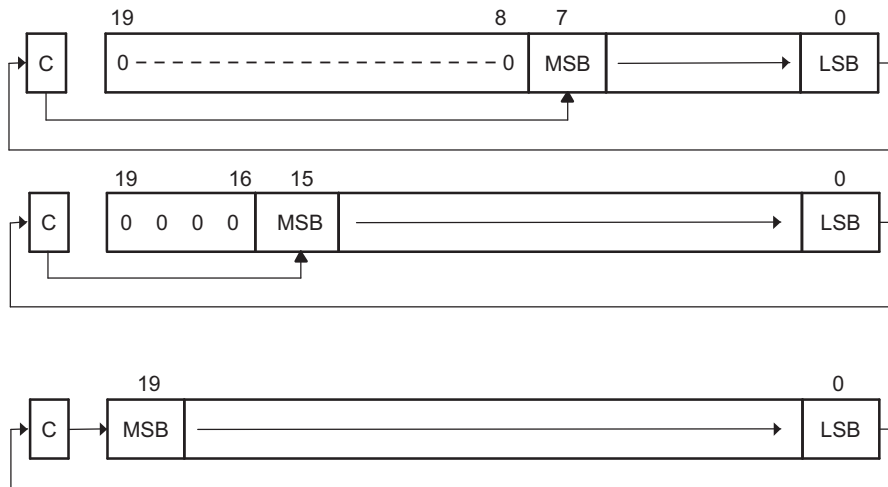


Figure 1-51. Rotate Right Through Carry RRCX(.B,.A) – Register Mode

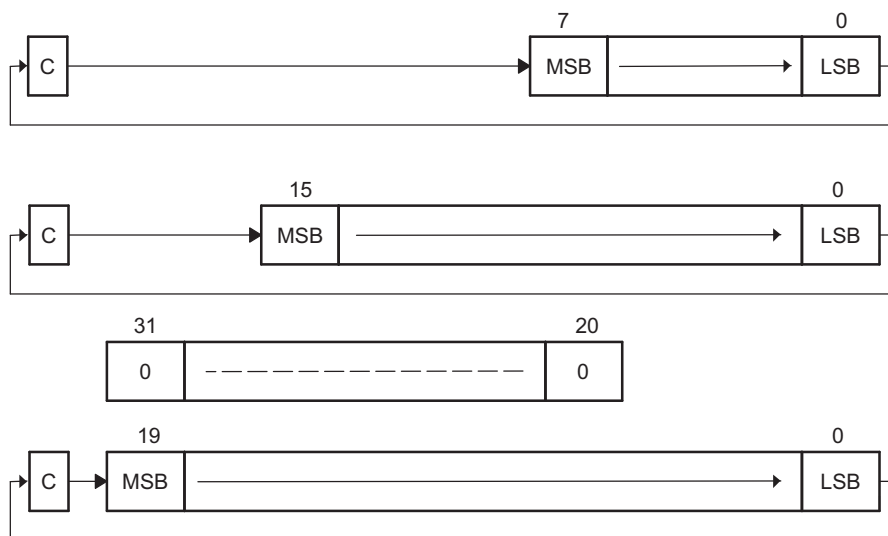


Figure 1-52. Rotate Right Through Carry RRCX(.B,.A) – Non-Register Mode

1.6.3.29 RRUM

RRUM.A Rotate right through carry the 20-bit CPU register content

RRUM.[W] Rotate right through carry the 16-bit CPU register content

Syntax `RRUM.A #n,Rdst` $1 \leq n \leq 4$

`RRUM.W #n,Rdst` OR `RRUM #n,Rdst` $1 \leq n \leq 4$

Operation $0 \rightarrow \text{MSB} \rightarrow \text{MSB}-1 \dots \text{LSB}+1 \rightarrow \text{LSB} \rightarrow \text{C}$

Description The destination operand is shifted right by one, two, three, or four bit positions as shown in Figure 1-53. Zero is shifted into the MSB, the LSB is shifted into the carry bit. RRUM works like an unsigned division by 2, 4, 8, or 16. The word instruction RRUM.W clears the bits Rdst.19:16.

Note : This instruction does not use the extension word.

Status Bits N: Set if result is negative

.A: Rdst.19 = 1, reset if Rdst.19 = 0

.W: Rdst.15 = 1, reset if Rdst.15 = 0

Z: Set if result is zero, reset otherwise

C: Loaded from the LSB (n = 1), LSB+1 (n = 2), LSB+2 (n = 3), or LSB+3 (n = 4)

V: Reset

Mode Bits OSCOFF, CPUOFF, and GIE are not affected.

Example The unsigned address-word in R5 is divided by 16.

```
RRUM.A #4,R5 ; R5 = R5 » 4. R5/16
```

Example The word in R6 is shifted right by one bit. The MSB R6.15 is loaded with 0.

```
RRUM.W #1,R6 ; R6 = R6/2. R6.19:15 = 0
```

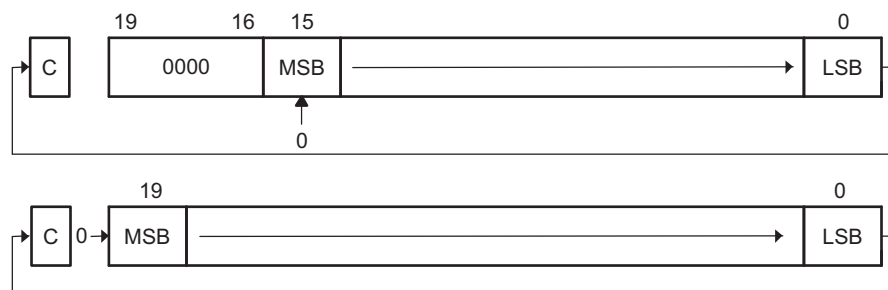


Figure 1-53. Rotate Right Unsigned RRUM[W] and RRUM.A

1.6.3.30 RRUX

RRUX.A Shift right unsigned the 20-bit CPU register content

RRUX.[W] Shift right unsigned the 16-bit CPU register content

RRUX.B Shift right unsigned the 8-bit CPU register content

Syntax RRUX.A Rdst

RRUX.W Rdst

RRUX Rdst

RRUX.B Rdst

Operation C=0 → MSB → MSB-1 ... LSB+1 → LSB → C

Description RRUX is valid for register mode only: the destination operand is shifted right by one bit position as shown in Figure 1-54. The word instruction RRUX.W clears the bits Rdst.19:16. The byte instruction RRUX.B clears the bits Rdst.19:8. Zero is shifted into the MSB, the LSB is shifted into the carry bit.

Status Bits

- N: Set if result is negative
 - .A: dst.19 = 1, reset if dst.19 = 0
 - .W: dst.15 = 1, reset if dst.15 = 0
 - .B: dst.7 = 1, reset if dst.7 = 0
- Z: Set if result is zero, reset otherwise
- C: Loaded from the LSB
- V: Reset

Mode Bits OSCOFF, CPUOFF, and GIE are not affected.

Example The word in R6 is shifted right by 12 positions.

```
RPT      #12
RRUX.W  R6      ; R6 = R6 >> 12. R6.19:16 = 0
```

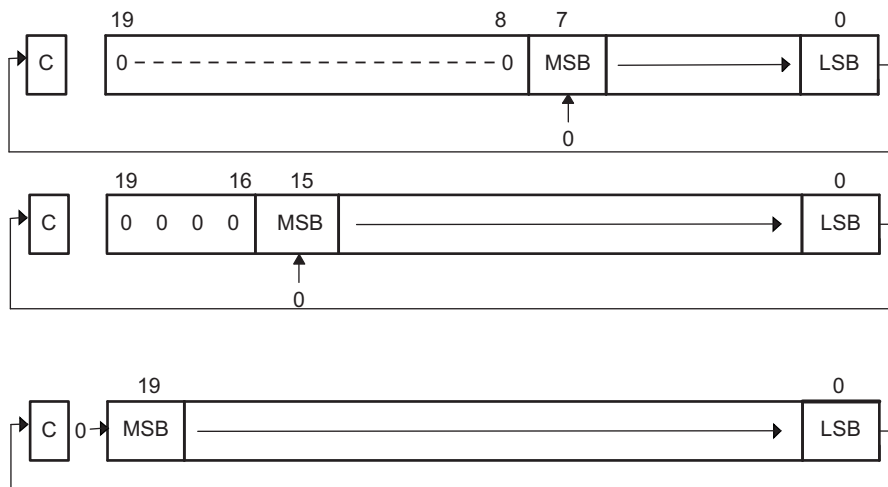


Figure 1-54. Rotate Right Unsigned RRUX(.B,.A) – Register Mode

1.6.3.31 SBCX

* SBCX.A	Subtract borrow (.NOT. carry) from destination address-word
* SBCX.[W]	Subtract borrow (.NOT. carry) from destination word
* SBCX.B	Subtract borrow (.NOT. carry) from destination byte
Syntax	SBCX.A dst SBCX dst OR SBCX.W dst SBCX.B dst
Operation	dst + 0FFFFFFh + C → dst dst + 0FFFFFFh + C → dst dst + 0FFh + C → dst
Emulation	SBCX.A #0, dst SBCX #0, dst SBCX.B #0, dst
Description	The carry bit (C) is added to the destination operand minus one. The previous contents of the destination are lost.
Status Bits	N: Set if result is negative, reset if positive Z: Set if result is zero, reset otherwise C: Set if there is a carry from the MSB of the result, reset otherwise Set to 1 if no borrow, reset if borrow V: Set if an arithmetic overflow occurs, reset otherwise
Mode Bits	OSCOFF, CPUOFF, and GIE are not affected.
Example	The 8-bit counter pointed to by R13 is subtracted from a 16-bit counter pointed to by R12.

```

SUBX.B   @R13,0(R12)       ; Subtract LSDs
SBCX.B   1(R12)           ; Subtract carry from MSD

```

NOTE: Borrow implementation

The borrow is treated as a .NOT. carry:

Borrow	Carry Bit
Yes	0
No	1

1.6.3.32 SUBX

SUBX.A	Subtract source address-word from destination address-word
SUBX.[W]	Subtract source word from destination word
SUBX.B	Subtract source byte from destination byte
Syntax	SUBX.A <i>src,dst</i> SUBX <i>src,dst</i> Or SUBX.W <i>src,dst</i> SUBX.B <i>src,dst</i>
Operation	(.not. <i>src</i>) + 1 + <i>dst</i> → <i>dst</i> or <i>dst</i> – <i>src</i> → <i>dst</i>
Description	The source operand is subtracted from the destination operand. This is done by adding the 1s complement of the source + 1 to the destination. The source operand is not affected. The result is written to the destination operand. Both operands may be located in the full address space.
Status Bits	N: Set if result is negative (<i>src</i> > <i>dst</i>), reset if positive (<i>src</i> ≤ <i>dst</i>) Z: Set if result is zero (<i>src</i> = <i>dst</i>), reset otherwise (<i>src</i> ≠ <i>dst</i>) C: Set if there is a carry from the MSB, reset otherwise V: Set if the subtraction of a negative source operand from a positive destination operand delivers a negative result, or if the subtraction of a positive source operand from a negative destination operand delivers a positive result, reset otherwise (no overflow)
Mode Bits	OSCOFF, CPUOFF, and GIE are not affected.
Example	A 20-bit constant 87654h is subtracted from EDE (LSBs) and EDE+2 (MSBs).

```
SUBX.A    #87654h,EDE        ; Subtract 87654h from EDE+2|EDE
```

Example A table word pointed to by R5 (20-bit address) is subtracted from R7. Jump to label TONI if R7 contains zero after the instruction. R5 is auto-incremented by two. R7.19:16 = 0.

```
SUBX.W    @R5+,R7           ; Subtract table number from R7. R5 + 2
JZ        TONI              ; R7 = @R5 (before subtraction)
...       ; R7 <> @R5 (before subtraction)
```

Example Byte CNT is subtracted from the byte R12 points to in the full address space. Address of CNT is within PC ± 512 K.

```
SUBX.B    CNT,0(R12)        ; Subtract CNT from @R12
```

Note: Use SUBA for the following two cases for better density and execution.

```
SUBX.A    Rsrc,Rdst
SUBX.A    #imm20,Rdst
```

1.6.3.33 SUBCX

SUBCX.A	Subtract source address-word with carry from destination address-word
SUBCX.[W]	Subtract source word with carry from destination word
SUBCX.B	Subtract source byte with carry from destination byte
Syntax	<p>SUBCX.A <i>src</i>,<i>dst</i></p> <p>SUBCX <i>src</i>,<i>dst</i> Or SUBCX.W <i>src</i>,<i>dst</i></p> <p>SUBCX.B <i>src</i>,<i>dst</i></p>
Operation	$(.not. src) + C + dst \rightarrow dst$ or $dst - (src - 1) + C \rightarrow dst$
Description	The source operand is subtracted from the destination operand. This is made by adding the 1s complement of the source + carry to the destination. The source operand is not affected, the result is written to the destination operand. Both operands may be located in the full address space.
Status Bits	<p>N: Set if result is negative (MSB = 1), reset if positive (MSB = 0)</p> <p>Z: Set if result is zero, reset otherwise</p> <p>C: Set if there is a carry from the MSB, reset otherwise</p> <p>V: Set if the subtraction of a negative source operand from a positive destination operand delivers a negative result, or if the subtraction of a positive source operand from a negative destination operand delivers a positive result, reset otherwise (no overflow).</p>
Mode Bits	OSCOFF, CPUOFF, and GIE are not affected.
Example	A 20-bit constant 87654h is subtracted from R5 with the carry from the previous instruction.

```
SUBCX.A    #87654h,R5        ; Subtract 87654h + C from R5
```

Example A 48-bit number (3 words) pointed to by R5 (20-bit address) is subtracted from a 48-bit counter in RAM, pointed to by R7. R5 auto-increments to point to the next 48-bit number.

```
SUBX.W    @R5+,0(R7)        ; Subtract LSBs. R5 + 2
SUBCX.W   @R5+,2(R7)        ; Subtract MIDs with C. R5 + 2
SUBCX.W   @R5+,4(R7)        ; Subtract MSBs with C. R5 + 2
```

Example Byte CNT is subtracted from the byte R12 points to. The carry of the previous instruction is used. 20-bit addresses.

```
SUBCX.B   &CNT,0(R12)       ; Subtract byte CNT from @R12
```


1.6.3.34 SWPBX

SWPBX.A Swap bytes of lower word

SWPBX.[W] Swap bytes of word

Syntax SWPBX.A dst
SWPBX dst OR SWPBX.W dst

Operation dst.15:8 ↔ dst.7:0

Description Register mode: Rn.15:8 are swapped with Rn.7:0. When the .A extension is used, Rn.19:16 are unchanged. When the .W extension is used, Rn.19:16 are cleared.
Other modes: When the .A extension is used, bits 31:20 of the destination address are cleared, bits 19:16 are left unchanged, and bits 15:8 are swapped with bits 7:0. When the .W extension is used, bits 15:8 are swapped with bits 7:0 of the addressed word.

Status Bits Status bits are not affected.

Mode Bits OSCOFF, CPUOFF, and GIE are not affected.

Example Exchange the bytes of RAM address-word EDE

```
MOVX.A #23456h, &EDE ; 23456h -> EDE
SWPBX.A EDE ; 25634h -> EDE
```

Example Exchange the bytes of R5

```
MOVA #23456h, R5 ; 23456h -> R5
SWPBX.W R5 ; 05634h -> R5
```

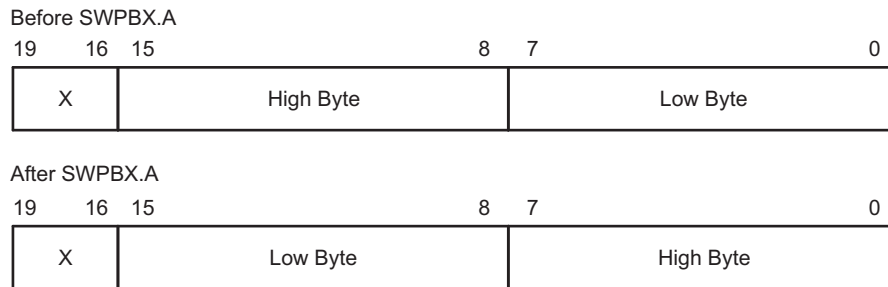


Figure 1-55. Swap Bytes SWPBX.A Register Mode

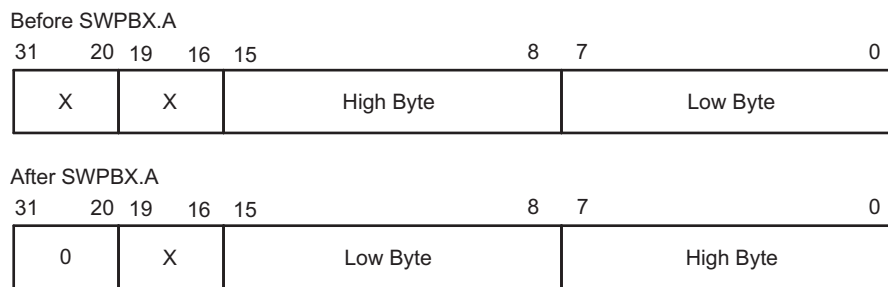


Figure 1-56. Swap Bytes SWPBX.A In Memory

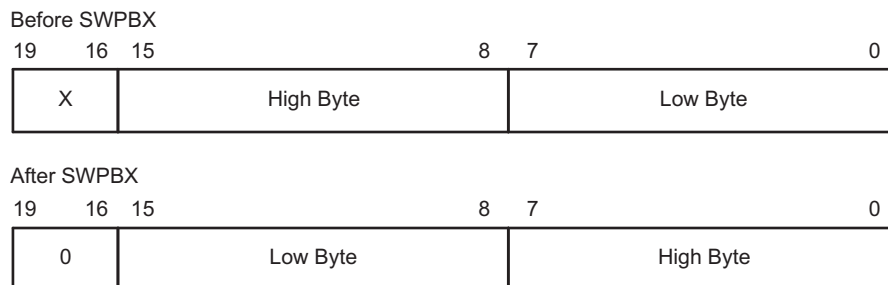


Figure 1-57. Swap Bytes SWPBX[.W] Register Mode

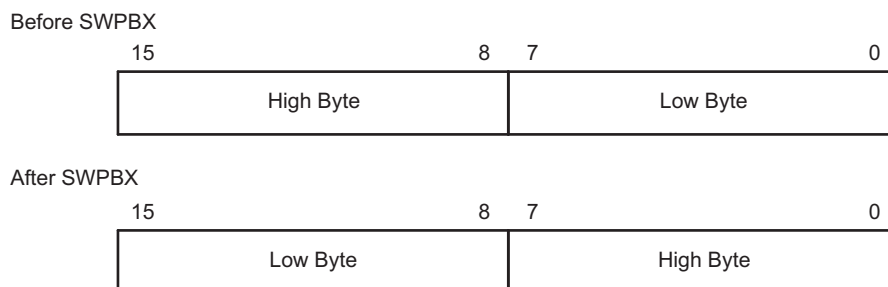


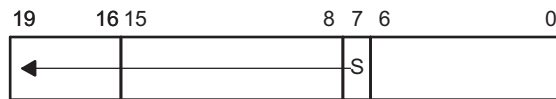
Figure 1-58. Swap Bytes SWPBX[.W] In Memory

1.6.3.35 SXTX

SXTX.A Extend sign of lower byte to address-word
SXTX.[W] Extend sign of lower byte to word
Syntax SXTX.A dst
 SXTX dst OR SXTX.W dst
Operation dst.7 → dst.15:8, Rdst.7 → Rdst.19:8 (Register mode)
Description Register mode: The sign of the low byte of the operand (Rdst.7) is extended into the bits Rdst.19:8.
 Other modes: SXTX.A: the sign of the low byte of the operand (dst.7) is extended into dst.19:8. The bits dst.31:20 are cleared.
 SXTX.[W]: the sign of the low byte of the operand (dst.7) is extended into dst.15:8.
Status Bits N: Set if result is negative, reset otherwise
 Z: Set if result is zero, reset otherwise
 C: Set if result is not zero, reset otherwise (C = .not.Z)
 V: Reset
Mode Bits OSCOFF, CPUOFF, and GIE are not affected.
Example The signed 8-bit data in EDE.7:0 is sign extended to 20 bits: EDE.19:8. Bits 31:20 located in EDE+2 are cleared.

SXTX.A &EDE ; Sign extended EDE -> EDE+2/EDE

SXTX.A Rdst



SXTX.A dst

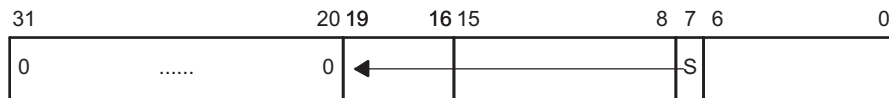
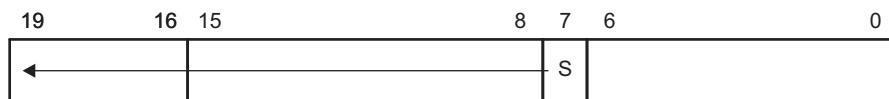


Figure 1-59. Sign Extend SXTX.A

SXTX.[W] Rdst



SXTX.[W] dst

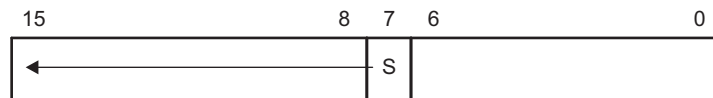


Figure 1-60. Sign Extend SXTX.[W]

1.6.3.36 TSTX

* **TSTX.A** Test destination address-word

* **TSTX.[W]** Test destination word

* **TSTX.B** Test destination byte

Syntax TSTX.A dst

TSTX dst **OR** TSTX.W dst

TSTX.B dst

Operation dst + 0FFFFFFh + 1

dst + 0FFFFFFh + 1

dst + 0FFh + 1

Emulation CMPX.A #0,dst

CMPX #0,dst

CMPX.B #0,dst

Description The destination operand is compared with zero. The status bits are set according to the result. The destination is not affected.

Status Bits N: Set if destination is negative, reset if positive

Z: Set if destination contains zero, reset otherwise

C: Set

V: Reset

Mode Bits OSCOFF, CPUOFF, and GIE are not affected.

Example RAM byte LEO is tested; PC is pointing to upper memory. If it is negative, continue at LEONEG; if it is positive but not zero, continue at LEOPOS.

```

TSTX.B LEO ; Test LEO
JN LEONEG ; LEO is negative
JZ LEOZERO ; LEO is zero
LEOPOS ..... ; LEO is positive but not zero
LEONEG ..... ; LEO is negative
LEOZERO ..... ; LEO is zero

```

1.6.3.37 XORX

XORX.A Exclusive OR source address-word with destination address-word

XORX.[W] Exclusive OR source word with destination word

XORX.B Exclusive OR source byte with destination byte

Syntax XORX.A *src,dst*

XORX *src,dst* OR XORX.W *src,dst*

XORX.B *src,dst*

Operation *src .xor. dst* → *dst*

Description The source and destination operands are exclusively ORed. The result is placed into the destination. The source operand is not affected. The previous contents of the destination are lost. Both operands may be located in the full address space.

Status Bits N: Set if result is negative (MSB = 1), reset if positive (MSB = 0)

Z: Set if result is zero, reset otherwise

C: Set if result is not zero, reset otherwise (carry = .not. Zero)

V: Set if both operands are negative (before execution), reset otherwise

Mode Bits OSCOFF, CPUOFF, and GIE are not affected.

Example Toggle bits in address-word CNTR (20-bit data) with information in address-word TONI (20-bit address)

```
XORX.A  TONI,&CNTR      ; Toggle bits in CNTR
```

Example A table word pointed to by R5 (20-bit address) is used to toggle bits in R6.

```
XORX.W  @R5,R6          ; Toggle bits in R6. R6.19:16 = 0
```

Example Reset to zero those bits in the low byte of R7 that are different from the bits in byte EDE (20-bit address)

```
XORX.B  EDE,R7          ; Set different bits to 1 in R7
INV.B   R7               ; Invert low byte of R7. R7.19:8 = 0.
```

1.6.4 Address Instructions

MSP430X address instructions are instructions that support 20-bit operands but have restricted addressing modes. The addressing modes are restricted to the Register mode and the Immediate mode, except for the MOVA instruction. Restricting the addressing modes removes the need for the additional extension-word op-code improving code density and execution time. The MSP430X address instructions are listed and described in the following pages.

1.6.4.1 ADDA

ADDA	Add 20-bit source to a 20-bit destination register
Syntax	ADDA Rsrc,Rdst ADDA #imm20,Rdst
Operation	src + Rdst → Rdst
Description	The 20-bit source operand is added to the 20-bit destination CPU register. The previous contents of the destination are lost. The source operand is not affected.
Status Bits	N: Set if result is negative (Rdst.19 = 1), reset if positive (Rdst.19 = 0) Z: Set if result is zero, reset otherwise C: Set if there is a carry from the 20-bit result, reset otherwise V: Set if the result of two positive operands is negative, or if the result of two negative numbers is positive, reset otherwise
Mode Bits	OSCOFF, CPUOFF, and GIE are not affected.
Example	R5 is increased by 0A4320h. The jump to TONI is performed if a carry occurs.
	<pre>ADDA #0A4320h,R5 ; Add A4320h to 20-bit R5 JC TONI ; Jump on carry ... ; No carry occurred</pre>

1.6.4.2 BRA

*** BRA** Branch to destination

Syntax BRA dst

Operation dst → PC

Emulation MOVA dst,PC

Description An unconditional branch is taken to a 20-bit address anywhere in the full address space. All seven source addressing modes can be used. The branch instruction is an address-word instruction. If the destination address is contained in a memory location X, it is contained in two ascending words: X (LSBs) and (X + 2) (MSBs).

Status Bits N: Not affected

Z: Not affected

C: Not affected

V: Not affected

Mode Bits OSCOFF, CPUOFF, and GIE are not affected.

Examples Examples for all addressing modes are given.

Immediate mode: Branch to label EDE located anywhere in the 20-bit address space or branch directly to address.

```
BRA    #EDE          ; MOVA    #imm20,PC
BRA    #01AA04h
```

Symbolic mode: Branch to the 20-bit address contained in addresses EXEC (LSBs) and EXEC+2 (MSBs). EXEC is located at the address (PC + X) where X is within ±32 K. Indirect addressing.

```
BRA    EXEC          ; MOVA    z16(PC),PC
```

Note: If the 16-bit index is not sufficient, a 20-bit index may be used with the following instruction.

```
MOVX.A EXEC,PC      ; 1M byte range with 20-bit index
```

Absolute mode: Branch to the 20-bit address contained in absolute addresses EXEC (LSBs) and EXEC+2 (MSBs). Indirect addressing.

```
BRA    &EXEC          ; MOVA    &abs20,PC
```

Register mode: Branch to the 20-bit address contained in register R5. Indirect R5.

```
BRA    R5             ; MOVA    R5,PC
```

Indirect mode: Branch to the 20-bit address contained in the word pointed to by register R5 (LSBs). The MSBs have the address (R5 + 2). Indirect, indirect R5.

```
BRA    @R5            ; MOVA    @R5,PC
```


Indirect, Auto-Increment mode: Branch to the 20-bit address contained in the words pointed to by register R5 and increment the address in R5 afterwards by 4. The next time the software flow uses R5 as a pointer, it can alter the program execution due to access to the next address in the table pointed to by R5. Indirect, indirect R5.

```
BRA    @R5+          ; MOVA    @R5+,PC. R5 + 4
```

Indexed mode: Branch to the 20-bit address contained in the address pointed to by register (R5 + X) (for example, a table with addresses starting at X). (R5 + X) points to the LSBs, (R5 + X + 2) points to the MSBs of the address. X is within $R5 \pm 32\text{ K}$. Indirect, indirect (R5 + X).

```
BRA    X(R5)        ; MOVA    z16(R5),PC
```

Note: If the 16-bit index is not sufficient, a 20-bit index X may be used with the following instruction:

```
MOVX.A X(R5),PC    ; 1M byte range with 20-bit index
```

1.6.4.3 CALLA

CALLA	Call a subroutine
Syntax	CALLA dst
Operation	dst → tmp 20-bit dst is evaluated and stored SP – 2 → SP PC.19:16 → @SP updated PC with return address to TOS (MSBs) SP – 2 → SP PC.15:0 → @SP updated PC to TOS (LSBs) tmp → PC saved 20-bit dst to PC
Description	A subroutine call is made to a 20-bit address anywhere in the full address space. All seven source addressing modes can be used. The call instruction is an address-word instruction. If the destination address is contained in a memory location X, it is contained in two ascending words, X (LSBs) and (X + 2) (MSBs). Two words on the stack are needed for the return address. The return is made with the instruction RETA.
Status Bits	N: Not affected Z: Not affected C: Not affected V: Not affected
Mode Bits	OSCOFF, CPUOFF, and GIE are not affected.
Examples	Examples for all addressing modes are given. Immediate mode: Call a subroutine at label EXEC or call directly an address.

```
CALLA #EXEC          ; Start address EXEC
CALLA #01AA04h      ; Start address 01AA04h
```

Symbolic mode: Call a subroutine at the 20-bit address contained in addresses EXEC (LSBs) and EXEC+2 (MSBs). EXEC is located at the address (PC + X) where X is within ±32 K. Indirect addressing.

```
CALLA EXEC          ; Start address at @EXEC. z16(PC)
```

Absolute mode: Call a subroutine at the 20-bit address contained in absolute addresses EXEC (LSBs) and EXEC+2 (MSBs). Indirect addressing.

```
CALLA &EXEC         ; Start address at @EXEC
```

Register mode: Call a subroutine at the 20-bit address contained in register R5. Indirect R5.

```
CALLA R5           ; Start address at @R5
```

Indirect mode: Call a subroutine at the 20-bit address contained in the word pointed to by register R5 (LSBs). The MSBs have the address (R5 + 2). Indirect, indirect R5.

```
CALLA @R5         ; Start address at @R5
```

Indirect, Auto-Increment mode: Call a subroutine at the 20-bit address contained in the words pointed to by register R5 and increment the 20-bit address in R5 afterwards by 4. The next time the software flow uses R5 as a pointer, it can alter the program execution due to access to the next word address in the table pointed to by R5. Indirect, indirect R5.

CALLA @R5+ ; Start address at @R5. R5 + 4

Indexed mode: Call a subroutine at the 20-bit address contained in the address pointed to by register (R5 + X); for example, a table with addresses starting at X. (R5 + X) points to the LSBs, (R5 + X + 2) points to the MSBs of the word address. X is within R5 ± 32 K. Indirect, indirect (R5 + X).

CALLA X(R5) ; Start address at @(R5+X). z16(R5)

1.6.4.4 CLRA

* CLRA	Clear 20-bit destination register
Syntax	CLRA Rdst
Operation	0 → Rdst
Emulation	MOVA #0, Rdst
Description	The destination register is cleared.
Status Bits	Status bits are not affected.
Example	The 20-bit value in R10 is cleared.

```
CLRA R10 ; 0 -> R10
```

1.6.4.5 CMPA

CMPA	Compare the 20-bit source with a 20-bit destination register
Syntax	CMPA Rsrc,Rdst CMPA #imm20,Rdst
Operation	(.not. src) + 1 + Rdst or Rdst – src
Description	The 20-bit source operand is subtracted from the 20-bit destination CPU register. This is made by adding the 1s complement of the source + 1 to the destination register. The result affects only the status bits.
Status Bits	N: Set if result is negative (src > dst), reset if positive (src ≤ dst) Z: Set if result is zero (src = dst), reset otherwise (src ≠ dst) C: Set if there is a carry from the MSB, reset otherwise V: Set if the subtraction of a negative source operand from a positive destination operand delivers a negative result, or if the subtraction of a positive source operand from a negative destination operand delivers a positive result, reset otherwise (no overflow)
Mode Bits	OSCOFF, CPUOFF, and GIE are not affected.
Example	A 20-bit immediate operand and R6 are compared. If they are equal, the program continues at label EQUAL.

```

CMPA #12345h,R6      ; Compare R6 with 12345h
JEQ  EQUAL          ; R6 = 12345h
...                 ; Not equal
    
```

Example The 20-bit values in R5 and R6 are compared. If R5 is greater than (signed) or equal to R6, the program continues at label GRE.

```

CMPA R6,R5          ; Compare R6 with R5 (R5 - R6)
JGE  GRE            ; R5 >= R6
...                 ; R5 < R6
    
```

1.6.4.6 DECDA

* DECDA	Double-decrement 20-bit destination register
Syntax	DECDA Rdst
Operation	Rdst – 2 → Rdst
Emulation	SUBA #2, Rdst
Description	The destination register is decremented by two. The original contents are lost.
Status Bits	<p>N: Set if result is negative, reset if positive</p> <p>Z: Set if Rdst contained 2, reset otherwise</p> <p>C: Reset if Rdst contained 0 or 1, set otherwise</p> <p>V: Set if an arithmetic overflow occurs, otherwise reset</p>
Mode Bits	OSCOFF, CPUOFF, and GIE are not affected.
Example	The 20-bit value in R5 is decremented by 2.

```
DECDA R5 ; Decrement R5 by two
```

1.6.4.7 INCD A

* INCD A	Double-increment 20-bit destination register
Syntax	INCD A Rdst
Operation	Rdst + 2 → Rdst
Emulation	ADDA #2, Rdst
Description	The destination register is incremented by two. The original contents are lost.
Status Bits	N: Set if result is negative, reset if positive Z: Set if Rdst contained 0FFFFFFh, reset otherwise Set if Rdst contained 0FFFEh, reset otherwise Set if Rdst contained 0FEh, reset otherwise C: Set if Rdst contained 0FFFFFFh or 0FFFFFFh, reset otherwise Set if Rdst contained 0FFFEh or 0FFFFh, reset otherwise Set if Rdst contained 0FEh or 0FFh, reset otherwise V: Set if Rdst contained 07FFFEh or 07FFFFh, reset otherwise Set if Rdst contained 07FFEh or 07FFFh, reset otherwise Set if Rdst contained 07Eh or 07Fh, reset otherwise
Mode Bits	OSCOFF, CPUOFF, and GIE are not affected.
Example	The 20-bit value in R5 is incremented by two.

```
INCD A    R5        ; Increment R5 by two
```

1.6.4.8 MOVA

MOVA	Move the 20-bit source to the 20-bit destination
Syntax	<pre> MOVA Rsrc,Rdst MOVA #imm20,Rdst MOVA z16(Rsrc),Rdst MOVA EDE,Rdst MOVA &abs20,Rdst MOVA @Rsrc,Rdst MOVA @Rsrc+,Rdst MOVA Rsrc,z16(Rdst) MOVA Rsrc,&abs20 </pre>
Operation	<pre> src → Rdst Rsrc → dst </pre>
Description	The 20-bit source operand is moved to the 20-bit destination. The source operand is not affected. The previous content of the destination is lost.
Status Bits	<pre> N: Not affected Z: Not affected C: Not affected V: Not affected </pre>
Mode Bits	OSCOFF, CPUOFF, and GIE are not affected.
Examples	Copy 20-bit value in R9 to R8
	<pre> MOVA R9,R8 ; R9 -> R8 </pre> <p>Write 20-bit immediate value 12345h to R12</p> <pre> MOVA #12345h,R12 ; 12345h -> R12 </pre> <p>Copy 20-bit value addressed by (R9 + 100h) to R8. Source operand in addresses (R9 + 100h) LSBs and (R9 + 102h) MSBs.</p> <pre> MOVA 100h(R9),R8 ; Index: + 32 K. 2 words transferred </pre> <p>Move 20-bit value in 20-bit absolute addresses EDE (LSBs) and EDE+2 (MSBs) to R12</p> <pre> MOVA &EDE,R12 ; &EDE -> R12. 2 words transferred </pre> <p>Move 20-bit value in 20-bit addresses EDE (LSBs) and EDE+2 (MSBs) to R12. PC index ± 32 K.</p> <pre> MOVA EDE,R12 ; EDE -> R12. 2 words transferred </pre> <p>Copy 20-bit value R9 points to (20 bit address) to R8. Source operand in addresses @R9 LSBs and @(R9 + 2) MSBs.</p> <pre> MOVA @R9,R8 ; @R9 -> R8. 2 words transferred </pre>

Copy 20-bit value R9 points to (20 bit address) to R8. R9 is incremented by four afterwards. Source operand in addresses @R9 LSBs and @(R9 + 2) MSBs.

MOVA @R9+,R8 ; @R9 -> R8. R9 + 4. 2 words transferred.

Copy 20-bit value in R8 to destination addressed by (R9 + 100h). Destination operand in addresses @(R9 + 100h) LSBs and @(R9 + 102h) MSBs.

MOVA R8,100h(R9) ; Index: +- 32 K. 2 words transferred

Move 20-bit value in R13 to 20-bit absolute addresses EDE (LSBs) and EDE+2 (MSBs)

MOVA R13,&EDE ; R13 -> EDE. 2 words transferred

Move 20-bit value in R13 to 20-bit addresses EDE (LSBs) and EDE+2 (MSBs). PC index \pm 32 K.

MOVA R13,EDE ; R13 -> EDE. 2 words transferred

1.6.4.9 RETA

* RETA	Return from subroutine
Syntax	RETA
Operation	<p>@SP → PC.15:0 LSBs (15:0) of saved PC to PC.15:0</p> <p>SP + 2 → SP</p> <p>@SP → PC.19:16 MSBs (19:16) of saved PC to PC.19:16</p> <p>SP + 2 → SP</p>
Emulation	MOVA @SP+,PC
Description	The 20-bit return address information, pushed onto the stack by a CALLA instruction, is restored to the PC. The program continues at the address following the subroutine call. The SR bits SR.11:0 are not affected. This allows the transfer of information with these bits.
Status Bits	<p>N: Not affected</p> <p>Z: Not affected</p> <p>C: Not affected</p> <p>V: Not affected</p>
Mode Bits	OSCOFF, CPUOFF, and GIE are not affected.
Example	Call a subroutine SUBR from anywhere in the 20-bit address space and return to the address after the CALLA

```

CALLA    #SUBR        ; Call subroutine starting at SUBR
...
SUBR     PUSHM.A     #2,R14 ; Save R14 and R13 (20 bit data)
...
        POPM.A      #2,R14 ; Restore R13 and R14 (20 bit data)
        RETA        ; Return (to full address space)

```

1.6.4.10 SUBA

SUBA	Subtract 20-bit source from 20-bit destination register
Syntax	SUBA Rsrc,Rdst SUBA #imm20,Rdst
Operation	$(\text{.not.src}) + 1 + \text{Rdst} \rightarrow \text{Rdst}$ or $\text{Rdst} - \text{src} \rightarrow \text{Rdst}$
Description	The 20-bit source operand is subtracted from the 20-bit destination register. This is made by adding the 1s complement of the source + 1 to the destination. The result is written to the destination register, the source is not affected.
Status Bits	N: Set if result is negative ($\text{src} > \text{dst}$), reset if positive ($\text{src} \leq \text{dst}$) Z: Set if result is zero ($\text{src} = \text{dst}$), reset otherwise ($\text{src} \neq \text{dst}$) C: Set if there is a carry from the MSB (Rdst.19), reset otherwise V: Set if the subtraction of a negative source operand from a positive destination operand delivers a negative result, or if the subtraction of a positive source operand from a negative destination operand delivers a positive result, reset otherwise (no overflow)
Mode Bits	OSCOFF, CPUOFF, and GIE are not affected.
Example	The 20-bit value in R5 is subtracted from R6. If a carry occurs, the program continues at label TONI.
	<pre> SUBA R5,R6 ; R6 - R5 -> R6 JC TONI ; Carry occurred ... ; No carry </pre>

1.6.4.11 TSTA

*** TSTA** Test 20-bit destination register

Syntax TSTA Rdst

Operation dst + 0FFFFFFh + 1
dst + 0FFFFFFh + 1
dst + 0FFh + 1

Emulation CMPA #0, Rdst

Description The destination register is compared with zero. The status bits are set according to the result. The destination register is not affected.

Status Bits
N: Set if destination register is negative, reset if positive
Z: Set if destination register contains zero, reset otherwise
C: Set
V: Reset

Mode Bits OSCOFF, CPUOFF, and GIE are not affected.

Example The 20-bit value in R7 is tested. If it is negative, continue at R7NEG; if it is positive but not zero, continue at R7POS.

```

TSTA   R7           ; Test R7
JN     R7NEG        ; R7 is negative
JZ     R7ZERO       ; R7 is zero
R7POS  .....       ; R7 is positive but not zero
R7NEG  .....       ; R7 is negative
R7ZERO .....       ; R7 is zero

```

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